CSE 451: Operating Systems
Winter 2022

Module 13
Page Table Management, TLBs, and Other Pragmatics

Gary Kimura
Address translation and page faults (refresher!)

What mechanism causes a page fault to occur?

Recall how address translation works
How does OS handle a page fault?

• Interrupt causes system to be entered
• System saves state of running process, then vectors to page fault handler routine
  – find or create (through eviction) a page frame into which to load the needed page (1)
    • if I/O is required, run some other process while it’s going on
  – find the needed page on disk and bring it into the page frame (2)
    • run some other process while the I/O is going on
  – fix up the page table entry
    • mark it as “valid,” set “referenced” and “modified” bits to false, set protection bits appropriately, point to correct page frame
  – put the process on the ready queue
• **(2) Find the needed page on disk and bring it into the page frame**
  – processor makes process ID and faulting virtual address available to page fault handler
  – process ID gets you to the base of the page table
  – VPN portion of VA gets you to the PTE
  – data structure analogous to page table (an array with an entry for each page in the address space) contains disk address of page
  – at this point, it’s just a simple matter of I/O
    • must be positive that the target page frame remains available!
      – or what?
• (1) Find or create (through eviction) a page frame into which to load the needed page
  – run page replacement algorithm
    • free page frame
    • assigned but unmodified (“clean”) page frame
    • assigned and modified (“dirty”) page frame
  – assigned but “clean”
    • find PTE (may be a different process!)
    • mark as invalid (disk address must be available for subsequent reload)
  – assigned and “dirty”
    • find PTE (may be a different process!)
    • mark as invalid
    • write it out
- OS may speculatively maintain lists of clean and dirty frames selected for replacement
  - May also speculatively clean the dirty pages (by writing them to disk)
“Issues”

• Memory reference overhead of address translation
  – 2 references per address lookup (page table, then memory)
  – solution: use a hardware cache to absorb page table lookups
    • translation lookaside buffer (TLB)

• Memory required to hold page tables can be huge
  – need one PTE per page in the virtual address space
  – 32 bit AS with 4KB pages = $2^{20}$ PTEs = 1,048,576 PTEs
  – 4 bytes/PTE = 4MB per page table
    • OS’s typically have separate page tables per process
    • 25 processes = 100MB of page tables
  – 48 bit AS, same assumptions, 64GB per page table!
Solution 1 to (2): Page the page tables

• Simplest notion:
  – Put user page tables in a pageable segment of the system’s address space
    • The OS page table maps the portion of the VAS in which the user process page tables live
  – Pin the system’s page table(s) in physical memory
    • So you can never fault trying to access them
  – When you need a user page table entry
    • It’s in the OS virtual address space, so need the OS page table to translate to a physical address
    • You cannot fault on accessing the OS page table (because it’s pinned)
    • The OS page table might indicate that the user page table isn’t in physical memory
      – That’s just a regular page fault

• This isn’t exactly what’s done any longer
  – Although it is exactly what VAX/VMS did!
  – And it’s a useful model, and a component, for what’s actually done
Solution 2 to (2): Multi-level page tables

• How can we reduce the physical memory requirements of page tables?
  – observation: only need to map the portion of the address space that is actually being used (often a tiny fraction of the total address space)
    • a process may not use its full 32/48/64-bit address space
    • a process may have unused “holes” in its address space
    • a process may not reference some parts of its address space for extended periods
  – all problems in CS can be solved with a level of indirection!
    • two-level (three-level, four-level) page tables
Two-level page tables

• With two-level PT’s, virtual addresses have 3 parts:
  – master page number, secondary page number, offset
  – master PT maps master PN to secondary PT
  – secondary PT maps secondary PN to page frame number
  – offset and PFN yield physical address
Two level page tables

