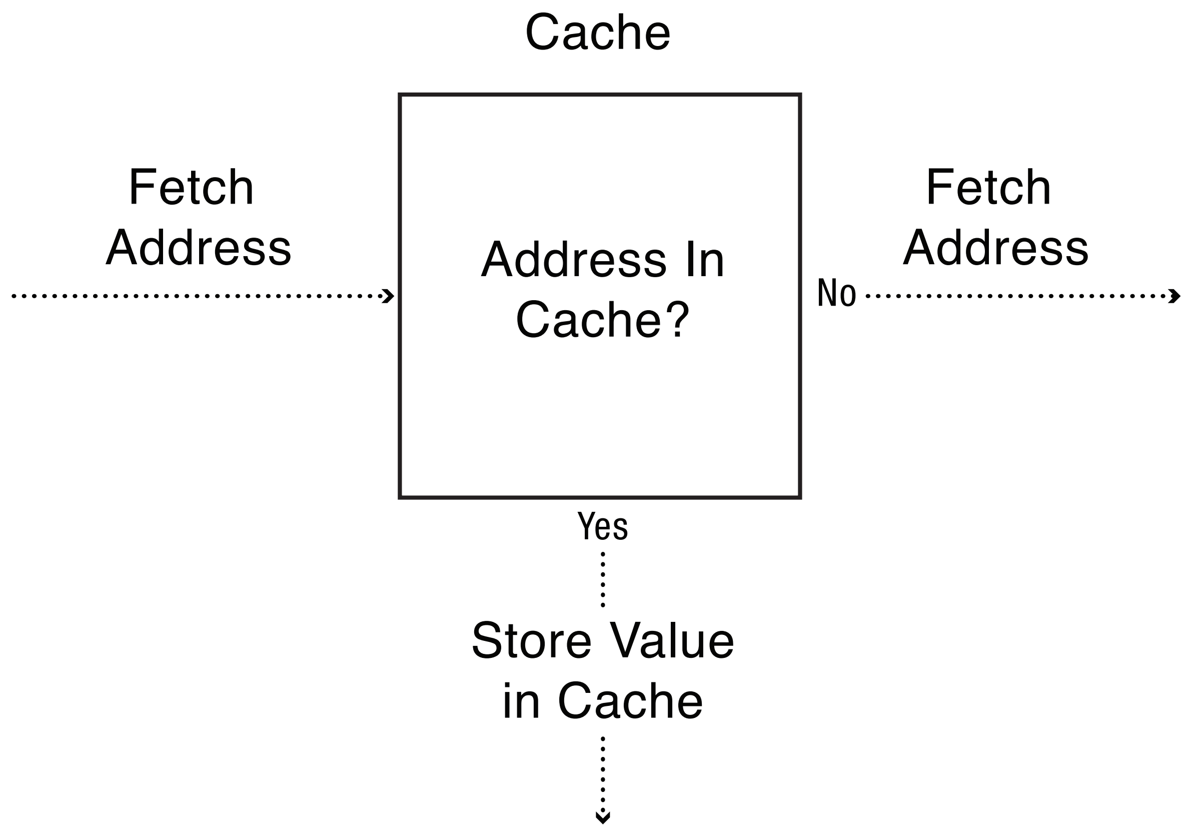
Caching and Demand-Paged Virtual Memory

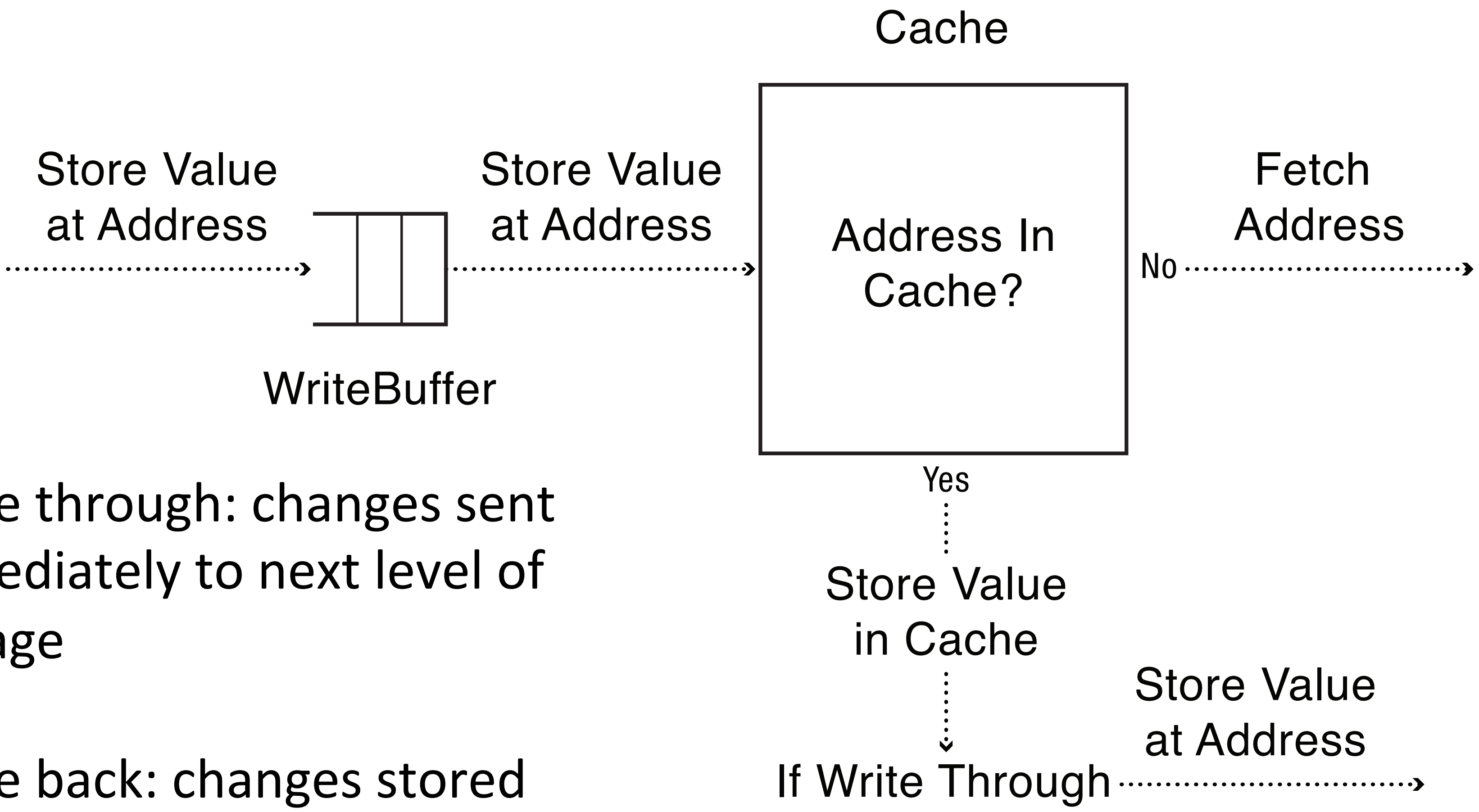
Definitions

- Cache
  - Copy of data that is faster to access than the original
  - Hit: if cache has copy
  - Miss: if cache does not have copy
- Cache block
  - Unit of cache storage (multiple memory locations)
- Temporal locality
  - Programs tend to repeatedly reference the same memory locations
  - Example: instructions in a loop
- Spatial locality
  - Programs tend to reference nearby locations
  - Example: data in a loop

Cache Concept (Read)



