CSE 451: Operating Systems Winter 2017

Module 19 File System Summary

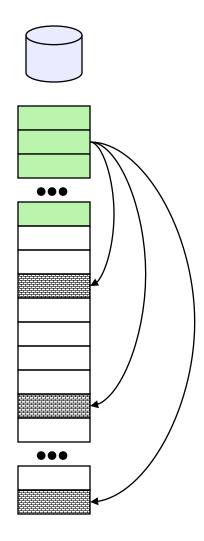
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UFS

Hardware Device

inodes

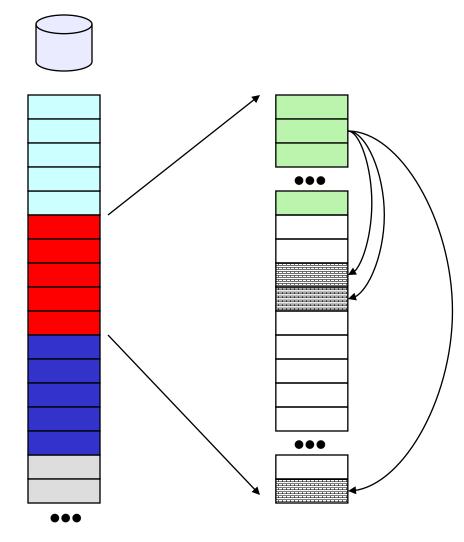
data blocks



Low throughput addressed by:

- larger blocks
- cylinder groups
- aggressive caching
- hardware awareness





Long post-crash boot times **JFS** addressed by: transactional journal of changes propagated back to "real" file system asynchronously Example: file creation Start t Alloc inode 1067 Write inode 1067 w/ [data] Write block 22731 w/ Main Journal [data] App Commit t memory requests cache file system If data block updates are not journaled, after a crash files may (To be clear: FS and FFS have garbage blocks have a cache too – I just didn't draw it.)

Write throughput addressed by: LFS • the file system is a log file 2 file1 Main Segments memory cache file1 inode maps © 2013 Gribble, Lazowska, Levy, Zahorjan 5

Supporting Multiple File Systems: vfs

