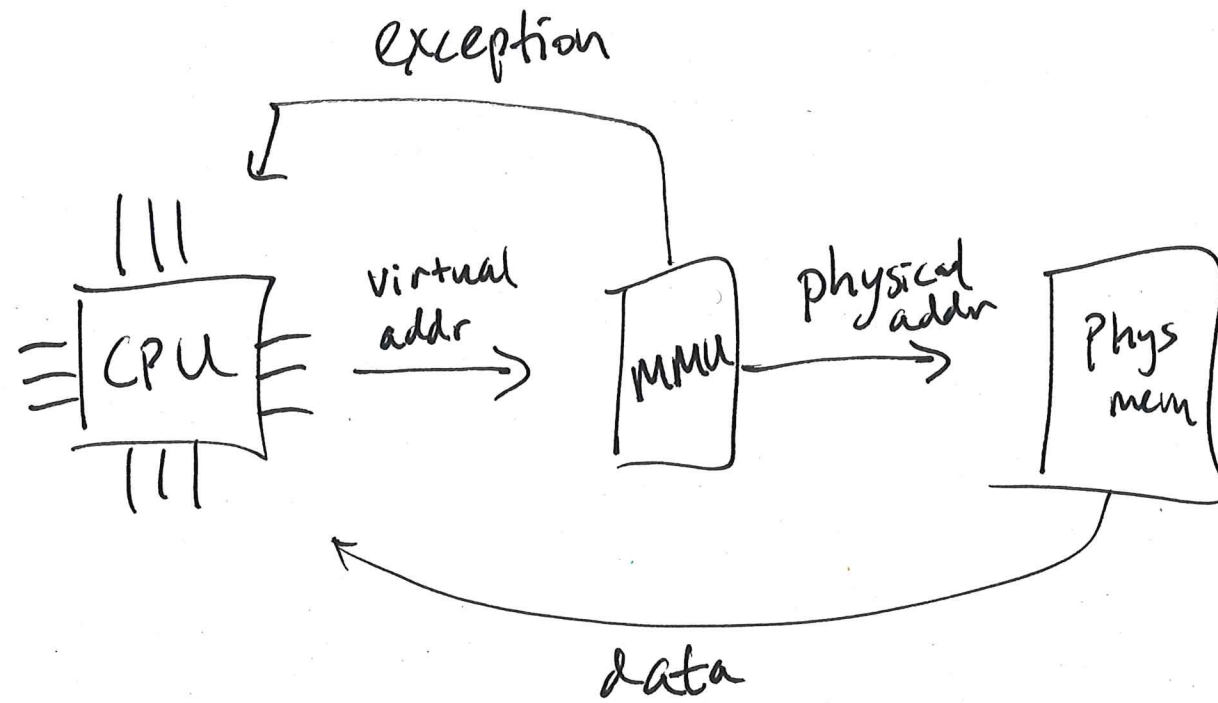


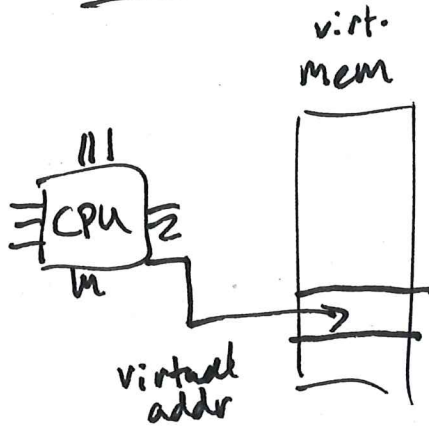
Lec 13

Address Translation



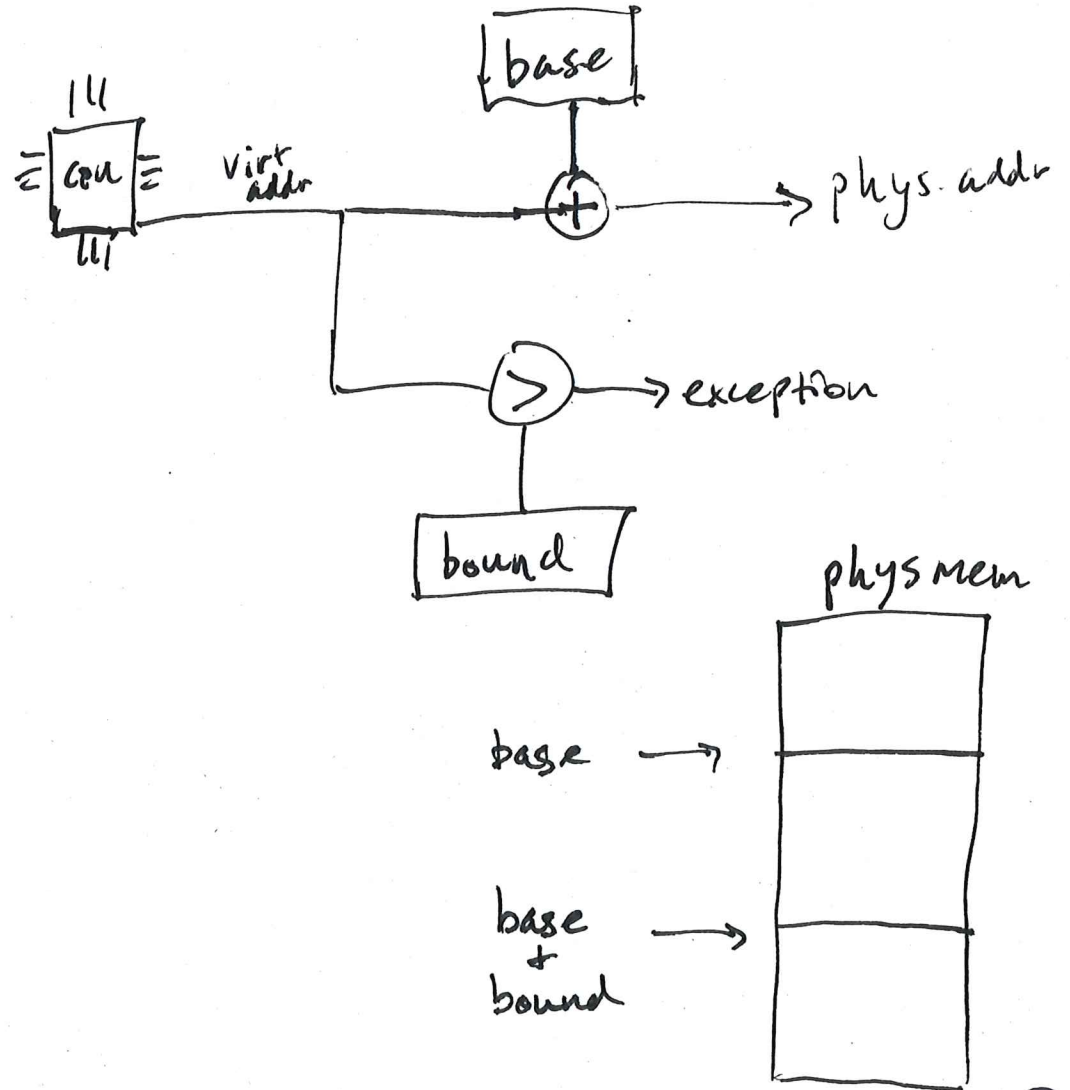
Base + Bounds

Virtual view



- + simple, fast, secure, easy impl in HW
- no protection within process
- no sharing
- difficult to extend region
- fragmentation