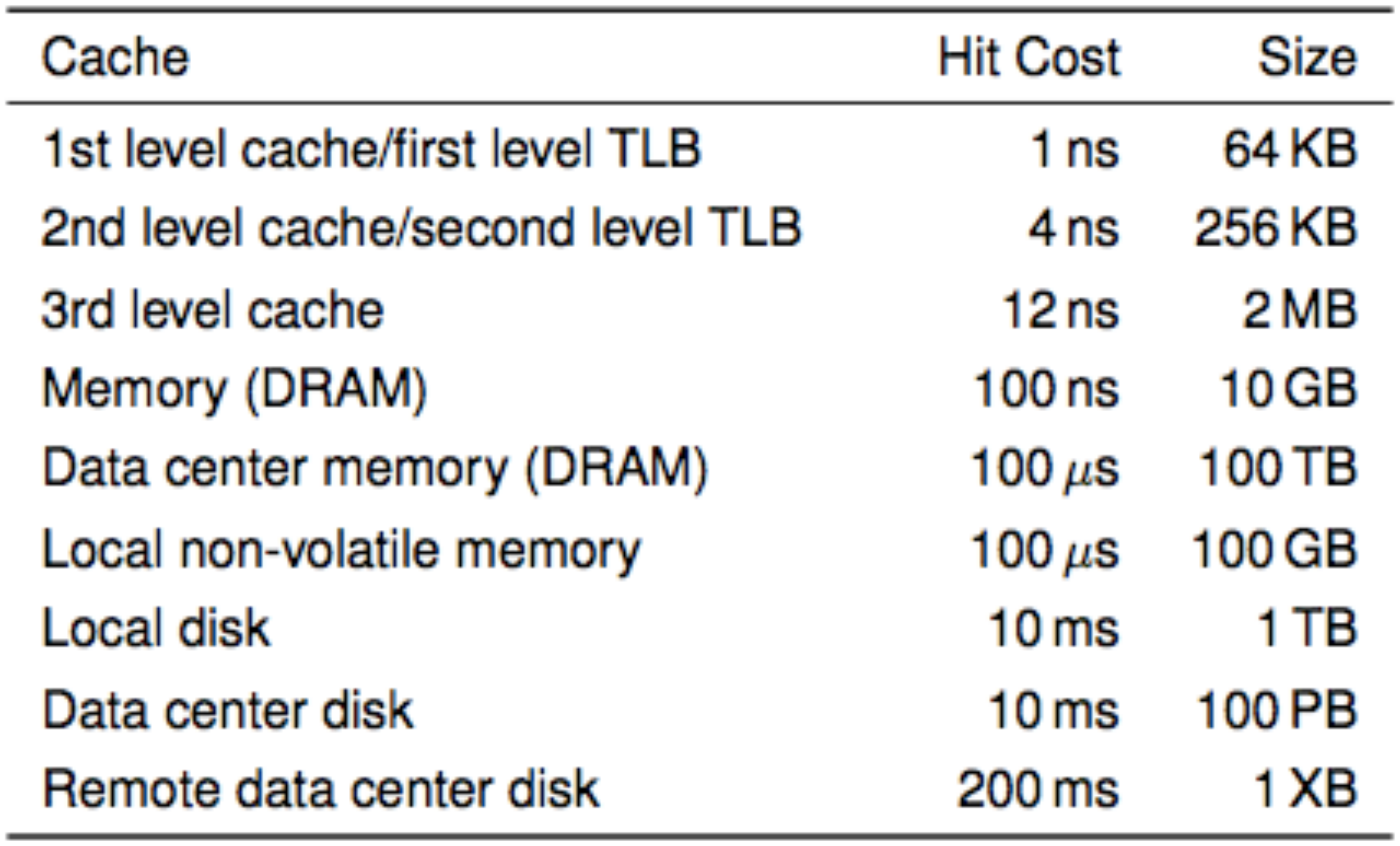
Cache Concept (Write)



Write through: changes sent immediately to next level of storage

Write back: changes stored in cache until cache block is replaced

Memory Hierarchy



Cache	Hit Cost	Size
1st level cache/first level TLB	1 ns	64 KB
2nd level cache/second level TLB	4 ns	256 KB
3rd level cache	12 ns	2 MB
Memory (DRAM)	100 ns	10 GB
Data center memory (DRAM)	100 μs	100 TB
Local non-volatile memory	100 μs	100 GB
Local disk	10 ms	1 TB
Data center disk	10 ms	100 PB
Remote data center disk	200 ms	1 XB

i7 has 8MB as shared 3rd level cache; 2nd level cache is per-core

Main Points

- Can we provide the illusion of near infinite memory in limited physical memory?
  - Demand-paged virtual memory
  - Memory-mapped files
- How do we choose which page to replace?
  - FIFO, MIN, LRU, LFU, Clock
- What types of workloads does caching work for, and how well?
  - Spatial/temporal locality vs. Zipf workloads

