CSE 451: Operating Systems Winter 2017

Module 24 Distributed Systems

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What is a "distributed system"?

- Nearly all systems today are distributed in some way
 - they use email
 - they access files over a network
 - they access printers over a network
 - they're backed up over a network
 - they share other physical or logical resources
 - they cooperate with other people on other machines
 - they access the web
 - they receive video, audio, etc.

Loosely-coupled systems

- Earliest systems used simple explicit network programs
 - FTP (rcp): file transfer program
 - telnet (rlogin/rsh): remote login program
 - mail (SMTP)
- Each system was a completely autonomous independent system, connected to others on the network

- Even today, most distributed systems are loosely-coupled (although not that loosely!):
 - each CPU runs an independent autonomous OS
 - computers don't really trust each other
 - some resources are shared, but most are not
 - the system may look differently from different hosts

Closely-coupled systems

- A distributed system becomes more "closely-coupled" as it
 - appears more uniform in nature
 - runs a "single" operating system
 - has a single security domain
 - shares all logical resources (e.g., files)
 - shares all physical resources (CPUs, memory, disks, printers, etc.)
- In the limit, a distributed system looks to the user as if it were a centralized timesharing system, except that it's constructed out of a distributed collection of hardware and software components

Tightly-coupled systems

- A "tightly-coupled" system usually refers to a multiprocessor
 - runs a single copy of the OS with a single workload queue
 - has a single address space
 - usually has a single bus or backplane to which all processors and memories are connected
 - has very low communication latency
 - processors communicate through shared memory

Some issues in distributed systems

- Availability (can you get it now?)
- Reliability (may you lose stuff permanently?)
- Performance (on a LAN, on a WAN, global)
- Scalability (can it grow modularly as you add users, without relying on ever-bigger/faster computers?)
- Transparency (how visible is the distribution to users?)
- Programming models (how visible is the distribution to programmers?)
- Communication models (messages, RPC, etc.)
- Security

Example: Grapevine distributed mail service

- Xerox PARC, 1980
 - cf. Microsoft Outlook/Exchange today!!!!!

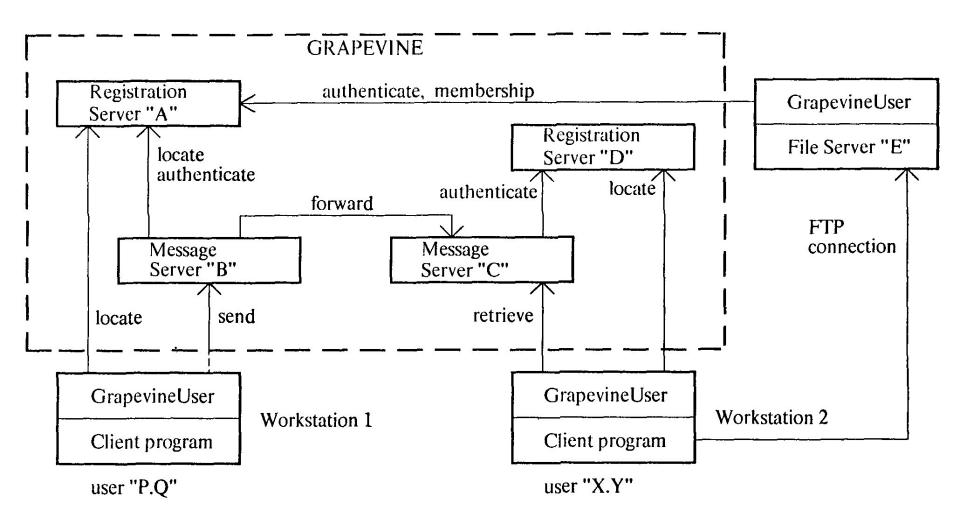
Goals

- cannot rely on integrity of client
- once the system accepts mail, it will be delivered
- no single Grapevine computer failure will make the system unavailable to any client either for sending or for receiving mail

Components

- GrapevineUser package on each client workstation
- Registration Servers
- Message Servers
- Implementation: Remote Procedure Call

