

## CSE 461: Bridging LANs

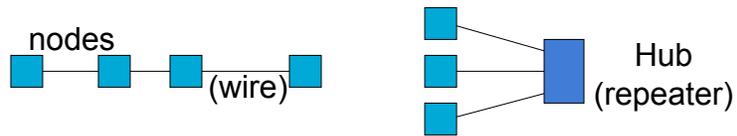
### Next Topic -- Switching (a.k.a. Bridging)

- Focus:
  - What to do when one shared LAN isn't big enough?
- Interconnecting LANs
  - Bridges and LAN switches
  - A preview of the Network layer

Application
Presentation
Session
Transport
Network
Data Link
Physical

## Limits of a LAN

- One shared LAN can limit us in terms of:
  - Distance
  - Number of nodes
  - Performance



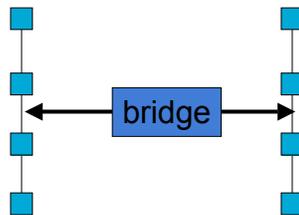
- How do we scale to a larger, faster network?
  - We must be able to interconnect LANs

## Switching (a.k.a. Bridging)

- Transferring a packet from one LAN to another LAN
  - Build an "extended LAN"
- Different varieties of switching
  - Packet switched vs. circuit switched
  - Connection vs. Connectionless
- We'll focus on connectionless, packet switched
  - Ethernet

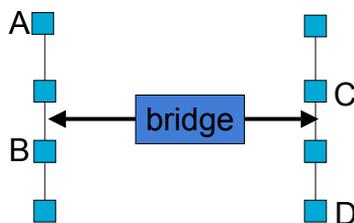
## Bridges and Extended LANs

- “Transparently” interconnect LANs with bridge
  - Receive frames from each LAN and forward to the other
  - Each LAN is its own collision domain; bridge isn’t a repeater
  - Could have many ports



## Learning Bridges

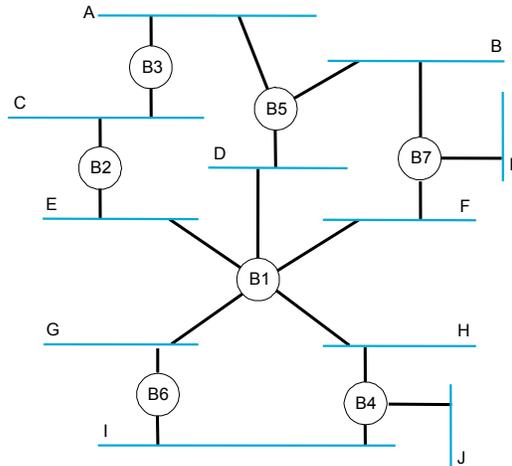
- To optimize overall performance:
  - Shouldn’t forward  $A \rightarrow B$  or  $C \rightarrow D$ , should forward  $A \rightarrow C$  and  $D \rightarrow B$



- How does the bridge know?
  - Learn who is where by observing source addresses and prune
  - Forward using destination address; age for robustness

## Why stop at one bridge?

- Allows you to incrementally build out network, across organizations
- What problems could arise?
- How to solve them?

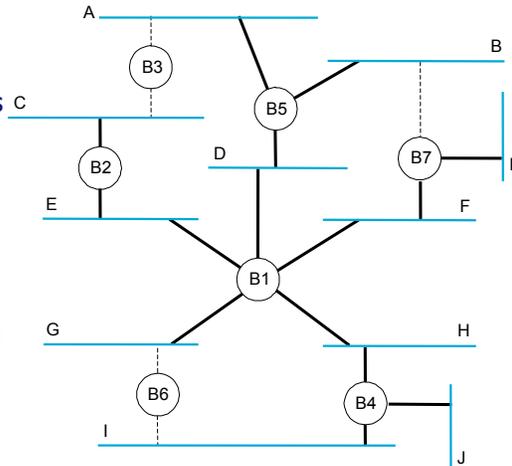


## High Level Outline of a Solution

- Take the set of networks and turn off some of the ports in the bridges
  - Goal is to have the bridges form a spanning tree
  - Spanning tree means that there are no cycles
- When a bridge doesn't know where is the destination of the packet:
  - It floods the packet
  - Flooding takes place along the spanning tree
- Over time, bridges learn which links of the spanning tree can be used to route to a given destination