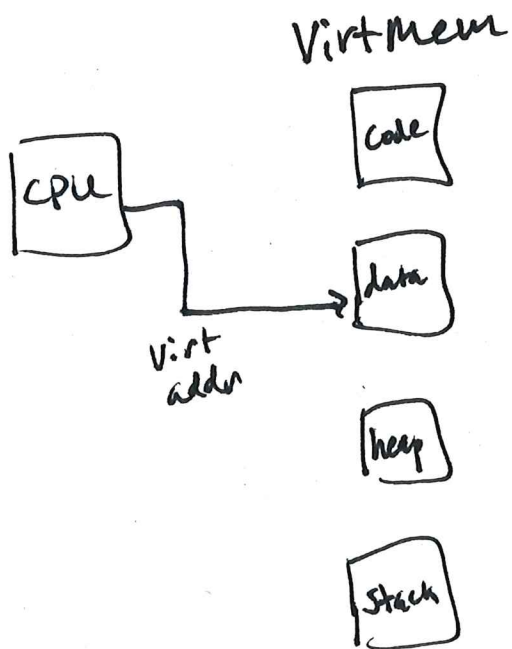
implementation



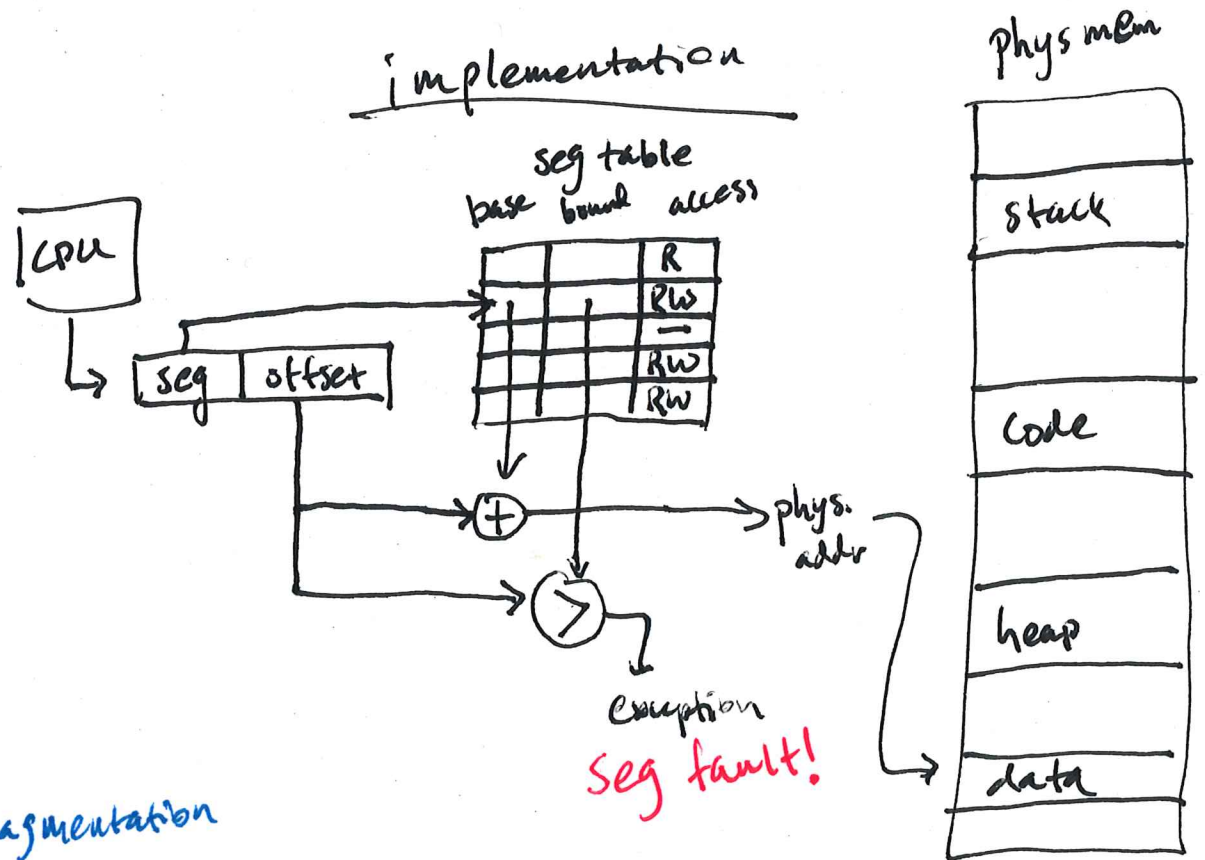
Segmentation

idea: multiple base + bound regions per process
"segments"