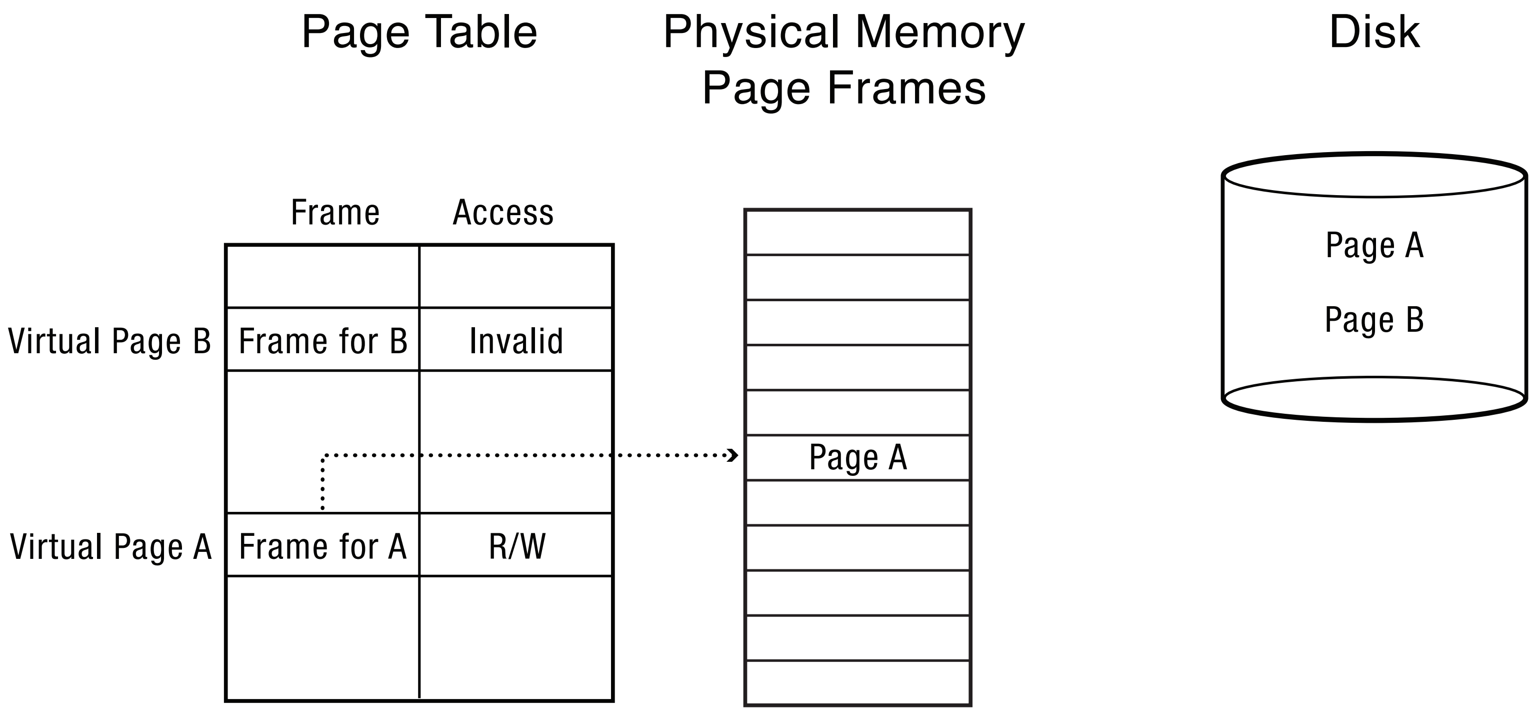
Hardware address translation is a power tool

- Kernel trap on read/write to selected addresses
  - Copy on write
  - Fill on reference
  - Zero on use
  - Demand paged virtual memory
  - Memory mapped files
  - Modified bit emulation
  - Use bit emulation

