You are to work on the following questions *alone*. Typeset your answers and submit as a `PDF`.

1. (10 points) **Paxos Acceptor States**

   Consider a deployment of single-instance Paxos with three acceptors. Decide whether each of the following is a valid state of the three acceptors. If the state is not valid, explain why in one sentence. (Hint: A state is valid if there is some sequence of message deliveries and message drops and node failures that leads to the state, assuming a correct implementation of proposers and acceptors.)

   For each part, we give you the highest accepted proposal at all three acceptors (A_1, A_2, A_3) at a single instance in time. Each acceptors highest accepted proposal is either in the form (n,v) where n is the proposal number (à la Paxos Made Simple) and v is a value or $\perp$ which indicates that the acceptor has not accepted any proposal.

   (a) A_1: $\perp$, A_2: $\perp$, A_3: $\perp$
   (b) A_1: $(1,x)$, A_2: $(2,y)$, A_3: $\perp$
   (c) A_1: $(2,x)$, A_2: $(2,y)$, A_3: $\perp$
   (d) A_1: $(1,x)$, A_2: $(2,y)$, A_3: $(3,z)$
   (e) A_1: $(1,x)$, A_2: $(2,x)$, A_3: $(3,x)$

2. (10 points) **Acceptor States in a Larger System**

   Consider a deployment with five acceptors. Is the following state valid? If it is valid, describe an execution that results in this state. If it is not valid, explain why.

   A_1: $(20,x)$, A_2: $\perp$, A_3: $(22,y)$, A_4: $(20,x)$, A_5: $(18,x)$

3. (10 points) **A Dubious Execution**

   Consider another Paxos deployment with acceptors A_1, A_2, and A_3, proposers P_1, P_2, and a distinguished learner L. According to the Paxos paper, a value is chosen when a majority of acceptors accept a proposal with that value, and only a single value is chosen. How does Paxos ensure that the following sequence of events cannot happen? What actually happens, and which value is ultimately chosen?

   – P_1 prepares proposal number 1, and gets responses from A_1, A_2, and A_3.
   – P_1 sends $(1,x)$ to A_1 and A_3 and gets responses from both. However, P_1's proposal to A_2 was dropped. Because a majority accepted, P_1 informs L that x has been chosen. P_1 then crashes.
   – P_2 prepares proposal number 2, and gets responses from A_2 and A_3.
   – P_2 sends $(2,y)$ messages to A_2 and A_3 gets responses from both, so P_2 informs L that y has been chosen.

4. (10 points) **Paxos Liveness**

   In the absence of a distinguished proposer, it is possible for Paxos to fail to make progress even if no messages are dropped and no nodes fail. Briefly describe how this can happen in a system with two proposers and three acceptors. Be specific about which messages are sent and in what order they are delivered.

5. (10 points) **Alternate Paxos Implementation**

   The *Paxos Made Simple* paper has the following definition in page 3.

   > A value is chosen when a single proposal with that value has been accepted by a majority of the acceptors.

   Consider pursuing an alternate implementation based on the following definition.

   > A value is chosen when proposals with that value have been accepted by a majority of the acceptors.

   Would the resulting implementation be correct? Justify your answer in a few sentences either with an informal proof or a scenario where this implementation would violate safety.