virtual view

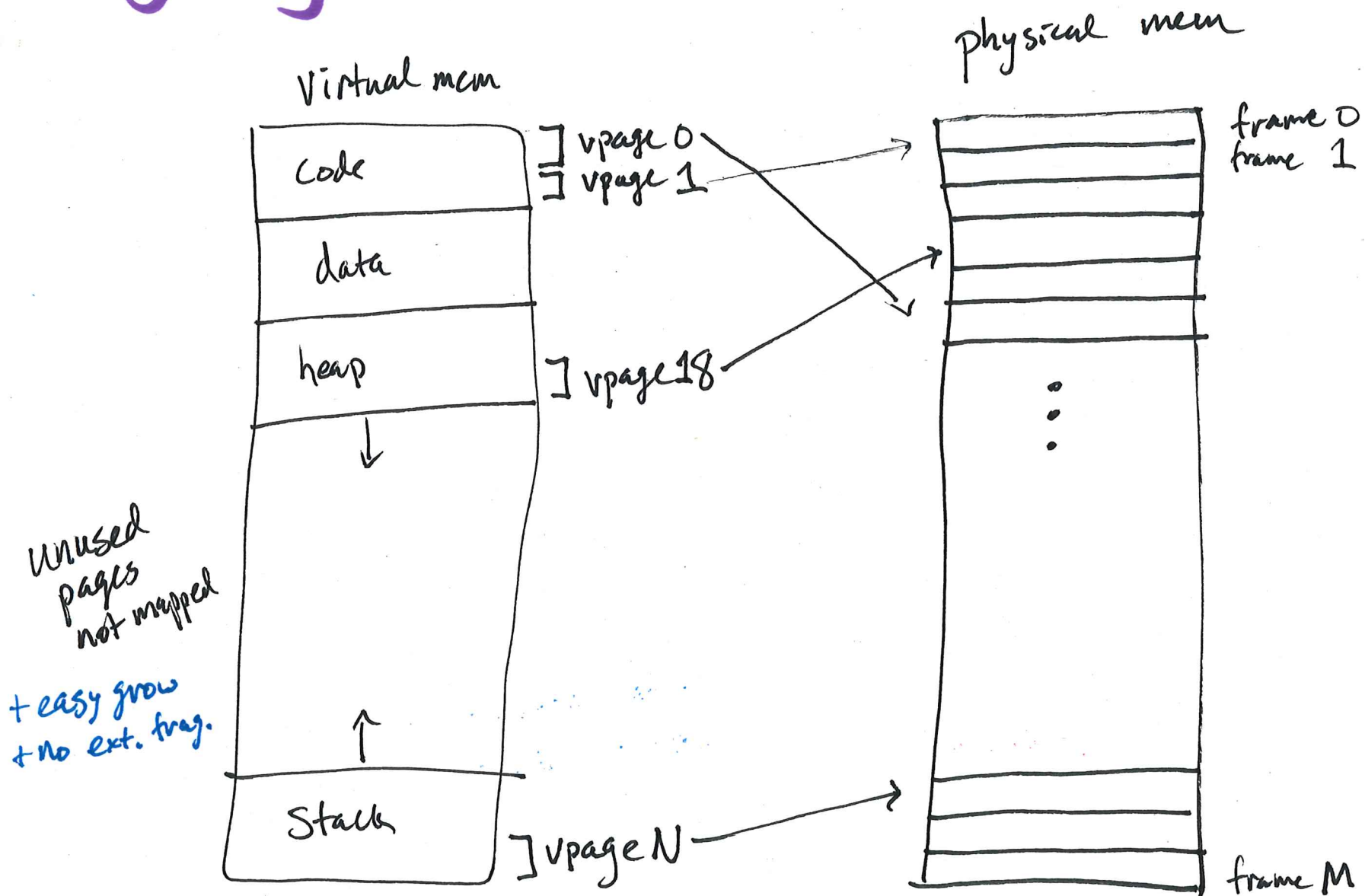


implementation

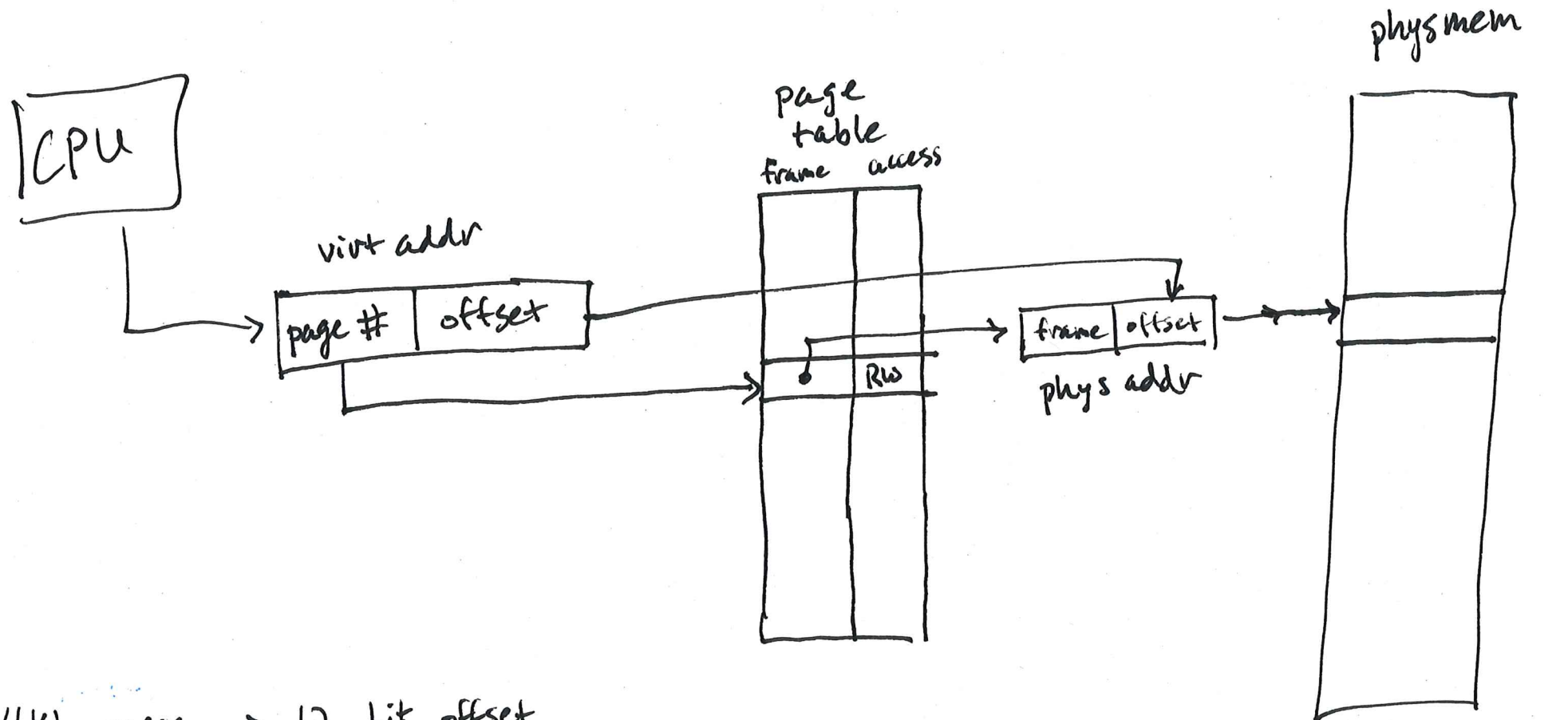


- fragmentation
- hard to grow

Paging (conceptual view)



Paging (implementation)



4Kb pages $\rightarrow$ 12 bit offset
- 64 bit word but **57 bit** virtual addresses
 $\rightarrow$ 45 bit page #

$\hookrightarrow$ how big is page table?

48 bit phys addr
on modern x86

Paging tradeoffs

- What happens when we switch processes?
- What if page size is very small?
- What if page size is very large?
- How big is page table?