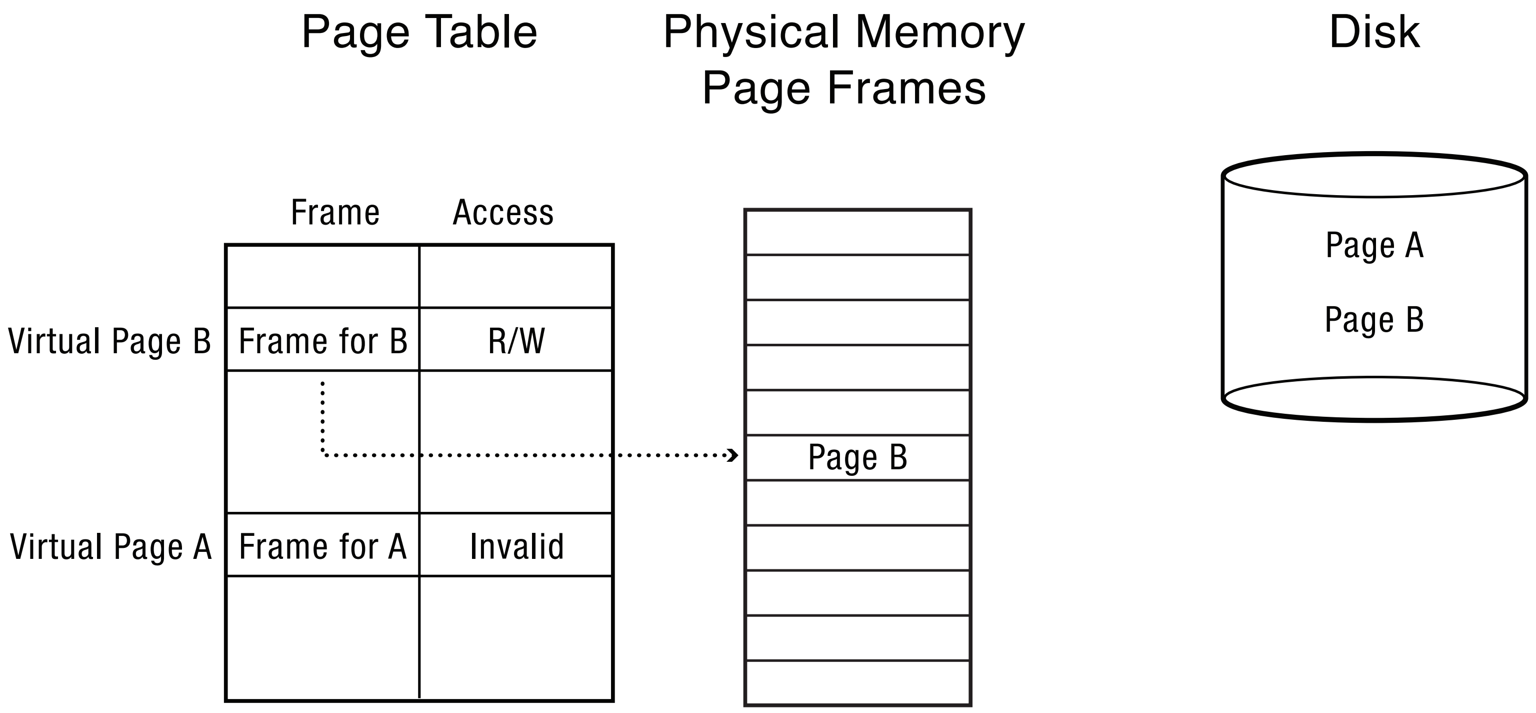
Demand Paging

- Illusion of (nearly) infinite memory, available to every process
- Multiplex virtual pages onto a limited amount of physical page frames
- Pages can be either
  - resident (in physical memory, valid page table entry)
  - non-resident (on disk, invalid page table entry)
- On reference to non-resident page, copy into memory, replacing some resident page
  - From the same process, or a different process

Demand Paging (Before)



Demand Paging (After)



Demand Paging Questions

- How does the kernel provide the illusion that all pages are resident?
- Where are non-resident pages stored on disk?
- How do we find a free page frame?
- Which pages have been modified (must be written back to disk) or actively used (shouldn't be evicted)?
- Are modified/use bits virtual or physical?
- What policy should we use for choosing which page to evict?

Demand Paging on MIPS

1. TLB miss
2. Trap to kernel
3. Page table walk
4. Find page is invalid
5. Locate page on disk
6. Allocate page frame
   - Evict page if needed
7. Initiate disk block read into page frame
8. Disk interrupt when DMA complete
9. Mark page as valid
10. Load TLB entry
11. Resume process at faulting instruction
12. Execute instruction

Demand Paging

1. TLB miss
2. Page table walk
3. Page fault (page invalid in page table)
4. Trap to kernel
5. Locate page on disk
6. Allocate page frame
   – Evict page if needed
7. Initiate disk block read into page frame
8. Disk interrupt when DMA complete
9. Mark page as valid
10. Resume process at faulting instruction
11. TLB miss
12. Page table walk to fetch translation
13. Execute instruction