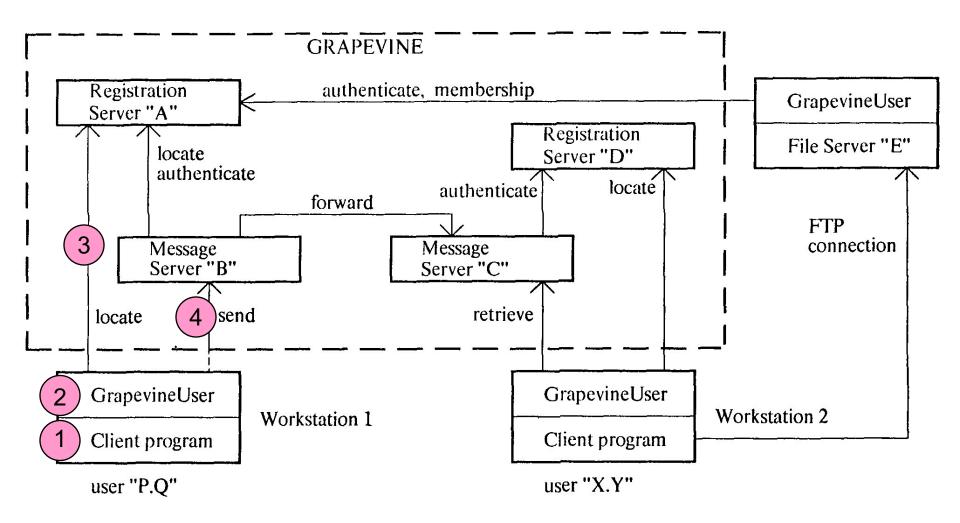
Grapevine: Functional diagram



Grapevine: Sending a message

- User prepares message using mail client
- Mail client contacts GrapevineUser package on same workstation to actually send message
- GrapevineUser package
 - Contacts any Registration Server to get a list of Message Servers
 - Contacts any Message Server to transmit message
 - presents source and destination userids, and source password, for authentication
 - Message Server uses any Registration Server to authenticate
 - sends message body to Message Server
 - Message Server places it in stable storage and acknowledges receipt

Grapevine: Functional diagram



Registries

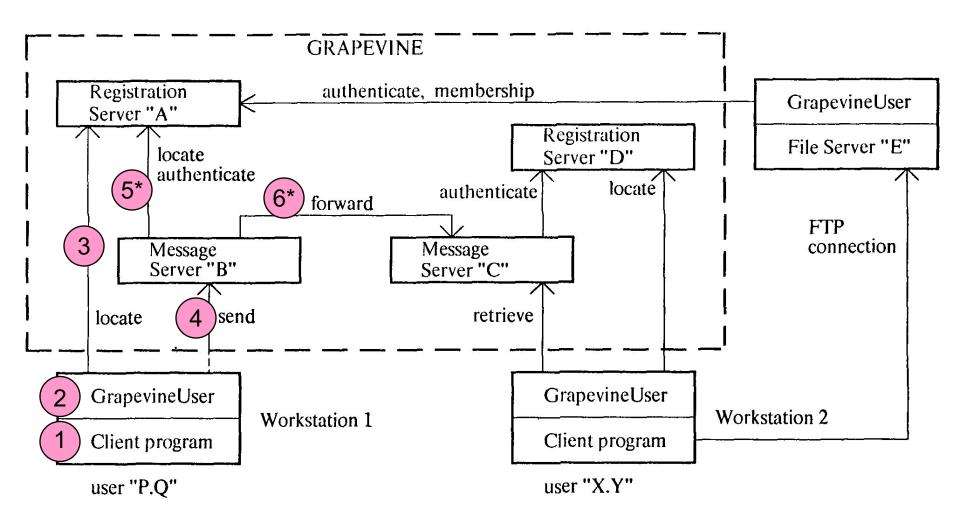
- Actually, I lied: There's an additional step.
 - For scalability, users are partitioned into registries "user 'P.Q'" is user P in registry Q.
 - Registries are replicated.
 - There is one registry that is replicated on every registration server: the registry of registries.
 - So, when I said:

Message Server uses any Registration Server to authenticate what actually happens is the Message Server contacts any Registration Server to obtain a list of those Registration Servers holding the registry of the user, then contacts one of those registration servers to authenticate the user

Grapevine: Transport and buffering

- For each recipient of the message, Message Server contacts any Registration Server to obtain list of Message Servers holding mail for that recipient
 - Same lie as before
- Sends a copy of the message to one of those Message Servers for that recipient

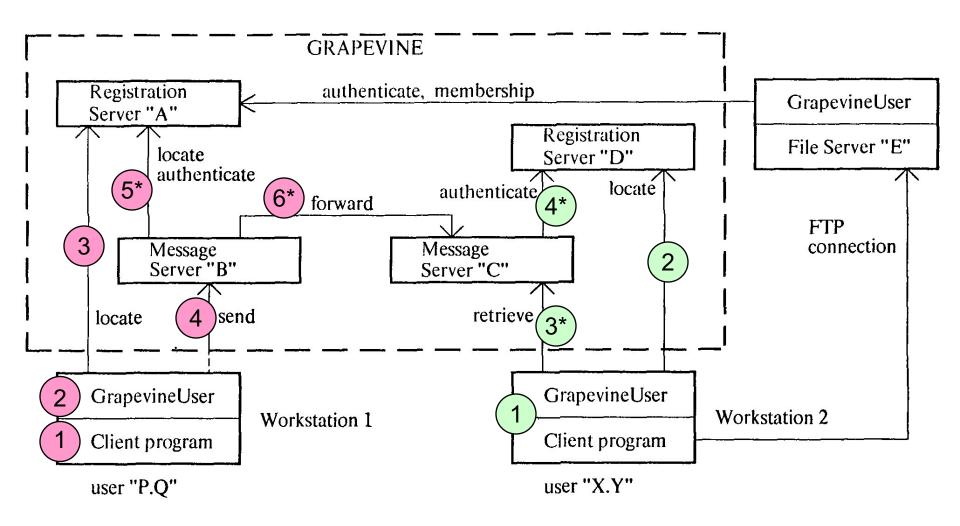
Grapevine: Functional diagram



Grapevine: Retrieving mail

- User uses mail client to contact GrapevineUser package on same workstation to retrieve mail
- GrapevineUser package
 - Contacts any Registration Server to get a list of each Message Server holding mail for the user ("inbox site")
 - Same lie as before
 - Contacts each of these Message Servers to retrieve mail
 - presents user credentials
 - Message Server uses any Registration Server to authenticate
 - acknowledges receipt of messages so that the server can delete them from its storage

Grapevine: Functional diagram



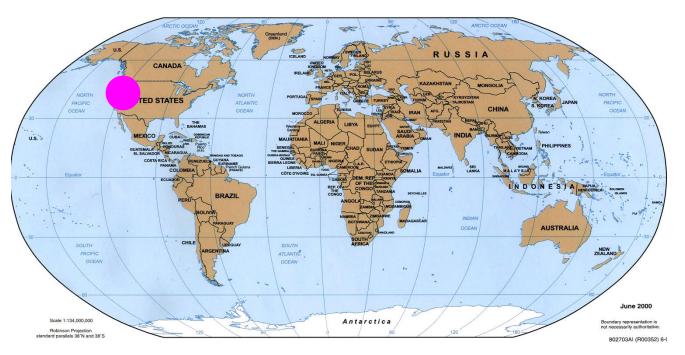
Grapevine: Scalability

- Can add more Registration Servers
- Can add more Message Servers
- Only thing that didn't scale was handling of distribution lists
 - the accepting Message Server was responsible for expanding the list (recursively if necessary) and delivering to an appropriate Message Server for each recipient
 - some distribution lists contained essentially the entire user community
- Jeff Dean (Google) told us they don't even think about more than two decimal orders of magnitude
 - fundamental design decisions will need to change
 - advances in technology will make it possible