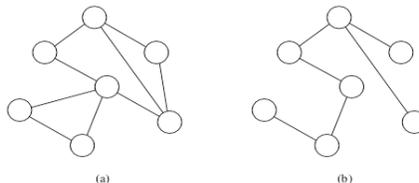
## Spanning Tree Example

- Spanning tree uses select bridge ports so there are no cycles
  - Prune some ports
  - Only one tree
- Q: How do we find a spanning tree?
  - Automatically with a distributed algorithm



## Spanning Tree

- Compute ST with a bridge as *root* such that
  - Root forwards onto all of its outgoing ports
  - Other bridges forward TO the root if a packet is coming from a bridge further from the root, else they forward away from the root
    - Packet traversal: forwards (UP)\* then (DOWN)\*



## Spanning tree with learning

---

- Once the spanning tree is in place...
  - the bridge uses the regular learning algorithm to figure out which ports to forward / flood packet on
- Job of spanning tree algorithm is to disable some ports to eliminate cycles

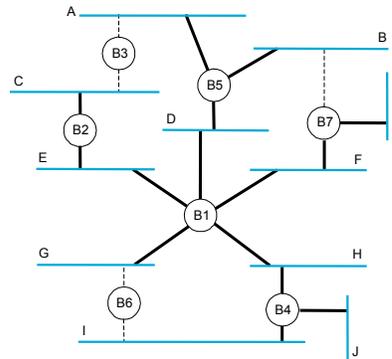
## Spanning Tree Algorithm

---

- Distributed algorithm to compute spanning tree
  - Robust against failures, needs no organization
  - Developed by Radia Perlman at DEC
    - IEEE 802.1 spec
    - <http://www1.cs.columbia.edu/~ji/F02/ir02/p44-perlman.pdf>
- Outline: Goal is to turn some bridge ports off
  1. Elect a root node of the tree (lowest address)
  2. Grow tree as shortest distances from the root (using lowest address to break distance ties)
    - All done by bridges sending periodic configuration messages over ports for which they are the "best" path
    - Then turn off ports that aren't on "best" paths

## Algorithm Overview

- Each bridge has a unique id
  - e.g., B1, B2, B3
- Select the bridge with the smallest id as root
- Select bridge on each LAN that is closest to the root as that LAN's designated bridge
  - use ids to break ties
- Each bridge forwards frames over each LAN on which it is the designated bridge



## Algorithm continued

- Bridges exchange configuration messages, containing:
  - id for bridge sending the message
  - id for what the sending bridge believes to be the root bridge
  - distance (hops) from sending bridge to root bridge
- Each bridge records current best configuration message for each port

## Algorithm continued

---

- Initially, each bridge believes it is the root
  - when learn not root, stop generating configuration messages
  - instead, forward root's configuration message
    - incrementing distance field by 1
  - in steady state, only root generates configuration messages

## Algorithm More...

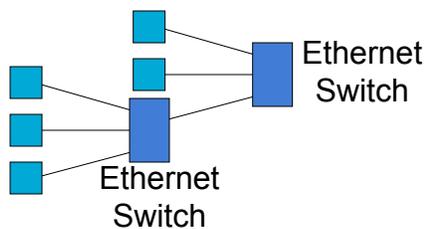
---

- When bridge learns it is not designated bridge on LAN, stop forwarding configuration messages
  - in steady state, only designated bridges forward configuration messages
- Root bridge continues to send configuration messages periodically
- If a bridge does not receive config. message after a period of time:
  - assumes topology has changed
  - starts generating configuration messages claiming to be root



## LAN Switches

- LAN switches are multi-port bridges
  - Modern, high performance form of bridged LANs
  - Looks like a hub, but frames are switched, not shared
  - Every host on a separate port, or can combine switches



## Limitations of Bridges/Switches

- LAN switches form an effective small-scale network
  - Plug and play for real!
- Why can't we build a large network using bridges?
  - Little control over forwarding paths
  - Size of bridge forwarding tables grows with number of hosts
  - Broadcast traffic flows freely over whole extended LAN
  - Spanning tree algorithm limits reconfiguration speed

## Key Concepts

---

- We can overcome LAN limits by interconnection
  - Bridges and LAN switches
  - But there are limits to this strategy ...