

CSE/EE 461 – Lecture 6

Wireless and Contention-Free Protocols

Last Time ...

- The multi-access problem
 - Medium Access Control (MAC) sublayer
- Random access protocols:
 - Aloha
 - CSMA variants
 - Classic Ethernet (CSMA/CD)

| |
|--------------|
| Application |
| Presentation |
| Session |
| Transport |
| Network |
| Data Link |
| Physical |

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This Lecture

More on multiple-access schemes:

1. Wireless schemes
2. Contention-free protocols

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|--------------|
| Application |
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| Data Link |
| Physical |

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1. Wireless Communication

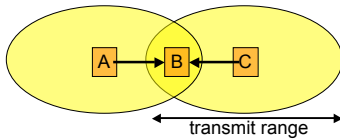
Wireless is more complicated than wired ...

1. Cannot detect collisions
 - Transmitter swamps co-located receiver
2. Different transmitters have different coverage areas
 - Asymmetries lead to hidden/exposed terminal problems

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Hidden Terminals

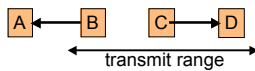


- A and C can both send to B but can't hear each other
 - A is a hidden terminal for C and vice versa
- CSMA will be ineffective - want to sense at receiver

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Exposed Terminals



- B, C can hear each other but can safely send to A, D
- Compare to spatial reuse in cell phones:



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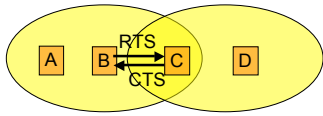
CSMA with Collision Avoidance

- Since we can't detect collisions, we avoid them
 - CSMA/CA as opposed to CSMA/CD
 - Not greedy like Ethernet
- CS: listen before transmitting.
 - When medium busy, choose random backoff interval
 - Wait for that many idle timeslots to pass before sending
- CA: transmit short "jamming" signal before sending frame
 - essentially reserves medium, let's others know your intent to transmit
- Collisions can be inferred
 - Use CRC and ACK from receiver to infer "no collision"
 - on collision, binary exponential backoff like Ethernet

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RTS / CTS Protocols (MACA)



1. B stimulates C with Request To Send (RTS)
2. A hears RTS and defers to allow the CTS
3. C replies to B with Clear To Send (CTS)
4. D hears CTS and defers to allow the data
5. B sends to C

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802.11 Wireless LANs

- Emerging standard with a bunch of options/features ...



- Wireless plus wired system or pure wireless (ad hoc)
- Avoids collisions (CSMA/CA (p-persistence), RTS/CTS)

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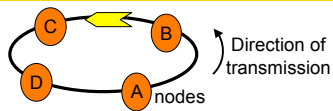
2. Contention-free Protocols

- Collisions are the main difficulty with random schemes
 - Inefficiency, limit to scalability
- Q: Can we avoid collisions?
- A: Yes. By taking turns or with reservations
 - Token Ring / FDDI, DQDB
- More generally, what else might we want?
 - Deterministic service, priorities/QOS, reliability

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Token Ring (802.5)



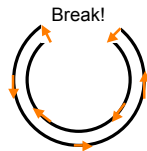
- Token rotates permission to send around node
- Sender injects packet into ring and removes later
 - Maximum token holding time (THH) bounds access time
 - Early or delayed token release
 - Round robin service, acknowledgments and priorities
- Monitor nodes ensure health of ring

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FDDI (Fiber Distributed Data Interface)

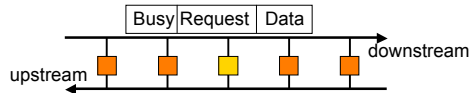
- Roughly a large, fast token ring
 - 100 Mbps and 200km vs 4/16 Mbps and local
 - Dual counter-rotating rings for redundancy
 - Complex token holding policies for voice etc. traffic
- Token ring advantages
 - No contention, bounded access delay
 - Supports fair, reserved, priority access
- Disadvantages
 - Complexity, reliability, scalability



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DQDB (Distributed Queue Dual Bus)



- Two unidirectional buses that carry fixed size cells
 - Cells are marked busy/free and can signal a request too
- Nodes maintain a distributed FIFO queue
 - By sending requests they are reserving future access

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DQDB Algorithm

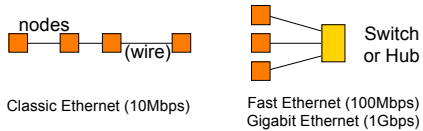
- Two counters per direction (UP, DN)
 - RC (request count), CD (countdown)
- Consider sending downstream (DN):
 - Always have RC count UP requests, minus free DN cells if larger than zero
 - This is a measure of how many others are waiting to send
 - To send, copy RC to CD and set RC to zero, then decrement CD for each free DN cell, send when zero
 - This waits for earlier requests to be satisfied before sending
- Highly scalable, efficient, but not perfectly fair

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Modern Ethernet

- A key concern is manageability
 - centralized vs. distributed layout
- Another is performance scalability
 - Switches vs. Hubs



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Key Concepts

- Wireless communication is relatively complex
 - No collision detection, hidden and exposed terminals
- There are contention-free MAC protocols
 - Based on turn taking and reservations, not randomization