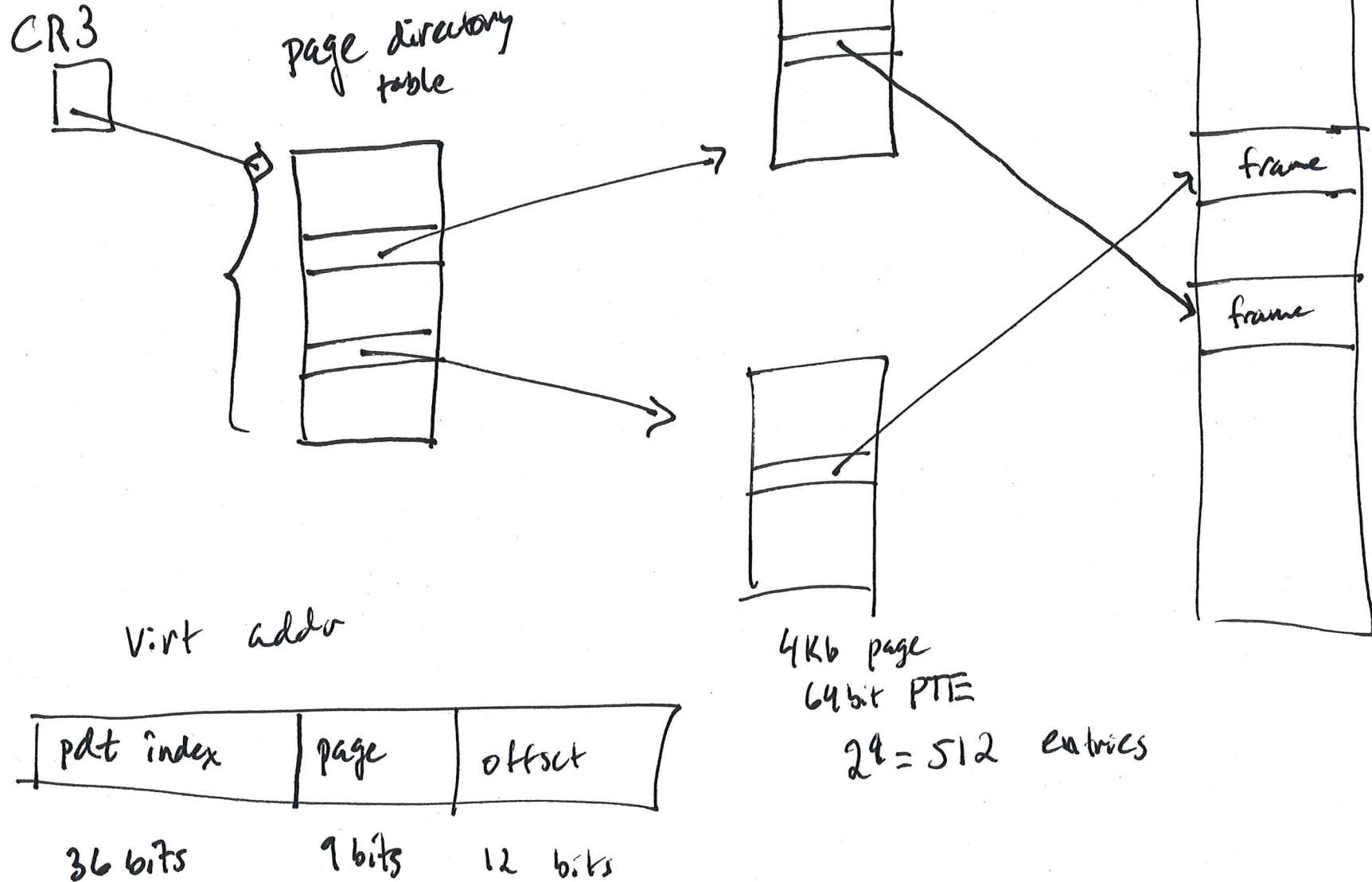
Table 5-19. Format of a Page-Directory Entry that References a Page Table

Bit Position(s)	Contents
0 (P)	Present; must be 1 to reference a page table
1 (R/W)	Read/write; if 0, writes may not be allowed to the 2-MByte region controlled by this entry (see Section 5.6)
2 (U/S)	User/supervisor; if 0, user-mode accesses are not allowed to the 2-MByte region controlled by this entry (see Section 5.6)
3 (P/W/T)	Page-level write-through; indirectly determines the memory type used to access the page table referenced by this entry (see Section 5.9.2)
4 (PCD)	Page-level cache disable; indirectly determines the memory type used to access the page table referenced by this entry (see Section 5.9.2)
5 (A)	Accessed; indicates whether this entry has been used for linear-address translation (see Section 5.8)
6	Ignored
7 (PS)	Page size; must be 0 (otherwise, this entry maps a 2-MByte page; see Table 5-18)
10:8	Ignored
11 (R)	For ordinary paging, ignored; for HLAT paging, restart (if 1, linear-address translation is restarted with ordinary paging)
(M-1):12	Physical address of 4-KByte aligned page table referenced by this entry
51:M	Reserved (must be 0)
62:52	Ignored
63 (XD)	If IA32_EFER_NXE = 1, execute-disable (if 1, instruction fetches are not allowed from the 2-MByte region controlled by this entry; see Section 5.6); otherwise, reserved (must be 0)

