

# The Two "Auth-" Operations

#### Authentication

- "Process of accepting *credentials* from a user and *validating* those
- credentials against some *authority*" - The result is an authenticated identity
- Authorization
  - "Process of *determining* whether the authenticated *identity* has *access* to a given resource"

Both steps follow this order and both are essential!

# What Can Go Wrong? Authentication breaks if:

- Credentials are forged
- Authority is subverted
- Validating function is replaced

#### Authorization breaks if:

- Authentication identity is forged
- Access matrix is tampered with
- Matrix lookup function is replaced

• Lesson: Security needs to be provisioned on each step!

## Types of Authentication

#### • Server authentication

- Necessary in e-commerce
- Achieved via:
- X.509 certificates, signed by known certificate authorities (CA)
- · Digital signatures using public/private key encryption

#### • Client authentication

- Necessary in e-commerce
- Majority of clients typically do not use X.509 certificates, or public/private key pairs
  - How many of you use one of these methods for authentication?

# How to Evaluate Proposed Approaches?

#### Ask:

- 1. What problem is the approach trying to solve?
- 2. What are the ways in which the approach can fail (including, be deliberately made to fail)?
- 3. Given the ways the approach can fail, does it really solve the problem at hand?
- 4. What are the costs (financial and otherwise) of deploying a real implementation of the approach?
- 5. Given the failure conditions and costs, is it worthwhile?

## **Client Authentication Methods**

- **Client certificates**
- No incentive for clients to have one  $\Rightarrow$  not widely deployed
- **Digital signatures** 
  - No PKI yet  $\Rightarrow$  hard to safely distribute public keys
- Passwords
  - Most primitive, pervasive method
  - Easy to use, easy to crack: passwords are guessable or users do not remember them
    - · Copy-and-store-in-wallet works well in practice with random passwords
    - Visual passwords random art; a drawing in lieu of a word

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<u>SKey protocol</u> - changing passwords on every communication <u>Smart cards</u> - store random password safely; PIN for theft protection; activated only by a special card reader; European invention

## **Client Authentication Methods**

#### • **Biometrics**

- Unique, inherently tied to the individual
- But:
  - · Fingerprinting non-permanent, could be tampered with
  - · Retina scans non-permanent, invasive, even dangerou
  - Face recognition high false positives rate, could be easily fooled
     Voice recognition high false positive and false negative rate, recordable
  - · DNA analysis slow, extremely invasive, may be non-permanent
  - (<u>Normal) Signature</u> varies widely (high false negative rate), more appropriate for non-repudiation that authentication

## Client Authentication on the Web

- What assumptions / constraints does the Web environment imply?
- Which of the above methods are unsuitable for authentication on the Web?
- What remains?

## **Motivation**

- Growing need for personalized, access-controlled Web-based services
  - E.g.: nytimes.com, myuw.washington.edu, hotmail.com
- Some popular authentication mechanisms not suitable for the Web environment
  - Designed for long-running connections
  - Involve expensive computations public/private key crypto - Authentication identities can be replayed - biometrics
- Developers lack proper background in security
- **Result: Proliferation of home-grown weak** authentication schemes

# Limitations on Web Authentication Schemes

- Must use only widely deployed, portable and lightweight technologies
- No smart cards or client certificates; JavaScript may be ok
- Must require minimum user involvement - No password re-typing or perpetual dialog boxes
- Must not unduly overload servers with expensive computations
- No public-key crypto; cryptographic hashes are fine
- Must store client state in a very limited space - E.g.: cookies on the client, (maybe) a database on server

# Not All Web Authentication Schemes Are Created Equal

#### **Designs differ depending on:**

- Type of service
  - General subscription
  - · Online newspapers and libraries
  - User customization · Online identities, per-user content filtering
- Security needs
  - Sensitivity of the client data
  - Store data on server and put an index to it in a client cookie
  - Load tolerance on the server Delicate tradeoff with clients' need for strong protection

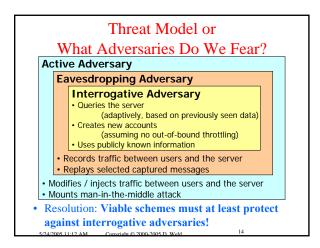
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# Threat Model or What Attacks Do We Fear?

