

## Image Scissors

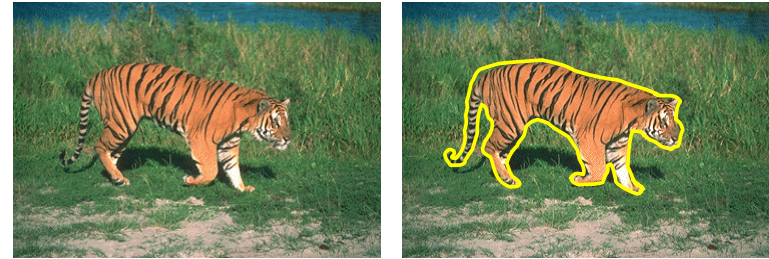


[Aging Helen Mirren](#)

### Today's Readings

- [Intelligent Scissors](#), Mortensen et. al, SIGGRAPH 1995

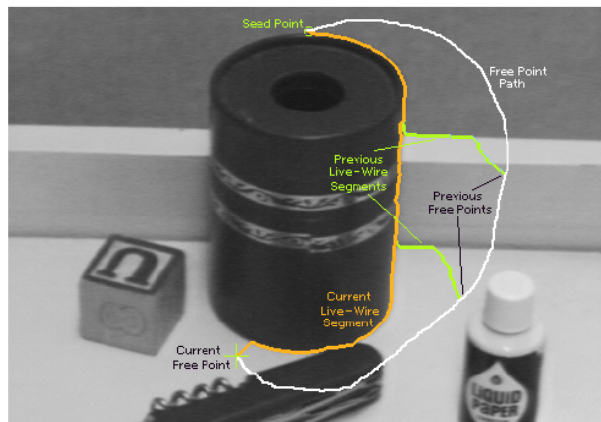
## Extracting objects



How could this be done?

- hard to do manually
- hard to do automatically ("image segmentation")
- easy to do *semi-automatically*

## Intelligent Scissors (demo)

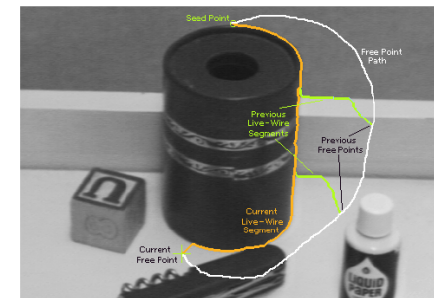


**Figure 2:** Image demonstrating how the live-wire segment adapts and snaps to an object boundary as the free point moves (via cursor movement). The path of the free point is shown in white. Live-wire segments from previous free point positions ( $t_0$ ,  $t_1$ , and  $t_2$ ) are shown in green.

## Intelligent Scissors

Approach answers a basic question

- Q: how to find a path from seed to mouse that follows object boundary as closely as possible?
- A: define a path that stays as close as possible to edges

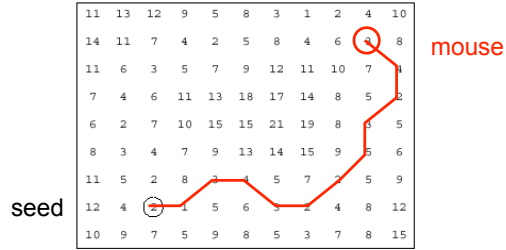


**Figure 2:** Image demonstrating how the live-wire segment adapts and snaps to an object boundary as the free point moves (via cursor movement). The path of the free point is shown in white. Live-wire segments from previous free point positions ( $t_0$ ,  $t_1$ , and  $t_2$ ) are shown in green.

# Intelligent Scissors

## Basic Idea

- Define edge score for each pixel
  - edge pixels have low cost
- Find lowest cost path from seed to mouse

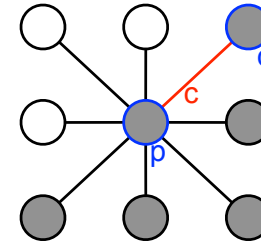


## Questions

- How to define costs?
- How to find the path?

