Operating Systems

Section 1 - Course Introduction, Misc. Topics
7 Jan 2021
Day 301 of Online UW
Who is this person talking at you all?

TAS - TELL US WHO YOU DO BE
Who is this person talking at you all (AA)?

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Hometown: Orange, CA

Education: 5th Year Master’s

Interests:
- Rockets!
- Other than that
  - The occasional video game
  - The occasional book
  - The occasional hike

Me, next to a Mars rover (not quite a rocket)
Schedule

1) Brief Lab 1 Intro
2) Brief recap of 351/333 topics
3) Office hours, discussion board
4) Misc. info
Lab 1 is out!

- Partner form on Ed
  - Only one partner needs to fill it out
- Due next Friday (1/15/21)
  - Start early!
Where to start?

Start by reading:

- lab/overview.md - A description of the xk codebase. A MUST-READ!
- lab/lab1.md - Assignment write-up
- lab/memory.md - An overview of memory management in xk
- lab1design.md - A design doc for the lab 1 code
  - You will be in charge of writing design docs for the future labs. Check out lab/designdoc.md for details.
File Information

Need a way to store the following information about a file:

- In memory reference count
- A reference to the inode of the file
- Current offset
- Access permissions (readable or writable) [for when we add pipes and file writability later]
There will be a global array of all the open files on the system (bounded by NFILE) placed in static memory.
fd = index into local File Descriptor Array

Process View

Global Array

- File Info Struct Index 0
- File Info Struct Index 1
- File Info Struct Index 2
- File Info Struct Index 3
- File Info Struct Index 4
- File Info Struct Index 5
- File Info Struct Index 6

Process 1’s File Descriptor Array

0 1 2 3 NOFILE

Process 2’s File Descriptor Array

0 1 2 3 NOFILE
Functions
**filewrite and fileread**

- Writing or reading of a "file", based on whether the file is an inode or a pipe.
  - Note that file is in quotes. A file descriptor can represent many different things. You could be reading from a file, or you could be reading from standard in or a pipe!
- Don’t need to worry about the pipe part for this lab, just the inode files.
- Check out the functions `readi` and `writei` defined in `kernel/fs.c`
  - file layer provides “policy” for accessing files, inode layer provides “mechanism” for reading/writing
**fileopen**

Finds an open file in the global file table to give to the process
fileclose

Release the file from this process, will have to clean up if this is the last reference
filedup

Duplicates the file descriptor in the process’ file descriptor table
Part 2: The C Programming Language
What was C, again? A brief recap
To jog your memory, not to re-teach C. Skimming over 351/333 isn't a bad idea

- Functions & Structs (they exist, and are about as complex as C gets)
- Forward Declarations & Header files (working with multi-file projects)
- The Preprocessor (and how it relates to header files)
- Pointers & Memory

- Assembly
int sum3(int x, int y, int z) { // A function, like in most programming languages
    return x + y + z;
}

typedef struct s_Point2D { // Not a class: only public fields. No inheritance, no methods
double x;
double y;
} Point2D; // Because of typedef, can use either type “struct s_Point2D”, or type “Point2D”

double dot(struct s_Point2D point1, Point2D point2) { // Pass by value
    return point1.x * point2.x + point1.y * point2.y;
}
Forward Declarations

- C compiler is single pass
  - If you define function A, then function B, the compiler doesn't know about B until it's done reading A
  - This will have a compiler error: when reading get4()'s implementation, get3() is unknown

```c
int get4() { return get3() + 1; }
```

```c
int get3() { return 3; }
```
Forward Declarations, Header Files

- The solution? Declare things before defining them

  ```c
  int get3(); // There will exist a function called get3 with no args, returning int
  int get4(); // Also one called get4()
  ```

  ```c
  int get4() { return get3() + 1; } // This is okay: we know that get3() will exist
  ```

  ```c
  int get3() { return 3; }
  ```

- We end up putting our forward declarations in a header file so that we know everything is declared first. As a bonus, other code can reference the header file to use functions it declares
Now we have two problems

1) Implementations don’t have the forward declarations anymore (we moved to a new file)
   Solution: The Preprocessor: `#include “MyHeader.h”`
   - In effect, replace this line with the entire contents of MyHeader.h
   - Now, implementation can reference things the header declares

2) Duplicate declarations: if we include the same header multiple times (in order to references its contents from multiple places), have repeated code -- since `#include` is copy-paste
   Solution: Header Guards (`#ifndef AAA \n #define AAA \n ...... #endif`)
   Everything between the ifndef and endif is only duplicated once
Header Files & The Preprocessor (2)

// mymath.h

#ifndef MYMATH // Header guard: in effect
#define MYMATH // only include declaration

int get4(); // once
int get3();
#endif

#include "mymath.h" // declare get3/4
int get4() { return get3() + 1; }
int get3() { return 3; } // Compiler binds
// implementation to declaration

#include <stdio.h>

#include "mymath.h" // Have get3’s declaration

void print3() { printf("%d", get3()); } // prints 3

// morecode.c
Preprocessor Macros to Know

#include: embed the given file here. As in, copy-paste the whole thing.