- Master page table
  - Secondary page table
    - Page frame number
- Virtual address
  - Master page #
  - Secondary page#
  - Offset
- Physical address
  - Page frame #
  - Offset
- Physical memory
  - Page frame 0
  - Page frame 1
  - Page frame 2
  - Page frame 3
  - ...
• Example:
  – 32-bit address space, 4KB pages, 4 bytes/PTE
    • how many bits in offset?
      – need 12 bits for 4KB ($2^{12}=4K$), so offset is 12 bits
    • want master PT to fit in one page
      – 4KB/4 bytes = 1024 PTEs
      – thus master page # is 10 bits ($2^{10}=1K$)
      – and there are 1024 secondary page tables
    • and 10 bits are left (32-12-10) for indexing each secondary page table
      – hence, each secondary page table has 1024 PTEs and fits in one page
Generalizing

- Early architectures used 1-level page tables
- VAX, P-II used 2-level page tables
- SPARC used 3-level page tables
- 68030 used 4-level page tables
- Key thing is that the outer level must be wired down (pinned in physical memory) in order to break the recursion – *no smoke and mirrors*
Alternatives

• Hashed page table (great for sparse address spaces)
  – VPN is used as a hash
  – collisions are resolved because the elements in the linked list at the hash index include the VPN as well as the PFN

• Inverted page table (really reduces space!)
  – one entry per page frame
  – includes process id, VPN
  – hard to search! (but IBM PC/RT actually did this!)
Making it all efficient

• Original page table scheme doubled the cost of memory lookups
  – one lookup into page table, a second to fetch the data
• Two-level page tables triple the cost!!
  – two lookups into page table, a third to fetch the data
• How can we make this more efficient?
  – goal: make fetching from a virtual address about as efficient as fetching from a physical address
  – solution: use a hardware cache inside the CPU
    • cache the virtual-to-physical translations in the hardware
    • called a translation lookaside buffer (TLB)
    • TLB is managed by the memory management unit (MMU)
TLBs

- Translation lookaside buffer
  - translates virtual page #s into PTEs (not physical address)

- TLB is implemented in hardware
  - is a fully associative cache (all entries searched in parallel)
  - cache tags are virtual page numbers
  - cache values are PTEs (including protection, valid bit!)
  - with PFN(from PTE) + offset, MMU can directly calculate the physical address

- TLBs exploit locality
  - processes only use a handful of pages at a time
    - can hold the “hot set” or “working set” of a process
  - hit rates in the TLB are therefore really important for performance
Intel i7 Skylake TLB

- TLB has cached levels, too
- Sizes
  - L1:
    - 32Kb for each I/D cache, 4/5 clocks to access
    - 128/64 I/D TLB entries, one clock to access, 9 clocks penalty
  - L2:
    - 256Kb, 12 clocks to access
    - 1536 TLB entries, 14 clocks to access, 17 clocks penalty
  - L3:
    - 8Mb, 42 clocks to access
Managing TLBs

- Address translations are mostly handled by the TLB
  - >99% of translations, but there are **TLB misses** occasionally
  - in case of a miss, translation is placed into the TLB, values are evicted. Selection algorithm is proprietary

- Hardware (memory management unit (MMU))
  - knows where page tables are in memory
    - OS maintains them, HW access them directly
  - tables have to be in HW-defined format
  - this is how x86 works
    - And that was part of the difficulty in virtualizing the x86 …

- Software loaded TLB (OS)
  - TLB miss faults to OS, OS finds right PTE and loads TLB
  - must be fast (but, 20-1000 cycles typically)
    - CPU ISA has instructions for TLB manipulation
    - OS gets to pick the page table format
Managing TLBs (2)

• OS must ensure TLB and page tables are consistent
  – when OS changes protection bits in a PTE, it needs to invalidate the PTE if it is in the TLB

• What happens on a process context switch?
  – remember, each process typically has its own page tables
  – need to invalidate all the entries in TLB! (flush TLB)
    • this is a big part of why process context switches are costly
  – can you think of a hardware fix to this?