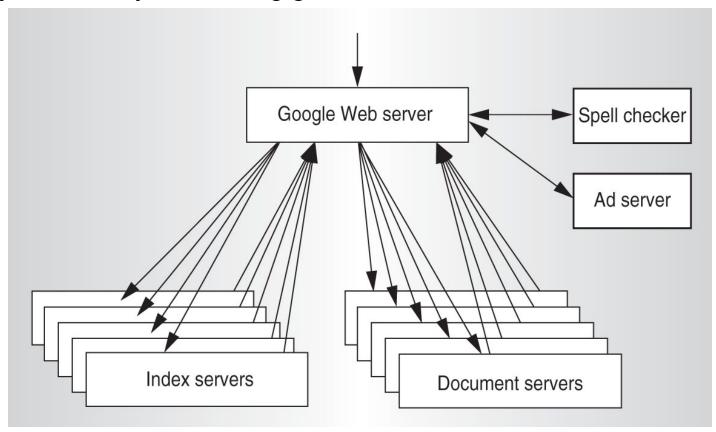
Example: Google search infrastructure

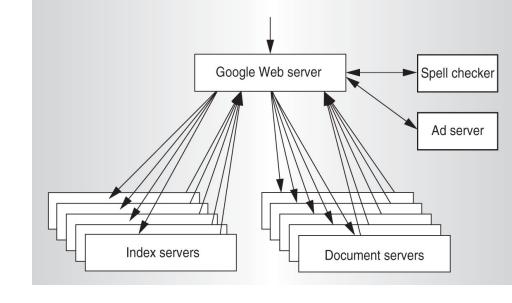
- It's likely that Google has several million machines
 - But let's be conservative 1,000,000 machines
 - A rack holds 176 CPUs (88 1U dual-processor boards), so that's about 6,000 racks
 - A rack requires about 50 square feet (given datacenter cooling capabilities), so that's about 300,000 square feet of machine room space (more than 6 football fields of real estate – although of course Google divides its machines among dozens of datacenters all over the world)
 - A rack requires about 10kw to power, and about the same to cool, so that's about 120,000 kw of power, or nearly 100,000,000 kwh per month (\$10 million at \$0.10/kwh)
 - Equivalent to about 20% of Seattle City Light's generating capacity

- There are multiple clusters (of thousands of computers each) all over the world
- Many hundreds of machines are involved in a <u>single</u>
 Google search request (remember, the web is 400+TB)
- 1. DNS routes your search request to a nearby cluster

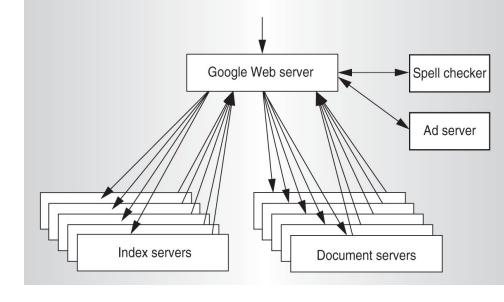


- A cluster consists of Google Web Servers, Index Servers, Doc Servers, and various other servers (ads, spell checking, etc.)
 - These are cheap standalone computers, rack-mounted, connected by commodity networking gear

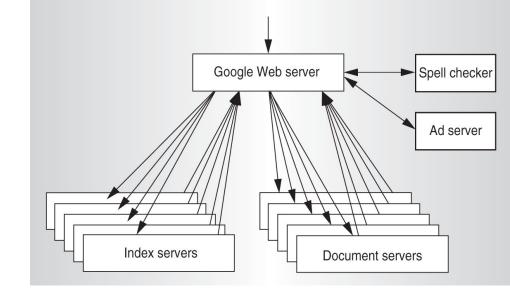




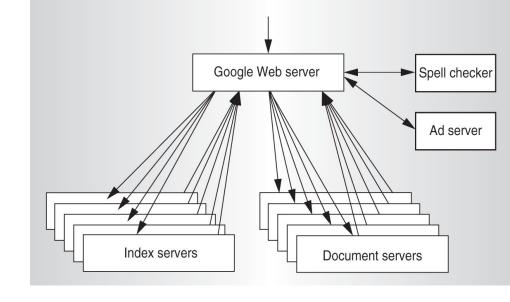
2. Within the cluster, load-balancing routes your search to a lightly-loaded Google Web Server (GWS), which will coordinate the search and response