Locating a Page on Disk

- When a page is non-resident, how do we know where to find it on disk?
- Option: Reuse page table entry
  - If resident, page frame
  - If non-resident, disk sector
- Option: Use file system
  - Code pages: executable image (read-only)
  - Data/Heap/Stack: per-segment file in file system, offset in file = offset within segment

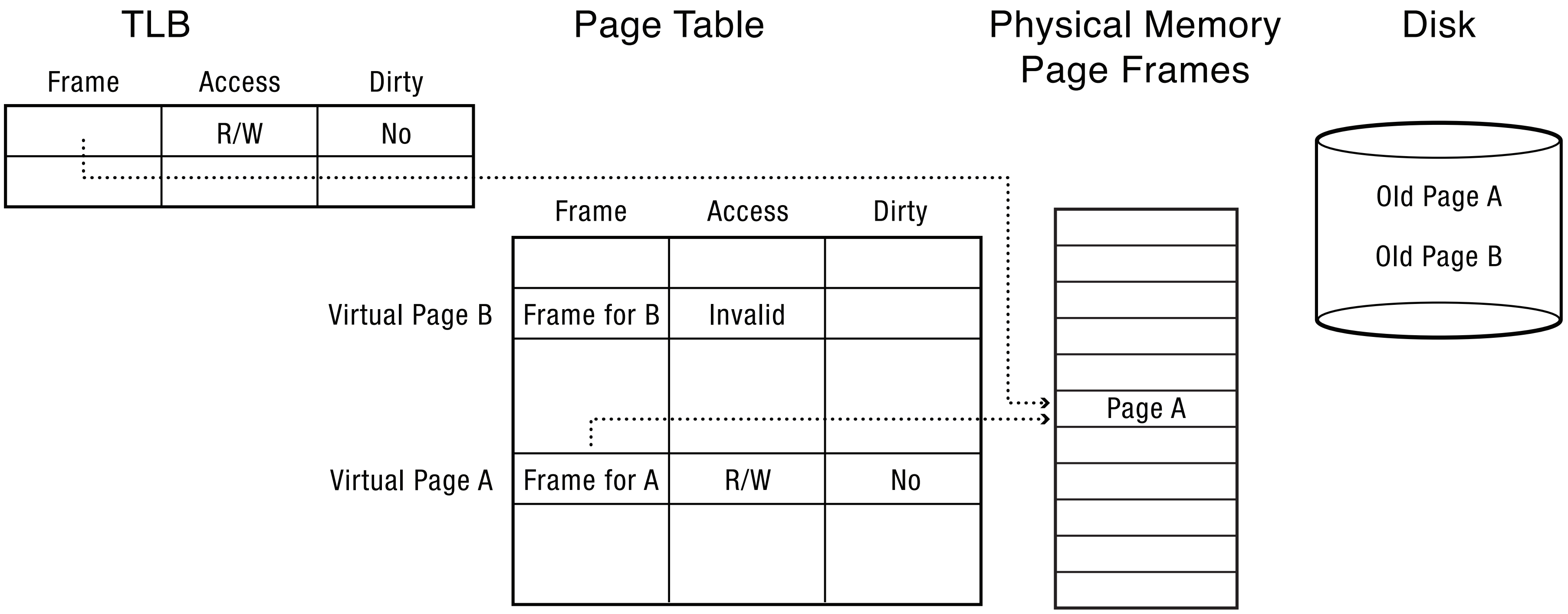
Allocating a Page Frame

- Select old page to evict
- Find all page table entries that refer to old page
  - If page frame is shared (hint: use a coremap)
- Set each page table entry to invalid
- Remove any TLB entries
  - Copies of now invalid page table entry
- Write changes on page back to disk, if necessary