- Forging\* an authentication token for

   A random user (a.k.a. existential forgery)
   Useful for free access to subscription services
  - A chosen user (a.k.a. selective forgery)
    Allows access to data for any selected user
  - All users (a.k.a. total break)
     Allows forging tokens for all users at any time

\* forging ≠ replay attack



# Hints for Designing Client Authentication Schemes

Disclaimer: Hints are useful, but following them is neither necessary, nor sufficient for security

# Hints: Use Cryptography Appropriately

- Using crypto is inescapable if you want to protect from adversaries!
- <u>Hint #1</u>: Assess your needs for protection – Tradeoffs between usability and complexity
- <u>Hint #2: Choose a "tried and true" existing scheme</u> – Home-grown schemes are almost always trivial to break

# Hints: Use Cryptography Appropriately

If you absolutely must design your own scheme:

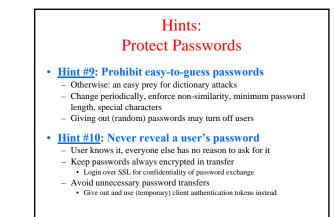
- <u>Hint #3</u>: Think twice! Ask those who know better!
  Hint #4: Have it reviewed by security experts
- <u>Hint #4</u>: Have it reviewed by security experts
   Announcing it loudly is good but not sufficient
- <u>Hint #5</u>: Keep the scheme simple – Makes it easier to analyze for security
- <u>Hint #6a</u>: Do not rely on the secrecy of the protocol
   Gives you false sense of security until someone figures it out
- Hint #6b: Instead, rely on the secrecy of keys

# Hints: Use Cryptography Appropriately

- <u>Hint #7</u>: Understand the properties and details of crypto primitives you use
  - Many provide some assurances, but not other (e.g., SSL)
     Many make fine-print assumptions
  - UNIX crypt() hash function truncates input beyond 8 characters
- <u>Hint #8</u>: Avoid composing security schemes
   May weaken the composite, even if secure in isolation
   • E.g., using the same secret key for multiple purposes

## Status on Using Passwords

- Users don't want passwords
  - Tradeoff between usability and security
  - Users tend to pick poor (easy) passwords
     Do not suggest ideas they will blindly follow it
- Users tend to reuse passwords across many sites – How many different passwords do you use?
  - How many of them do you commit to memory?
  - How many of them do you have written somewhere (as a backup)?
- Compromising a password leads to impersonation



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# Hints:

# Protect Passwords

 <u>Hint #11</u>: Redo authentication before securitysensitive operations

- E.g.: changing passwords
- Avoids attacks through replayed authentication tokens

# Hints: Handle Authentication Tokens Wisely <u>Hint #12</u>: Avoid predictable authentication tokens

- E.g.: publicly available info, sequential ID numbers, etc.
- <u>Hint #13</u>: Protect tokens from tampering – Tokens may contain sensitive user info
  - Use only strong cryptographic hash functions (e.g., no CRC)
     Use a keyed message digest (e.g., MAC, no MD5)
- <u>Hint #14</u>: If combining multiple data into a token, separate components unambiguously

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Avoids a splicing attack:
"Alice" • "213" • "Bob" == "Alice2" • "13" • "Bob"

## Hints: Handle Authentication Tokens Wisely

#### • <u>Hint #15</u>: Encrypt tokens

- For tokens stored in cookies and sent over SSL, set Secure flag
- Prevents eavesdroppers from capturing and replaying tokens
- <u>Hint #16</u>: Do not include a token as part of a URL
   Otherwise, token may leak through plaintext channels
  - E.g.: cross-site scripting attack using the HTTP Referer field

#### • <u>Hint #17</u>: Avoid using persistent cookies

- If cookie (file) is leaked, attacker can impersonate user
- Can users defend against this threat (the authentication scheme designer may have been negligent)?

# Hints: Handle Authentication Tokens Wisely

- Hint #18: Make authentication tokens expire:

   Store a tamper-resistant timestamp in cookie, or keep token expiration time on the server
   Limits the potential damage in case a token leaks out
- Hint #19: Do not trust the client...
   ... to enforce token expiration (manipulating a cookie is easy)
   ... (in general) for anything that the client can possibly forge
- <u>Hint #20</u>: To prevent replays of leaked tokens:
   \_ Keep tokens confidential and mint new ones after each use
  - Bind tokens to network addresses
     But DHCP users' tokens may expire prematurely