- The index is partitioned into "shards." Each shard indexes a subset of the docs (web pages). Each shard is replicated, and can be searched by multiple computers – "index servers"
- 3. The GWS routes your search to one index server associated with each shard, through another load-balancer
- 4. When the dust has settled, the result is an ID for every doc satisfying your search, rank-ordered by relevance



- The docs, too, are partitioned into "shards" the
 partitioning is a hash on the doc ID. Each shard
 contains the full text of a subset of the docs. Each shard
 can be searched by multiple computers "doc servers"
- 5. The GWS sends appropriate doc IDs to one doc server associated with each relevant shard
- 6. When the dust has settled, the result is a URL, a title, and a summary for every relevant doc



- 7. Meanwhile, the ad server has done its thing, the spell checker has done its thing, etc.
- 8. The GWS builds an HTTP response to your search and ships it off
- Many hundreds of computers have enabled you to search 400+TB of web in ~100 ms.

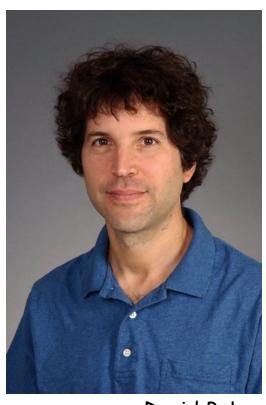
Google: The Big Picture

- Enormous volumes of data
- Extreme parallelism
- The cheapest imaginable components
 - Failures occur all the time
 - You couldn't afford to prevent this in hardware
- Software makes it
 - Fault-Tolerant
 - Highly Available
 - Recoverable
 - Consistent
 - Scalable
 - Predictable
 - Secure

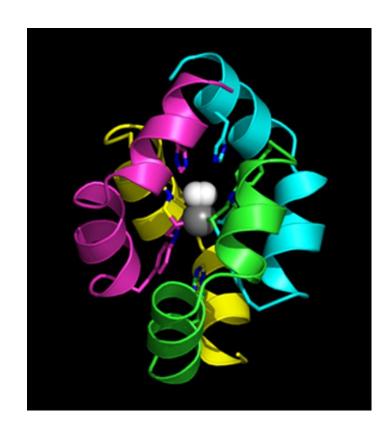
How on earth would you enable mere mortals write hairy applications such as this?

- Recognize that many Google applications have the same structure
 - Apply a "map" operation to each logical record in order to compute a set of intermediate key/value pairs
 - Apply a "reduce" operation to all the values that share the same key in order to combine the derived data appropriately
- Build a runtime library that handles all the details, accepting a couple of customization functions from the user – a Map function and a Reduce function
- That's what MapReduce is
 - Supported by the Google File System and the Chubby lock manager
 - Augmented by the BigTable not-quite-a-database system

