

2.

The network is saturated if there are 10 or more people sending. The cumulative probability is obtained by computing the probability that EXACTLY i nodes are transmitting, and then summing them up over all $i \geq 10$

$$\text{So Pr[exactly } i \text{ of the } n \text{ nodes transmit]} = \binom{n}{i} * (p^i) * (1-p)^{(n-i)}$$

where p = probability that a node transmits

3. Book Questions

1.5

$$\begin{aligned} \text{(a) total_time} &= (\text{handshake}) + (\text{propagation delay}) + (\text{transfer time}) \\ &= (2*100)\text{ms} + (100/2)\text{ms} + \frac{(1000*(10^3)*8) \text{ bits}}{1.5 * (10^6) \text{ bps}} = 5.583 \text{ sec} \end{aligned}$$

(b) The total time will be the same as above, except that we wait for an extra 999 RTTs, while sending the 1000 packets.

$$\text{So, total_time} = 5.583 + 999*(0.1) = 105.483 \text{ sec}$$

(c) We send 20 pkts in 1 RTT => 1000 pkts can be sent in 50 RTTs
total_time = (handshake) + 50*RTT
= 52*RTT = 5.2 sec

(d) Num pkts sent out till k th RTT = $(2^k) - 1$
So, num RTTs required for 1000 pkts = $\text{ceil}(\log_2(1000+1)) = 10$ RTTs
total_time = (handshake) + 10*RTT
= 12*RTT = 1.2 sec

Note: The answer for (c) and (d) may be different depending on the way in which the propagation delay is considered.

1.6

$$\begin{aligned} \text{(a) total_time} &= (\text{handshake}) + (\text{propagation delay}) + (\text{transfer time}) \\ &= (2*80)\text{ms} + (80/2)\text{ms} + \frac{(1500*(10^3)*8) \text{ bits}}{10 * (10^6) \text{ bps}} = 1.4 \text{ sec} \end{aligned}$$

(b) The total time will be the same as above, except that we wait for an extra 1499 RTTs, while sending the 1500 packets.

$$\text{So, total_time} = 1.4 + 1499*(0.08) = 121.32 \text{ sec}$$

(c) We send 20 pkts in 1 RTT => 1500 pkts can be sent in 75 RTTs
total_time = (handshake) + 75*RTT
= 77*RTT = 6.16 sec

(d) Num pkts sent out till k th RTT = $(2^k) - 1$
So, num RTTs required for 1500 pkts = $\text{ceil}(\log_2(1500+1)) = 11$ RTTs
total_time = (handshake) + 11*RTT
= 13*RTT = 1.04 sec

1.7

$$\text{Propagation delay} = \frac{2*(10^3)\text{km}}{2*(10^8)\text{km/sec}} = 10 \text{ microseconds}$$

$$\begin{aligned} 100 \text{ byte packets: } 10^{*-5} &= (100*8 \text{ bits})/(B \text{ bps}) \\ \Rightarrow B &= 8*(10^7)\text{bps} = 80 \text{ Mbps} \end{aligned}$$

$$\begin{aligned} 512 \text{ byte packets: } 10^{*-5} &= (512*8 \text{ bits})/(B \text{ bps}) \\ \Rightarrow B &= 4096*(10^5)\text{bps} = 409.6 \text{ Mbps} \end{aligned}$$

1.13

$$\begin{aligned} \text{Width of a bit} &= (1 \text{ bit}) / (10^{-9} \text{ bps}) = 10^{*-9} \text{ sec} = 1 \text{ ns} \\ \text{Length of the bit} &= 2.3*(10^8) \text{ m/sec} * 1 \text{ ns} = 0.23 \text{ m} \end{aligned}$$

2.16

To find the new checksum:

Complement the old checksum,
Subtract 1 or 256 (depending on lower or higher order byte decrement)
Complement the result

Or equivalently: Take the old checksum, and add 1 or 256 : this gives the new checksum :)