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Monitors
(block/uablock)
$\rightarrow$ design pattern for synchronizing threads based on events \& conditions
$\rightarrow$ condition, condition variable, lock
$\rightarrow$ cv-wait, CV-signal, CV-bwadcast ak $\Downarrow$ sleep(chan, IN) $\times$ wakeup(chan)
Basic Pattern


Synchovization Problem: coffee "Spuinons wakeup" int coffee $=0$; wait in a mhle condition
condvar coffee-cu; bop ble the Condtin anymore lock IK; $\uparrow$ might not we wase up.
function get-woffeel) $\{$
lk. acquire ()-;
while Loffee $==0)\{$
coffee-cv.wait (Ik);
$\}$
wffee--;
lk. release();
\}

## Implementation Considerations

When waiting upon a Condition, a "spurious wakeup" is permitted to occur, in general, as a concession to the underlying platform semantics This has little practical impact on most application programs as a Condition should always be waited upon in a loop, testing the state predicate that is being waited for. An implementation is free to remove the possibility of spurious wakeups but it is recommended that applications programmers always assume that they can occur and so always wait in a loop.

Synchowization problem: making breakfast pancalce

Stanter wode:
Lock $1 k$;
Londvar breakfast_cu;
int pancake $=0$;
int benies $=0$;

11 needs I pancalce \& 2 beries for breakfast function plate-breakfast $)\}$

Il produe I pancoke at a tine function make_pancatel) $\}$.
"refill 10 benies at a time function refill-beriesl) $\}$

Lock ik;
Condvar breakfast_cu;
int pancale $=0$;
int beries $=0$;
function plate brealefostl) $\{$
IK. acquirel);
while (pancalce <1 11 beries <2) $\{$ breakfast-cu, wait (lk))
\}
II assert $(\underbrace{\text { pancake }<1 \& d!\underbrace{\text { beries }}_{\text {berries }}<2}_{\text {pancance }>=1}<2$;
pancale --;
bemies - $=2$;
lk.releasel);
\}
make_pancakel) $\{$
(K.acquirel); pancalce $t t$;
breakfast_cu. signall ();
lk. seleasell;
3
Il similar code for beries

Bounded Buffer Problem
$||||||||\mid$ fixed size buffer: $N$ elements
Producer: produce item and pat into an empty slot,
blocks if no room to put item (buffer is full)
Consumer: consume item from a slot, blocks if no item to consume
starter code
char buffer $[N]$;
int consume_ofs = 0; II consumer reads here
int produce-ofs $=0$; 11 producer writes here
int count $=0$, 11 \# of items in the buffer.
function produce( item) function consunre L)