# Let's look at this more closely

Treat the image as a graph



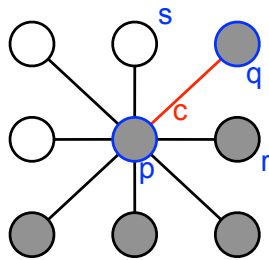
## Graph

- node for every pixel **p**
- link between every adjacent pair of pixels, **p,q** (8-connected)
- cost **c** for each link

Note: each *link* has a cost

- this is a little different than the figure before where each pixel had a cost

# Defining the costs

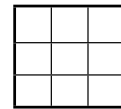
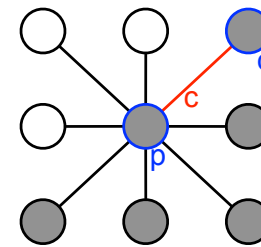


Want to hug image edges: how to define cost of a link **c**?

- good (low-cost) links follow the intensity edge
  - want intensity to change rapidly  $\perp$  to the link

- $c \approx -\frac{1}{\sqrt{2}} |\text{intensity of } r - \text{intensity of } s|$

# Defining the costs

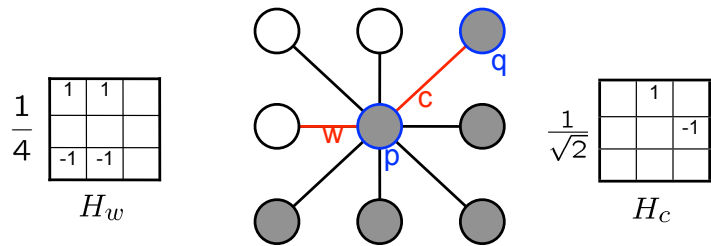


$H_c$

**c** can be computed using a cross-correlation filter

- assume it is centered at **p**

## Defining the costs



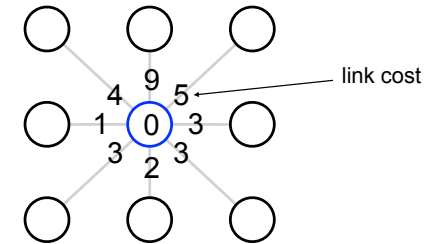
$c$  can be computed using a cross-correlation filter

- assume it is centered at  $p$

A couple more modifications

- Scale the filter response by length of link  $c$ . Why?
- Make  $c$  positive
  - Set  $c = (\max\text{-|filter response|} \cdot \text{length})$
  - where max = maximum |filter response| \* length over all pixels in the image

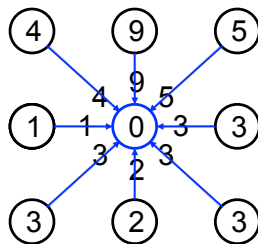
## Dijkstra's shortest path algorithm



Algorithm

1. init node costs to  $\infty$ , set  $p$  = seed point,  $\text{cost}(p) = 0$
2. expand  $p$  as follows:
  - for each of  $p$ 's neighbors  $q$  that are not expanded
    - » set  $\text{cost}(q) = \min(\text{cost}(p) + c_{pq}, \text{cost}(q))$

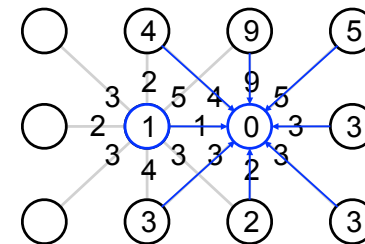
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3. set  $r$  = node with minimum cost on the ACTIVE list

## Dijkstra's shortest path algorithm

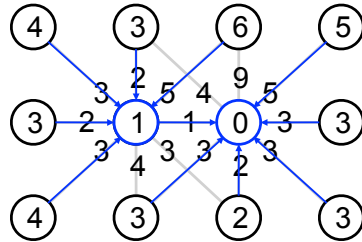


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## Dijkstra's shortest path algorithm

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## Dijkstra's shortest path algorithm

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### Properties

- It computes the minimum cost path from the seed to every node in the graph. This set of minimum paths is represented as a *tree*
- Running time, with  $N$  pixels:
  - $O(N^2)$  time if you use an active list
  - $O(N \log N)$  if you use an active priority queue (heap)
  - takes fraction of a second for a typical (640x480) image
- Once this tree is computed once, we can extract the optimal path from any point to the seed in  $O(N)$  time.
  - it runs in real time as the mouse moves
- What happens when the user specifies a new seed?

## Results

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[Jason Dang](#)



[Peter Davis](#)

<http://www.cs.washington.edu/education/courses/455/12wi/projects/project1/artifacts/index.html>

## Help session—Avanish

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