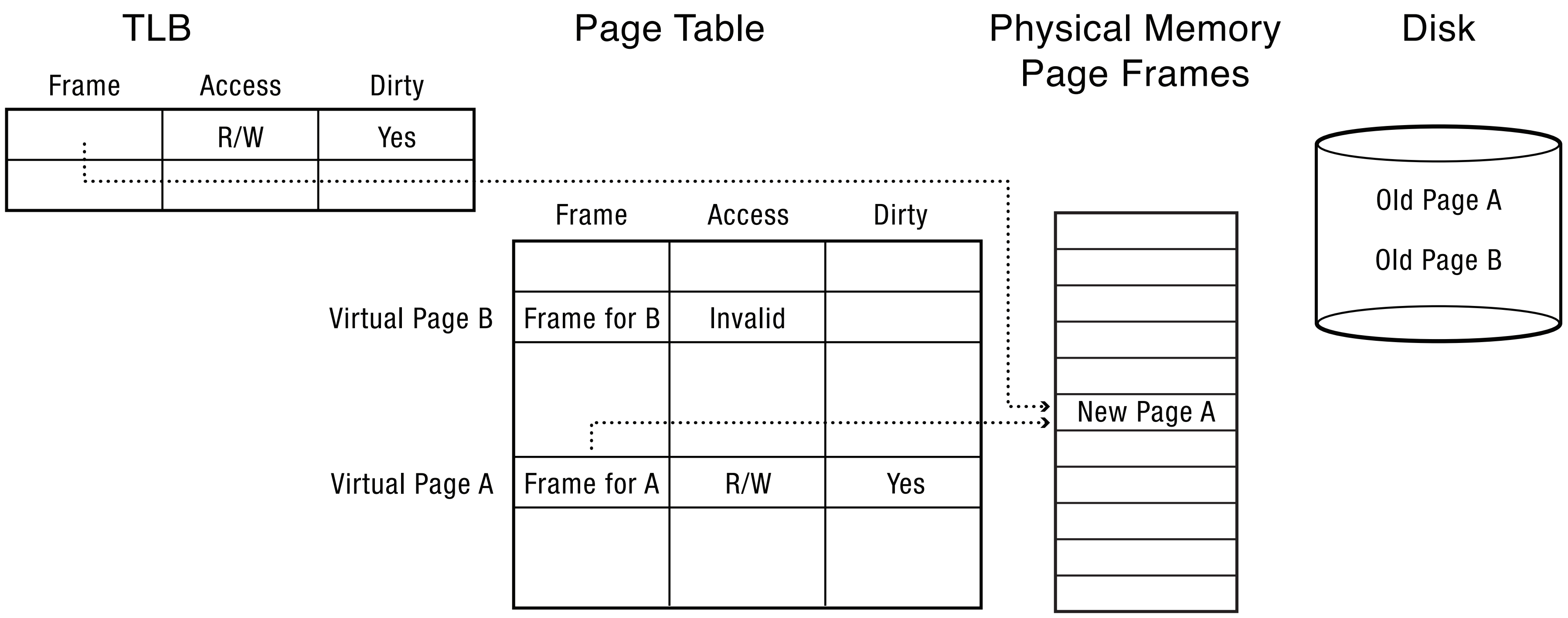
Has page been modified/recently used?

- Every page table entry has some bookkeeping
  - Has page been modified?
    - Set by hardware on store instruction
    - In both TLB and page table entry
  - Has page been recently used?
    - Set by hardware on in page table entry on every TLB miss
- Bookkeeping bits can be reset by the OS kernel
  - When changes to page are flushed to disk
  - To track whether page is recently used

Tracking Page Modifications (Before)



Tracking Page Modifications (After)



Modified/Use Bits are (often) Virtual

- Most machines keep modified/use bits in the page table entry – why?
- Physical page is
  - Modified if *any* page table entry that points to it is modified
  - Recently used if *any* page table entry that points to it is recently used
- On MIPS, simpler to keep modified/use bits in the core map (map of physical page frames)

Use Bits are Fuzzy

- Modified bit must be ground truth
  - What happens if we evict a modified page without writing the changes back to disk?
- Use bit can be approximate
  - What happens if we evict a page that is currently being used?
  - “Evict any page not used for a while” is nearly as good as “evict the single page not used for the longest”

Emulating Modified/Use Bits w/ MIPS Software Loaded TLB

- MIPS TLB entries can be read-only or read-write
- On a TLB read miss:
  - If page is clean (in core map), load TLB entry as read-only
  - if page is dirty, load as read-write
  - Mark page as recently used in core map
- On a TLB write to an unmodified page:
  - Mark page as modified/recently used in core map
  - Reset TLB entry to be read-write
- On TLB write miss:
  - Mark page as modified/recently used in core map
  - Load TLB entry as read-write

Emulating a Modified Bit (Hardware Loaded TLB)

