

## CSE/EE 461 – Lecture 3

### Bits, Encoding and Framing

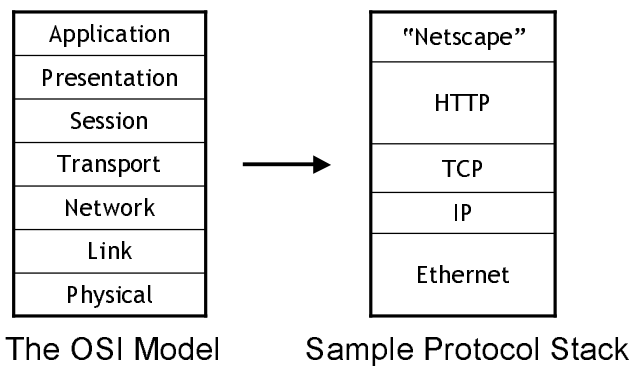
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### Last Time ...

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- Protocols, layering and reference models



## This Lecture

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- A look at the physical/link layers:
  1. Different kinds of Media
  2. Encoding bits with signals
  3. Framing messages
- Key Focus: How do we send a message across a wire?

Application
Presentation
Session
Transport
Network
Data Link
Physical

## 1. Different Kinds of Media

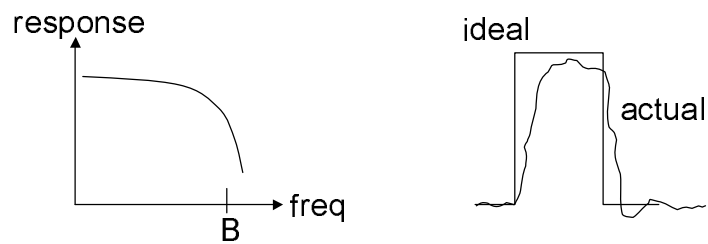
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- Wire
  - Twisted pair, e.g., CAT5 UTP, 10 → 100Mbps, 100m
  - Coaxial cable, e.g, thin-net, 10 → 100Mbps, 200m
- Fiber
  - Multi-mode, 100Mbps, 2km
  - Single mode, 100 → 2400 Mbps, 40km
- Wireless
  - Infra-red, e.g., IRDA, ~1Mbps
  - RF, e.g., 802.11 wireless LANs, Bluetooth (2.4GHz)
  - Microwave, satellite, cell phones, ...

## Wires

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- Signal subject to:
  - Attenuation (repeaters)
  - Distortion (frequency and delay)
  - Noise (thermal, crosstalk, impulse)



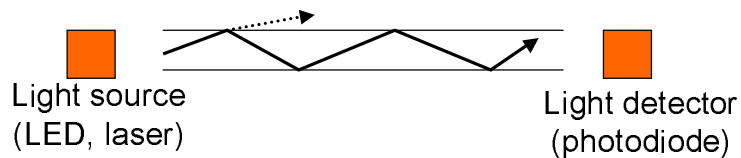
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## Fiber

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- Long, thin, pure strand of glass
  - Enormous bandwidth available (terabits)



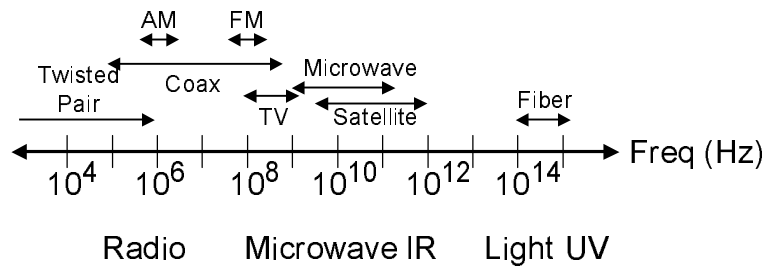
- Multi-mode allows many different paths, dispersion
- Chromatic dispersion if multiple frequencies

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# Wireless

- Different frequencies have different properties
- Signals subject to atmospheric/environmental effects



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## 2. Encoding Bits with Signals

- Generate analog waveform (e.g., voltage) from digital data at transmitter and sample to recover at receiver



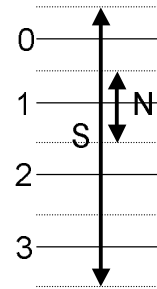
- We send/recover symbols that are mapped to bits
  - Signal transition rate = baud rate, versus bit rate
- This is baseband transmission ... take a signals course!

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## Aside: Bandwidth of a Channel

- EE: Bandwidth (B) (Hz) is the width of the pass-band in the frequency domain
- CS: "Bandwidth" (bps) is the information carrying capacity (C) of the channel
- Shannon showed how they are related by noise, which limits how many signal levels we can safely distinguish.