Table 5-20. Format of a Page-Table Entry that Maps a 4-KByte Page

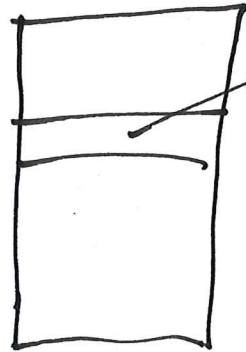
Bit Position(s)	Contents
0 (P)	Present; must be 1 to map a 4-KByte page
1 (R/W)	Read/write; if 0, writes may not be allowed to the 4-KByte page referenced by this entry (see Section 5.6)
2 (U/S)	User/supervisor; if 0, user-mode accesses are not allowed to the 4-KByte page referenced by this entry (see Section 5.6)
3 (PWT)	Page-level write-through; indirectly determines the memory type used to access the 4-KByte page referenced by this entry (see Section 5.9.2)
4 (PCD)	Page-level cache disable; indirectly determines the memory type used to access the 4-KByte page referenced by this entry (see Section 5.9.2)
5 (A)	Accessed; indicates whether software has accessed the 4-KByte page referenced by this entry (see Section 5.8)
6 (D)	Dirty; indicates whether software has written to the 4-KByte page referenced by this entry (see Section 5.8)
7 (PAT)	Indirectly determines the memory type used to access the 4-KByte page referenced by this entry (see Section 5.9.2)
8 (G)	Global; if CR4.PGE = 1, determines whether the translation is global (see Section 5.10); ignored otherwise
10:9	Ignored
11 (R)	For ordinary paging, ignored; for HLAT paging, restart (if 1, linear-address translation is restarted with ordinary paging)
(M-1):12	Physical address of the 4-KByte page referenced by this entry
51:M	Reserved (must be 0)
58:52	Ignored
62:59	Protection key; if CR4.PKE = 1 or CR4.PKS = 1, this may control the page's access rights (see Section 5.6.2); otherwise, it is ignored and not used to control access rights.
63 (XD)	If IA32_EFER.NXE = 1, execute-disable (if 1, instruction fetches are not allowed from the 4-KByte page controlled by this entry; see Section 5.6); otherwise, reserved (must be 0)

