

10/25/23

Deadlocks

→ cycle of waiting threads blocked on each other

t1	t2
lock(A)	lock(B)
lock(B)	lock(A)

mutually recursive locking

t1	t2
buf_a.get()	buf_b.get()
buf_b.put()	buf_a.put()

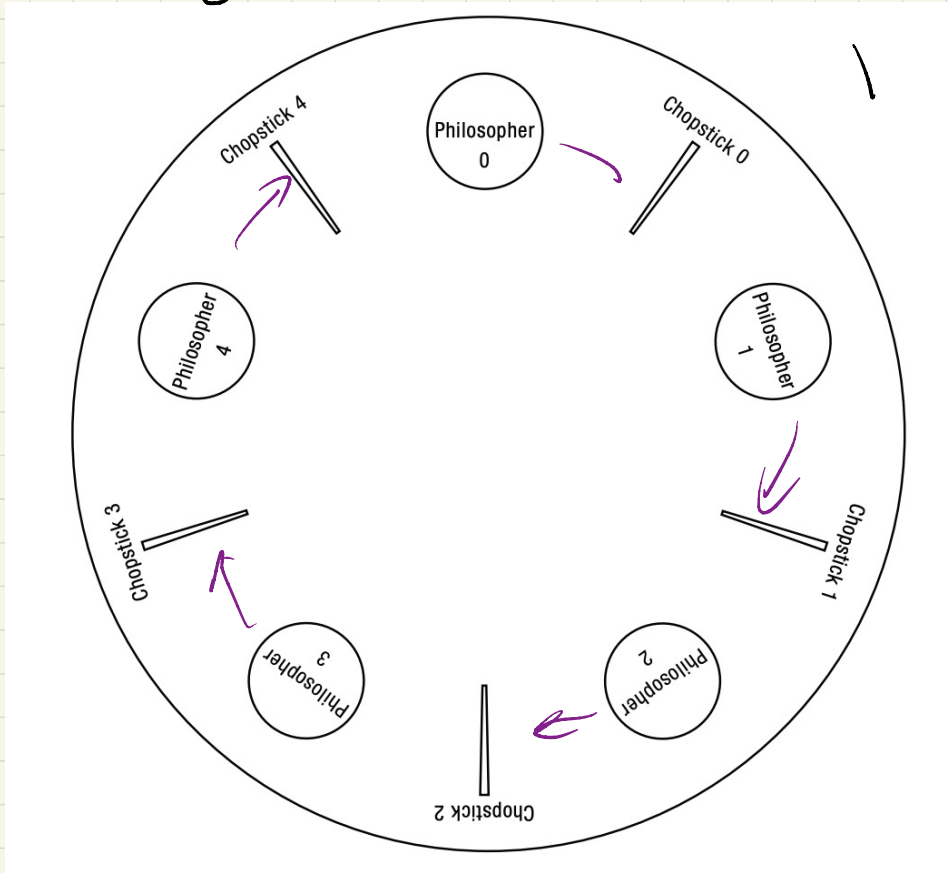
deadlock w/out locks.

resource = producing & consuming power.

→ Necessary Conditions For Deadlock

- ① Bounded Resources: finite instances of each resource
- ② No Preemption: resource can't be forcibly taken away
- ③ Hold & wait: hold on to resource while waiting
- ④ Circular wait: cycle of waiting

Dining Philosophers



5 philosophers
5 chopsticks (resource)
each needs 2 chopsticks to eat
→ grab one from left
& grab one from right

→ Necessary condition for
Readlock but not sufficient
for deadlock.

- single instance \Rightarrow sufficient
- multi instance \Rightarrow not sufficient
(+1 to chopstick)

What to do?

(limit code behavior)

- ① Prevention: when writing code, break one of the 4 conditions for deadlock.
- ② Avoidance: system controls resource access & scheduling to avoid ↑
- ③ Detection: let things be, detect & recover when there's deadlock.

Deadlock Prevention

Bounded resources: provide sufficient resources (release some resources to deal w/ cases before running out of resources)

No Preemption: have preemption

Hold & wait: release while wait

Circular wait: ordering.

Locking ordering

Deadlock Avoidance

- system grants access to resource, can delay fulfilling requests
- needs to know maximum resources each thread needs.

★ have a way to grant all max requests = Safe.

	A	B	C	D	E	max
allobs.	0	0	0	0	0	(2)
A ^{tl}	1 ⁽¹⁾	0	0	0	0	
B ^{tl}	1 ⁽¹⁾	1 ⁽¹⁾	0	0	0	
C ^{tl}	1 ⁽¹⁾	1 ⁽¹⁾	1 ⁽¹⁾	0	0	
D ^{tl}	1 ⁽¹⁾	1 ⁽¹⁾	1 ⁽¹⁾	1 ⁽¹⁾	0	
E ^{tl}	1 ⁽¹⁾	1 ⁽¹⁾	1 ⁽¹⁾	1 ⁽¹⁾	1 ⁽¹⁾	

★ unsafe state.

avail = 5

avail = 4, to be avail = 5

avail = 3

avail = 2

avail = 1

avail = 0

overheadly

to be avail = 2, 3, 4, 5

to be avail = 0 → 1