An extremely loosely coupled system: BOINC



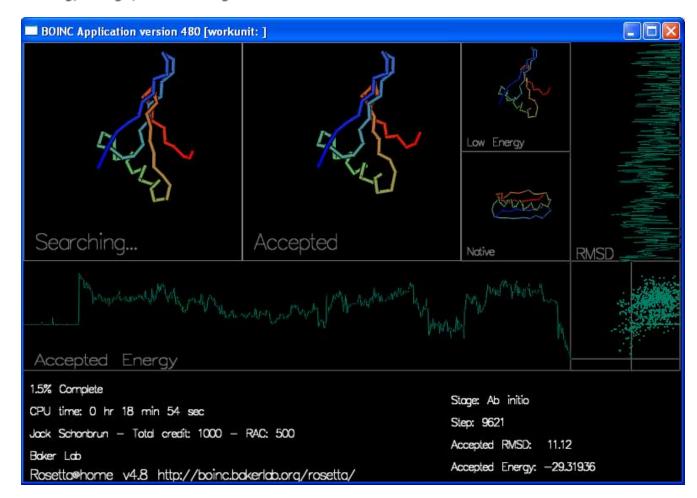
David Baker





Rosetta@home

Protein Folding, Design, and Docking

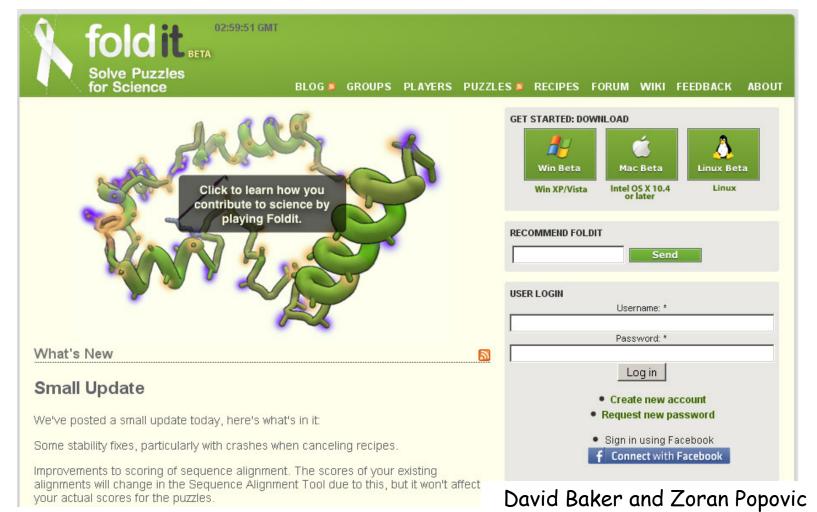




Berkeley Open Infrastructure for Network Computing



Totally off the subject of OS: Human Computation



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BootsMcGraw

Global Soloist Rank: #6 Global Soloist Score: 3784

Cases

Profile

Name: BootsMcGraw

Location: Dallas, Texas USA

Started Folding: 12/06/08

About me: An educated redneck here, from Dallas, Texas.

When I was in grad school in 1985 at the State University of New York at Buffalo, my master's thesis was to construct and present a computer program that predicted the secondary structures (helix, sheet, loop) of proteins based on their amino acid sequences. Tertiary structure (i.e. folding) prediction was

a pie-in-the-sky fantasy.

Imagine my delight, a quarter century later, to find out that not only are people determining tertiary structures of proteins, but they've made a *game* of it.

Hobbies: Licensed Massage Therapist; also a photographer, videographer, and

webmaster. I have studied health and nutrition for over twenty years. Ask me

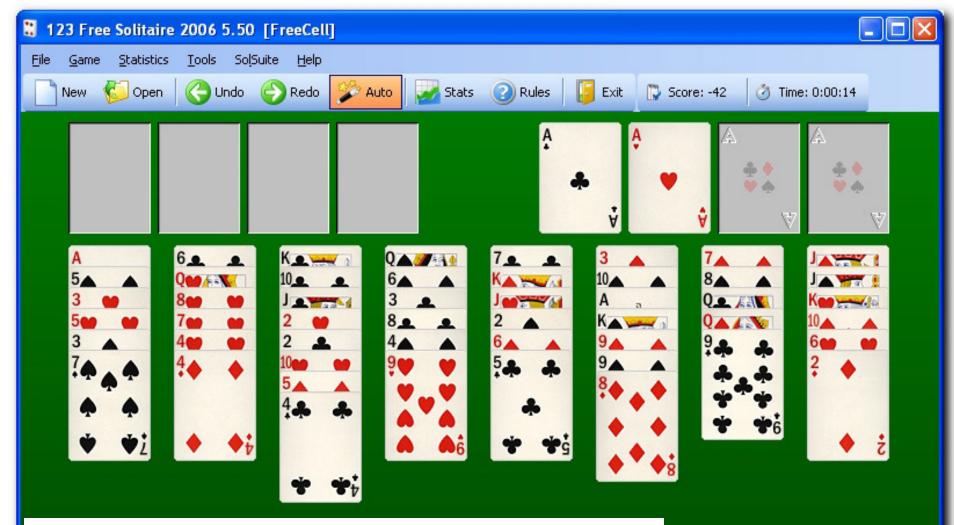
my opinions about the subject.

Group: Contenders



Luis von Ahn

- Humans and computers have different computational strengths
- Can we exploit these differences?
 - To differentiate computers from humans?
 - E.g., to make it harder for spambots to acquire new email accounts from which to send spam
 - To create human/machine computational systems that combine the best of each?