Two-level paging



Three level paging

page directory
pointer
table



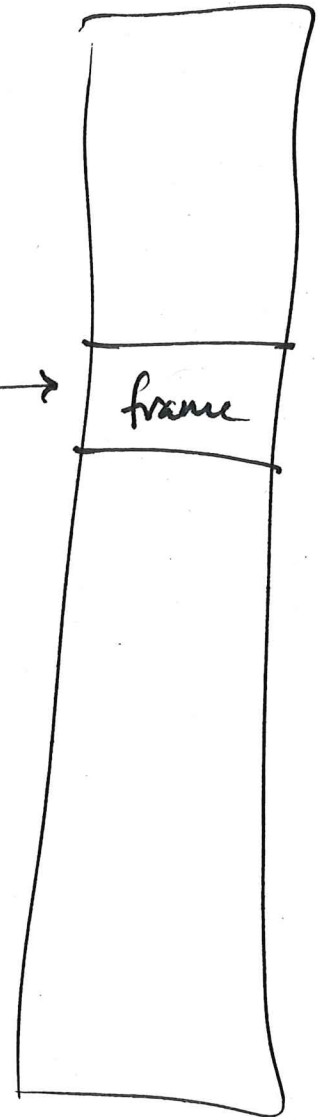
page
directory
table



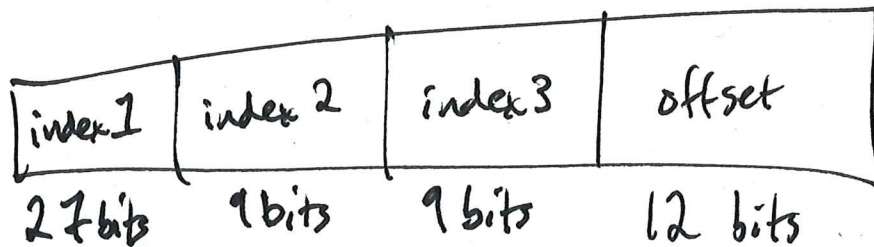
page
table



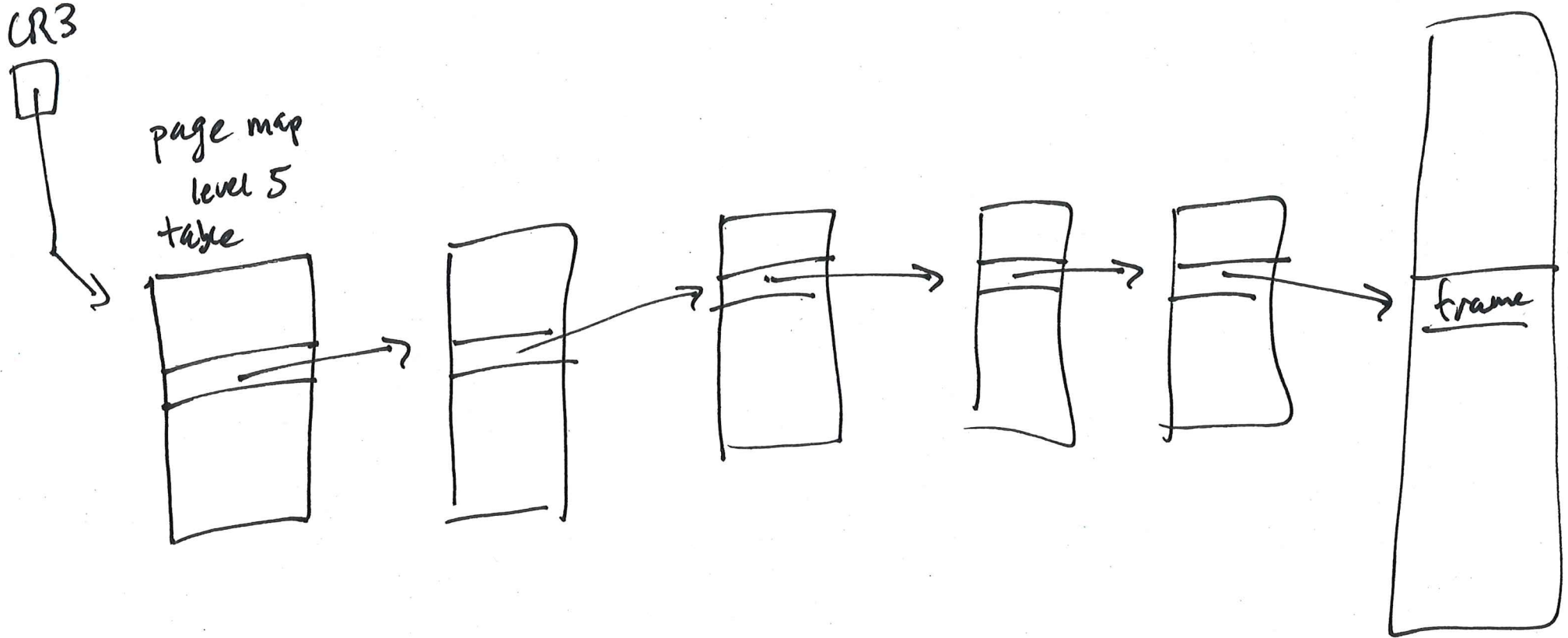
phys mem



virt addr



Five-level paging



Virt addr

i1	i2	i3	i4	i5	offset
9	9	9	9	9	12

$$5 \times 9 + 12 = 45 + 12 = 57 \text{ bit virt addr}$$

13.9