• When the TLB misses, and a new PTE is loaded, a cached PTE must be evicted
  – choosing a victim PTE is called the “TLB replacement policy”
Functionality enhanced by page tables

• Code (instructions) is read-only
  – A bad pointer can’t change the program code
• Dereferencing a null pointer is an error caught by hardware
  – Don’t use the first page of the virtual address space – mark it as invalid – so references to address 0 cause an interrupt
• Inter-process memory protection
  – My address XYZ is different that your address XYZ
• Shared libraries
  – All running C programs use libc
  – Have only one (partial) copy in physical memory, not one per process
  – All page table entries mapping libc point to the same set of physical frames
    • DLL’s in Windows
More functionality

• Generalizing the use of “shared memory”
  – Regions of two separate processes’s address spaces map to the same physical frames
  – Why? Faster inter-process communication
    • Just read/write from/to shared memory
    • Don’t have to make a syscall
  – Will have separate PTE’s per process, so can give different processes different access rights
    • E.g., one reader, one writer

• Copy-on-write (CoW), e.g., on fork()
  – Instead of copying all pages, create shared mappings of parent pages in child address space
    • Make shared mappings read-only for both processes
    • When either process writes, fault occurs, OS “splits” the page
Less familiar uses

- **Memory-mapped files**
  - instead of using open, read, write, close
    - “map” a file into a region of the virtual address space
      - e.g., into region with base ‘X’
    - accessing virtual address ‘X+N’ refers to offset ‘N’ in file
    - initially, all pages in mapped region marked as invalid
  - Using that “table that looks like a page table”…
    - OS reads a page from file whenever invalid page accessed
    - OS writes a page to file when evicted from physical memory
      - only necessary if page is dirty
Memory mapped files

- Imagine you have a pointer-based, in-memory data structure, like a tree
- You want to preserve it across runs
- Usual approach:
  - Serialize on way from memory to a disk file, deserialize on way from file back to memory
    - E.g., to serialize, perform a depth-first traversal, writing each node to disk as you go; to deserialize, do the opposite
- Potentially easier
  - Allocate tree nodes in a “region”
  - Treat the memory region as a file, using the memory-mapped file facility
  - Normal paging causes changes to be pushed to disk; the file is still there next time you run
  - What happens if you crash? Uh oh…
More unusual uses

• We saw that page replacement algorithms use the fact that “soft faults” are relatively cheap
  – Soft faults: faults on pages that are actually in memory, but whose PTE entries have artificially been marked as invalid
• That idea can be used whenever it would be useful to trap on a reference to some data item
• Example: debugger watchpoints
  – How?
• (The utility of this idea is limited by the fact that the granularity of detection is the page)
Summary

• We know how address translation works in the “vanilla” case (single-level page table, no fault, no TLB)
  – hardware splits the virtual address into the virtual page number and the offset; uses the VPN to index the page table; concatenates the offset to the page frame number (which is in the PTE) to obtain the physical address

• We know how the OS handles a page fault
  – find or create (through eviction) a page frame into which to load the needed page
  – find the needed page on disk and bring it into the page frame
  – fix up the page table entry
  – put the process on the ready queue
• We’re aware of two “gotchas” that complicate things in practice
  – the memory reference overhead of address translation
    • the need to reference the page table doubles the memory traffic
    • solution: use a hardware cache (TLB = translation lookaside buffer) to absorb page table lookups
  – the memory required to hold page tables can be huge
    • solution: use multi-level page tables; can page the lower levels, or at least omit them if the address space is sparse
      – this makes the TLB even more important, because without it, a single user-level memory reference can cause two or three or four page table memory references … and we can’t even afford one!
• TLB details
  – Implemented in hardware
    • fully associative cache (all entries searched in parallel)
    • cache tags are virtual page numbers
    • cache values are page table entries (page frame numbers)
    • with PTE + offset, MMU can directly calculate the physical address
  – Can be small because of locality
    • 16-48 entries can yield a 99% hit ratio
  – Searched before the hw or OS walks the page table(s)
    • hit: address translation does not require an extra memory reference (or two or three or four) – “free”
    • miss: walk the page table(s) to translate the address; this translation is put into the TLB, evicting some other translation; typically managed LRU
– On context switch
  • TLB must be purged/flushed (using a special hardware instruction) unless entries are tagged with a process ID
    – otherwise, the new process will use the old process’s TLB entries and reference its page frames!

• Cool tricks
  – Read-only code
  – Dereferencing a null pointer is an error
  – Inter-process memory protection
  – Shared libraries
  – Inter-process communication
  – Shared memory
  – Copy-on-write
  – Memory-mapped files
  – Soft faults (e.g., debugger watchpoints)