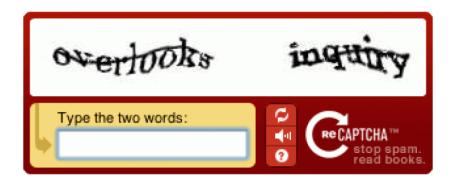


Hours per year, world-wide, spent playing computer solitaire: 9 billion

Hours spent building the Panama Canal: 20 million (less than a day of solitaire)





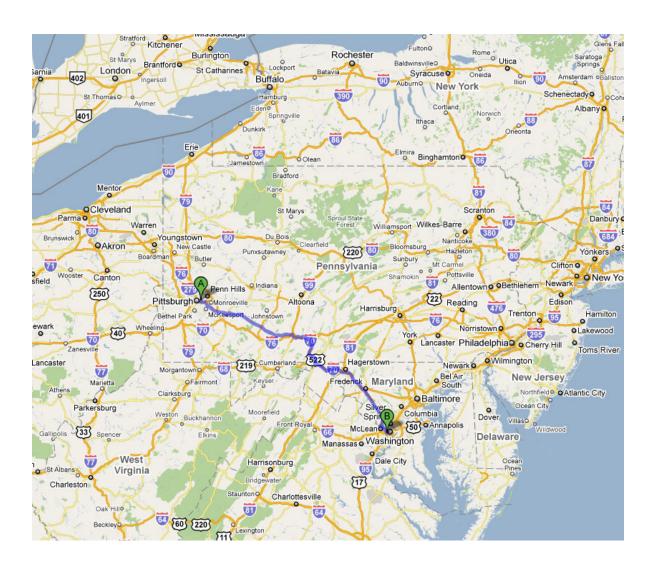


Where do the words come from?

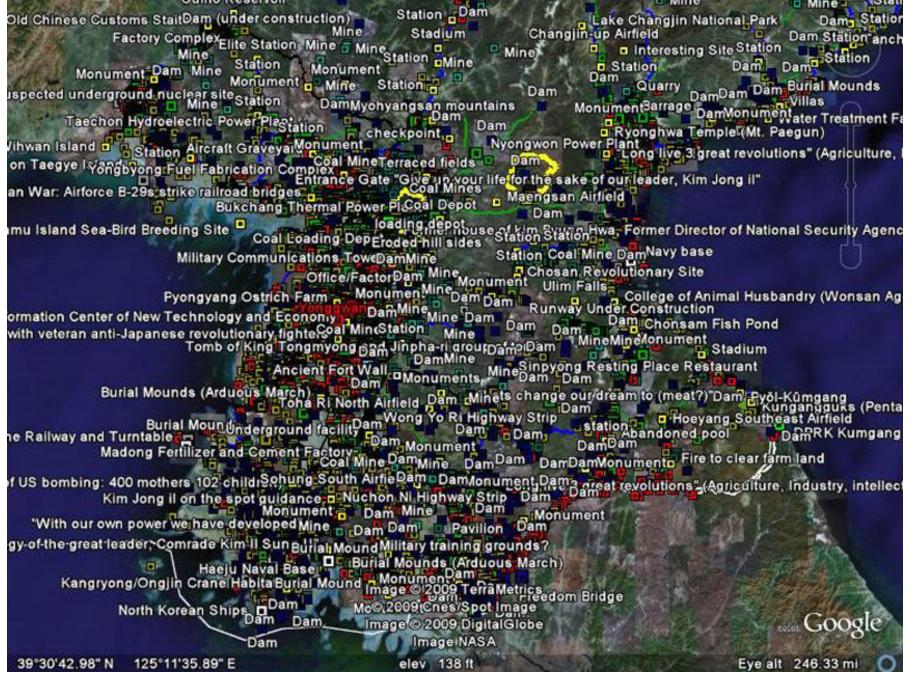
The New York Times

Entire photo archive (years 1851-1980) was completed in 2009









DARPA NETWORK CHALLENGE DARPA





[Peter Lee, DARPA]

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