# Sample Authentication Scheme

- Goals
  - Statelessly verify authenticity of request and its contents
  - Explicitly control lifetime of token Portability
- Design choice
  - Authentication cookies
  - · Anyone with a valid cookie has access to protected server content
- Claim

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- Secure against an interrogative adversary

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- If layered over SSL with server authentication, secure against an active adversary

## **Cookie Basics**

- HTTP is a stateless protocol
- Client IDs generated by server, stored on client
- Sent back to server with subsequent requests
- Cookie attributes:
  - Data used to uniquely identify client
  - Domain cookie only applies to this server domain
  - Path server path
  - Secure flag should cookie data be encrypted? - Expiration - current session or physical time

Suggested Cookie Structure

#### exp=t&data=s&digest=MAC<sub>k</sub>(exp=t&data=s)

 $t \rightarrow expiration time (seconds past 1970 GMT)$  $s \rightarrow$  data, associated with the client  $k \rightarrow$  server secret kev

 $MAC \rightarrow$  strong cryptographic hash function

 $HMAC_{k}(M) ::= H(k \oplus 0x5c \bullet H(k \oplus 0x36 \bullet M))$ 

where  $H \in \{SHA1, MD5\}$ , M is the message

# Disecting the Scheme

#### • Expiration time:

- Avoids keeping server state
- Tradeoff between potential damage and frequent
- reauthentication (security vs. usability) - Should users be allowed to control it?

#### • Data:

- Sensitive data should not be stored here · If needed, store cryptographically random session ID, while keeping important data on server
- Balance between respecting users' privacy and saving server resources
  - · Likely to be biased in favor of the latter

## Disecting the Scheme

#### Key:

- Recommended length is twice that of block encryption ciphers (~160 bits or more)
  - · Fends off birthday attacks

## Disecting the Scheme

#### **Strengths:**

- Simplicity
- Authenticating clients:
  - Requires O(1) server state (for the key)
- Takes O(1) time
- Would depend on number of clients if server state were kept
- Easier to deploy multiserver systems
  - No need for dynamically shared data between servers

## Disecting the Scheme

#### Weaknesses:

- · Server is vulnerable against colluding clients
  - Clients more likely to share temporary tokens than passwords
  - How many other people's passwords do you know?
- No mechanism for selective secure token revocation
  - Unnecessary for short sessions Separation of policy and mechanism? - If needed, keep session status on server
    - Yahoo does it
    - But, allows simultaneous revocation of all tokens · By changing the secret server key

## Security Analysis

#### Strength of authentication scheme depends on:

- Strength of MAC function
- · Secrecy of server key
- Strength of server key and frequency of changing it - Longer keys adversely affect performance of hash functions
- · Strength of client passwords against guessing and dictionary attacks

## **Performance Factors**

#### • HMAC-SHA1

- 1.2 ms / request
- Runs on small chunks of data
- SSL
  - 90 ms / request
  - Runs on the entire HTTP stream
  - New connections are costly to setup, session resumption helps

## **Other Authentication Schemes**

#### • HTTP Basic Authentication

- Sends username and password repeatedly in cleartext Falls prey to eavesdropping adversaries
   dsniff - automated tool for sniffing authentication exchanges
- HTTP Digest Authentication
  - Encrypts username and password before transmitting - Little client support yet

#### • SSL

- Requires public-key crypto in X.509 certificates
- No global PKI → no wide support for client certificates
- Involves heavyweight operations

## Conclusions

- No single authentication scheme can effectively and efficiently meet the requirements of all Web sites and Web clients
- There are clear guidelines (but no standards yet) for designing secure authentication schemes

# **Open Issues**

- What can end users do to protect themselves? - Those who can provide a solution (i.e., vendors) have no incentive to do so.
  - Those who really care about finding a solution (i.e., clients) cannot create one.
- Should there be a standard for authentication protocols? What factors play against establishing such a standard?
- Would you trust a centralized authentication service (such as Microsoft Passport) with your data? A step in which direction is this - forward or backward?