S=signal  
N=noise

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## The Shannon Limit (1948)

- Define Signal to Noise Ratio (SNR):  
$$\text{SNR} = 10\log_{10}(\text{signal} / \text{noise}) \text{ decibels (dB)}$$

e.g, 30 dB means signal 1000 times noise
- For a noisy channel with bandwidth B (Hz) and given SNR, the maximum rate at which it is possible to send information, the channel capacity, is:  
$$C = B \log_2(1 + \text{SNR}) \text{ (bits/sec)}$$

e.g 3KHz and 30dB SNR  $\rightarrow$  30Kbps

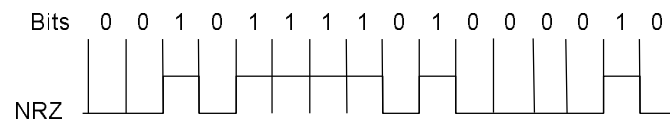
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## NRZ and NRZI

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- Simplest encoding, NRZ (Non-return to zero)
  - Use high/low voltages, e.g., high = 1, low = 0
- Variation, NRZI (NRZ, invert on 1)
  - Use transition for 1s, no transition for 0s



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## Clock Recovery

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- Problem: How do we distinguish consecutive 0s or 1s?
- If we sample at the wrong time we get garbage ...
- If sender and receiver have exact clocks no problem
  - But in practice they drift slowly
- This is the problem of clock recovery
  
- Possible solutions:
  - Send separate clock signal → expensive
  - Keep messages short → limits data rate
  - Embed clock signal in data signal → other codes

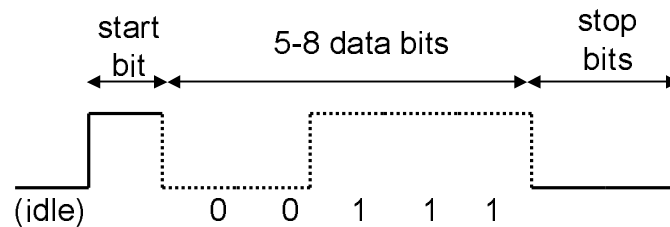
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## "Asynchronous" Transmission

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- Avoid timing problem by sending short, delimited data
  - E.g., UARTs (typically used to connect your keyboard)



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## Manchester Coding

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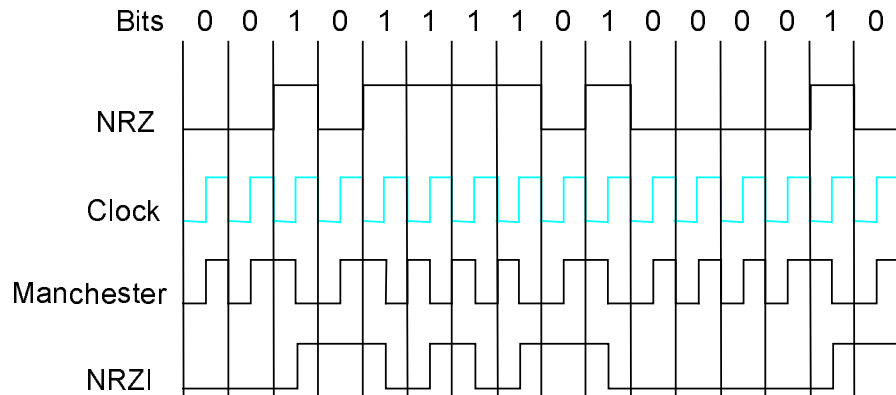
- Make transition in the middle of every bit period
  - Low-to-high is 0; high-to-low is 1
  - Signal rate is twice the bit rate
  - Used on 10 Mbps Ethernet
- Advantage: self-clocking
- Disadvantage: 50% efficiency

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## Coding Examples

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## 4B/5B Codes

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- We want transitions \*and\* efficiency ...
- Solution: map data bits (which may lack transitions) into code bits (which are guaranteed to have them)
- 4B/5B code:
  - 0000  $\rightarrow$  11110, 0001  $\rightarrow$  01001, ... 1111  $\rightarrow$  11101
  - Never more than three consecutive 0s back-to-back
  - 80% efficiency
- This code is used by LANs such as FDDI

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### 3. Framing

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- Need to send message, not just bits
  - Even if we know where the bits are we still need to synchronize on the start of the message
  - Complete Link layer messages are called frames
- Common approach: Sentinels
  - Look for special control code that marks start of frame

### Point-to-Point Protocol (PPP)

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- IETF standard, used for dialup and leased lines

Flag 01111110	(header)	Payload (variable)	(trailer)	Flag 01111110
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- Flag indicates start/end of frame
- Occurrences of flag inside payload must be "stuffed"
  - Replace "flag" with "flag flag"
  - Length of payload is data-dependent!

## Other Approaches

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- Use explicit byte count after preamble
  - More susceptible to errors?
- Use “invalid” codeword
  - E.g., pick non-data 4B/5B symbol, used for FDDI
- SONET: “clock”-based framing
  - Periodic sync bits plus very accurate clock
  - Used extensively in the telecommunications industry

## Key Concepts

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- Different media have different properties that affect higher layer protocols
- To send messages in practice we must solve the problems of clock recovery and framing