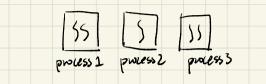
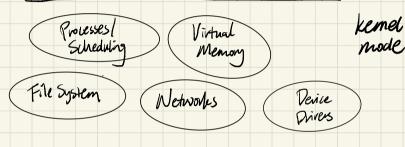
11/30 OS Structures





user mode

-> 50 far all the OS components ne've learned rens in Kernel mode

40 OS = Kernel

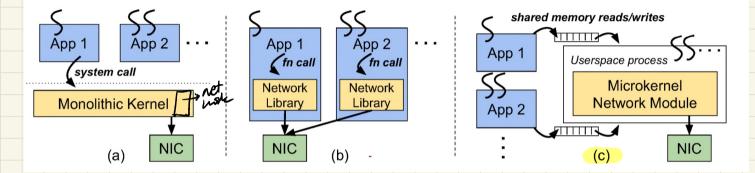
monolithic kennel

all components drave the same address space & same faiture domain How can ne make kemel more reliable? D Modulanze it 50 that when one component fails the entire kemel doesn't consh O Use programming languages with Stronger promises. -> most guarantees memory Safety (no dota races)

How de ne isolate fault for each component ? Process is a unit of fault domain, so maybe more lach component into their own user process !

TRC TRC & Uhat are some challenges Filesys process 1 process 2 process 3 if he more an OS component user Mode into a user praess? Filesys Process 1). Availability kernel mode Processes/ Schedwing Wirtual Meanory -> What if the fs process terminates? no one com use filesys anymore? FS Mechanism Networks Device Drivers 2). Commication -> How do ne request service from the fs process? IPC? Can do this for other OS components as well? 3). Access Perisson is lots of the industry in thigh overhead -> this approach is called microbernel -> Access to certain HW still requires us to be in the kernel male ! Keeps the kemel Small & implement most is needs both a Kernel part for priviledged access OS components naser Le a user provers paint for policies de designs Space !

Example of Minokernel in use = Grogle's Snap



-> mased off of linux, moves network stack into a userspace process!

https://research.google/pubs/pub48630/