# **SPAM**

#### Problem

- Zero marginal cost of sending an email
- **Solutions**
- Machine learning client to detect spam
- Brightmail

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- Dummy accounts
  Correlate SPAM messages
- · Supply fingerprint to enterprise customers
- Client refuses messages from unknown senders, until

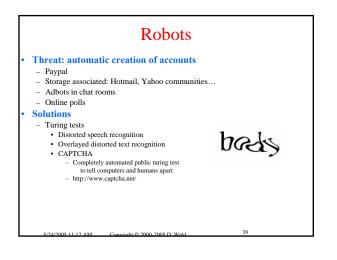
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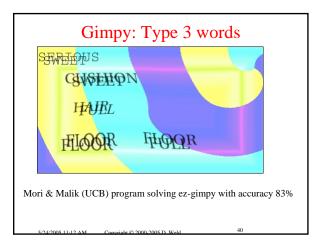
- They respond to a Turing test query
   They execute a computationally expensive applet
   Micropayment

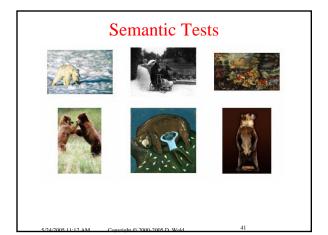


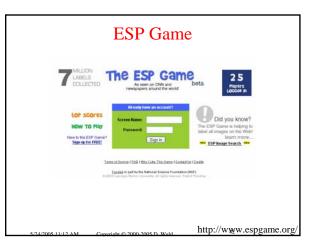
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## Viruses

#### **Defn**

- Requires human action to spread
- Infects most files on local computer
- Doesn't automatically spread across network
- Carries payload (destructive or annoying messages)
- Common MO
- Macro attached to office document

## **Solutions**

- Fingerprint based (to detect viruses)
- Application checksums (to detect tampering)

## Worms

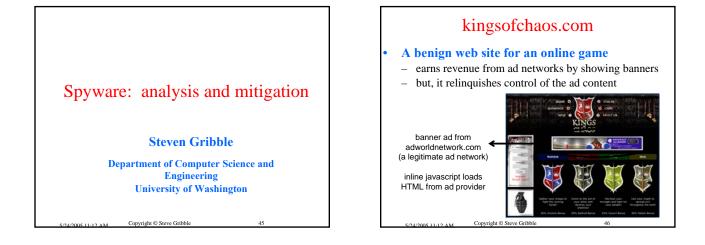
- Automatically spreads to other systems

#### • MO

Defn

- Protocol worms
- Hybrid virus / worms

#### **Solutions**



# Incident: January 2004

#### kingsofchaos.com was given this "ad content"

<script type="text/javascript">document.write(' \u003c\u0062\u006f\u0064\u0079\u0020\u006f\u006e\u0055\u 006f\u0077\u0050\u006f\u0070\u0075\u0070\u0028\u0029\u0 03b\u0073\u0068\u006f\u0077\u0048\u0069 ...etc.

- This "ad" initiated a cascade of redirections through many sites, and ultimately:
  - bombarded the user with pop-up ads

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- hijacked the user's homepage
- exploited an IE vulnerability to install spyware

# What's going on?

#### • The advertiser was really an ad-spammer

- his goal: force users to see ads from his servers
  revenue from ad "affiliate programs"
  - paid to show ads for bogus anti-spyware software

#### • Why install spyware?

- to show ads whether or not the victim is on the Web
- to make escape hard
  - his spyware shows his ads
  - the hijacked home page shows his ads
  - some of his ads re-install spyware and re-hijack

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# Take-away lessons

#### • Your PC has value to third parties

- spyware tries to steal this value from you
  - adware: eyeballs and demographic information
    spyware: sensitive data, PC resources
- Web content should never be trusted
  - even if its direct provider is
- Consumer software and OSs are weak
  - browsers are bug-ridden

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- OSs cannot deal with malicious software

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# What is Spyware?

#### Incredibly difficult to define precisely

- no clean line between good and bad behavior
- hard to define 'informed consent'

#### **Spyware is a** software parasite that:

- collects info of value and relays it to a third party
- hijacks functions or resources of PC

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- installs without consent of user, resists de-installation
- Spyware provides value to others, but not to you

| Types of spyware                             |                        |  |
|----------------------------------------------|------------------------|--|
| Class                                        | # signatures           |  |
| Cookies and web bugs                         | 47                     |  |
| Browser hijackers                            | 272                    |  |
| Adware                                       | 210                    |  |
| Keyloggers                                   | 75                     |  |
| Dialers                                      | 201                    |  |
| Backdoors / monitors                         | 279                    |  |
| From the "Spybot S&                          | D" database, Feb. 2005 |  |
| 5/24/2005 11:12 AM Copyright © Steve Gribble | 51                     |  |