#define A (or #define A B): register A as a known symbol. If B is given, replace all occurrences of A with B
  -> Used for constants! (e.g. “#define SIZE 20”)
  -> Also used for macros. e.g. “#define MAX(a,b) a > b ? a : b”
  This is a find/replace operation

#if ___ / #endif: Only include the code between the #if and #endif if the condition is true

ifdef ___ / ifndef ___ / #endif: Only include the code between this and endif if the symbol is/isn’t defined
Pointers & Addresses

Get the address of where something is stored in (virtual) computer memory
- a 32/64 bit (4/8 byte) number.
- Just a number. You can do arbitrary math to a pointer value. Might end up with an invalid address......

Dereferencing: “Give me whatever is stored in memory at this address”.
If you use an address outside of what your program has claim to, segfault.

Array subscripting is pointer math:
```
typename* myPtr = some address
myPtr[3] == *(myPtr + 3 * sizeof(typename))
```
void increment(int* ptr) {
    *ptr = *ptr + 1;
}

int x = 3;
increment(&x);
// x is now 4

← Pass in a pointer: the address at which some int is stored
*ptr gets the value stored at the address stored by ptr
So we assign to the memory at ptr’s address:
   “whatever was there before + 1”
The pointer (address) is passed by value: “ptr = ptr + 1” only changes the local “ptr” variable

← Use the address at which ‘x’ resides in memory
Memory

Program memory:
- The Stack
  - Function data
  - Tied to function lifetime
  - Grows down
- The Heap
  - Anything dynamically allocated
  - Persists until deleted
  - Grows up... ish

System memory:
- Registers
  - The CPU has a few places to store actively-used data
- RAM
  - Volatile data associated with a running program
- Cache
  - CPU has caches to make lookups faster
All CPU operations must be done on data in registers
Compilation Process

1) Preprocessor
   Scan all files top-to-bottom, doing text substitution as appropriate

2) Compiler
   Compile C files into assembly files (intermediate form)

3) Assembler
   “Assemble” assembly code into machine code (creating object files)

4) Linker
   Combine all source files together into executable. Include external libraries, too
Compilation & Assembly

C code is human-readable (mostly). It's not machine-readable -- you can't just feed the CPU some C.

Need to **compile** the code from C into machine code (ones and zeros)

-> It’s just a bunch of simple instructions. Add *this* and *that*. Move this data over there.

Assembly is raw CPU instructions in a slightly-readable format.

This class:
- You won’t need to write assembly
- You might need to read a bit of assembly
- You **will** need to use registers, understand how they work, and understand how x86 does function calls

We'll go over some of the details for relevant assignments, but skimming over CSE351 would not hurt.
Wow, that was a lot of review, very quickly

I assume many of you have questions

Please ask.
Your best friend this quarter

This is a systems class and you'll be doing a LOT of debugging
Also lots of pointers.
Really, the pointers are the main reason for the debugging
GDB commands to know: a non-exhaustive list

`gdb path/to/exe`
run: start execution of the given executable
n: run the next line of code. If it’s a function, execute it entirely.
s: run the next line of code. If it’s a function, step into it
c: run the rest of the program until it hits a breakpoint or exits

b ____: set a breakpoint for the given function or line (e.g. “b myfile.c:foo” or “b otherfile.c:43”)
b(t): get the stack trace to the current point. Can be ran after segfaults!
up/down: go up/down function stack frames in the backtrace
(r)watch ____: set a breakpoint for the given thing being accessed
p ____: print the value of the given thing
x ____: examine the memory at an address. Many flags
General Debugging Tips (more to come later)

- Get familiar with GDB
  - Stepping through line by line and printing out variables is slow, but will find the bug.
- Make sure you know what the code is supposed to do first
  - There are a lot of complicated systems, with limited framework. Unlike 333, this isn't fill-in-the-blank
- Don't be afraid of println debugging
  - It can be an efficient way to find what section of code is wrong so your GDB debugging can be more focused
  - Some timing-based bugs don't really show up in GDB. Need printlns to figure out problematic orderings
- Get familiar with GDB
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- Get familiar with GDB
Part 3: Procedure
Regarding office hours

- Zoom is not a great platform for helping multiple people at the same time

- There are a lot of strange ways you can break xk
- Unlike in other classes, there are many functional ways to structure your code
- Going through GDB in office hours is way too slow

What are we saying?

- If there's a queue, we might not be able to find your bugs for you, just give you a push in a direction
- GDB (single process) and println-debugging (multiple processes) will find the problem. It just takes time.
Discussion Board

If you’ve tried debugging and have come up against a wall that would take too long for office hours, consider posting on the discussion board. Specifics of what to include TBD, but along the lines of

- What is the problem
- In which methods does the problem manifest
- What code are you confident works
- What debugging have you tried

But once again -- for some of the weirdest bugs there’s no solution other than taking a couple hours to run through GDB line-by-line
Resources

We’ll be putting out a signup form for partners. Please try to find people via the discussion board (or other platforms) so you don't get tossed into random matching.

All labs are in the git repo. There'll be one due roughly every two weeks, but you can work ahead if you really want to (Note that TAs will not be assisting with problems for not-yet-assigned labs).

There’s some documentation in the project repository. Read it!
Questions (C, the class, your TA, etc)
Meeting people is hard

No in-person interaction for the next weeks/months

That is it for section, so if anyone wants to say hi/chat, feel free. I’ll leave the room open.