# Infection methods Piggybacking on legitimate software provides revenue stream for free software vendors Drive-by downloads malicious Web content exploits browser vulnerability software is installed and run silently Installed during remote attack some worms now carry spyware payload Snowball effect from existing spyware trojan downloaders

# Spyware trends

#### Most Internet PCs have it

- June '03: 80% of Internet-connected PCs are infected
- [AOL/NCSA online safety study]

#### It's getting more vicious

- December '04: 14% of enterprise PCs have backdoor or monitor spyware
- doubled between October '04 and December '04
- [Webroot reported scan statistics]

## **Convergence of threats**

- worms, viruses, spyware, botnets are fusing

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Two research studies

- November 2003 study of adware within UW
- passive network measurement of entire campus
- measured spread of four adware programs
- Sneak preview of crawler-based study

   active retrieval of content from the Web

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- how much is "out there," and who is spreading it

# UW adware study

#### Examined four programs

- Gator, Cydoor, SaveNow, and eZula
- piggyback installation, adware and HTML rewriting

#### Derived network signatures

- look for the spyware "phoning home"
   e.g., Gator traffic contains Gator/x.xx UserAgent and is sent to a \*.gator.com host
- signatures permit passive network monitoring

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# Method

#### Network monitor deployed at UW

- sniffs packets sent between UW hosts and Internet
- gathered a 7 day Web traceAug. 26th Sept. 2nd, 2003
- looked for packets that match signatures
   traffic matches signature ⇒ sender has spyware

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The major result WWW Gator Cydoor SaveNow eZula 31,303 1,077 399 406 63 # clients (100%) (3.4%) (1.3%) (1.3%) (0.2%) (% clients) 5.1% of UW hosts have  $\geq 1$  of these programs This may appear small, but:

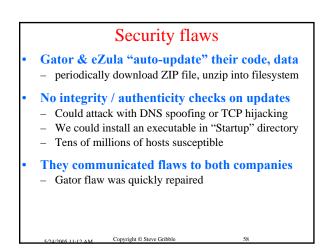
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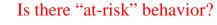
- Only considers 4 spyware programs out of thousands
- University may be non-representativemodem pool has 2.5x higher infection rate

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a Gator vulnerability => 1000+ UW hosts at risk





#### # of web objects downloaded

- fewer than 1000 requests per week: 1.8% have Gator
- more than 12,000 requests per week: 8.9% have Gator

#### **# of executables downloaded**

- none downloaded over week:
- one or more over week:
  - week: 8.4% have Gator

0.9% have Gator

#### using the Kazaa P2P client

 issued one or more Kazaa request: 38% have Gator but...62% of spyware infections are in hosts that didn't issue a Kazaa request

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Two research studies

#### November 2003 study of adware within UW

- passive network measurement of entire campus
- measured spread of four adware programs
- Sneak preview of crawler-based study
   active retrieval of content from the Web

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- how much is "out there," and who is spreading it

# Two studies

- November 2003 study of adware within UW – passive network measurement of entire campus
  - passive network measurement of entire early
     measured spread of four adware programs

#### • Sneak preview of crawler-based study

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- active retrieval and analysis of Web content
- how much is "out there," and where is it coming from?

# Method

- Crawl subsets of Internet to find spyware
- used "heritrix" public domain crawler
- downloaded .zip, .exe, .cab, etc. (programs)
- Cluster of virtual machines to analyze programs
  - "forked" a clean Windows VM per program
    installed program, ran anti-spyware tool to analyze
  - O(1 min) per program

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- on 10-node cluster, O(15,000) programs per day
- many performance optimizations possible

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Major result

Web sites crawled: 12,000
URLs retrieved: 23,714,927

# of executable files downloaded: 9,330

# infected with spyware: 766 (8.21%)
unique spyware programs found: 137

1 in 12 executables on the Internet have spyware!

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# What kind of spyware is out there?

| Behavior                                                  | % spyware |
|-----------------------------------------------------------|-----------|
| Adware                                                    | 57%       |
| Browser hijackers                                         | 56%       |
| Keyloggers                                                | 0.06%     |
| Dialers                                                   | 0.1%      |
| Backdoors / monitors                                      | 15%       |
| • Other stats:                                            |           |
| <ul> <li>58% try to evade discovery or removal</li> </ul> |           |
| <ul> <li>32% monitor Web browsing behavior</li> </ul>     |           |

- most popular: eZula, 180 solutions, SaveNow

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