- Some processor architectures do not keep a modified bit per page
  - Extra bookkeeping and complexity
- Kernel can *emulate* a modified bit:
  - Set all clean pages as read-only
  - On first write to page, trap into kernel
  - Kernel set modified bit in core map
  - Kernel set page table entry as read-write
  - Resume execution
- Kernel needs to keep track
  - Current page table permission (e.g., read-only)
  - True page table permission (e.g., writeable, clean)

Emulating a Recently Used Bit (Hardware Loaded TLB)

- Some processor architectures do not keep a recently used bit per page
  - Extra bookkeeping and complexity
- Kernel can emulate a recently used bit:
  - Set all recently unused pages as invalid
  - On first read or write, trap into kernel
  - Kernel set recently used bit in core map
  - Kernel mark page table entry as read or read/write
  - Resume execution
- Kernel needs to keep track
  - Current page table permission (e.g., invalid)
  - True page table permission (e.g., read-only, writeable)

Models for Application File I/O

- Explicit read/write system calls
  - Data copied to user process using system call
  - Application operates on data
  - Data copied back to kernel using system call
- Memory-mapped files
  - Open file as a memory segment
  - Program uses load/store instructions on segment memory, implicitly operating on the file
  - Page fault if portion of file is not yet in memory
  - Kernel brings missing blocks into memory, restarts process

Advantages to Memory-mapped Files

- Programming simplicity, esp for large files
  - Operate directly on file, instead of copy in/copy out
- Zero-copy I/O
  - Data brought from disk directly into page frame
- Pipelining
  - Process can start working before all the pages are populated
- Interprocess communication
  - Shared memory segment vs. temporary file

Implementing Memory-Mapped Files

- Memory mapped file is a (logical) segment
  - Per segment access control (read-only, read-write)
- File pages brought in on demand
  - Using page fault handler
- Modifications written back to disk on eviction, file close
  - Using per-page modified bit
- Transactional (atomic, durable) updates to memory mapped file requires more mechanism

From Memory-Mapped Files to Demand-Paged Virtual Memory

- Every process segment backed by a file on disk
  - Code segment -> code portion of executable
  - Data, heap, stack segments -> temp files
  - Shared libraries -> code file and temp data file
  - Memory-mapped files -> memory-mapped files
  - When process ends, delete temp files
- Unified memory management across file buffer and process memory

Cache Replacement Policy

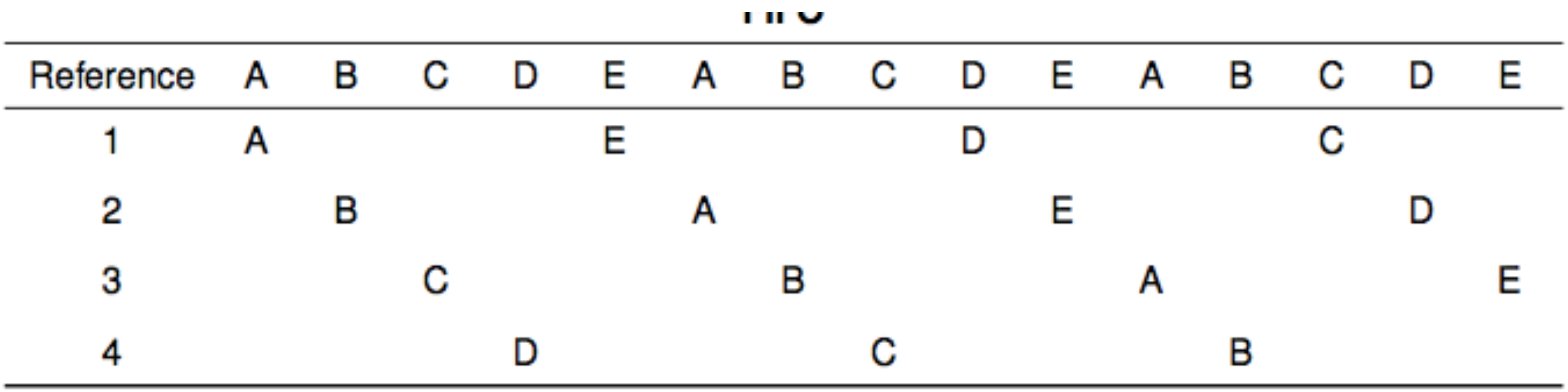
- On a cache miss, how do we choose which entry to replace?
  - Assuming the new entry is more likely to be used in the near future
  - In direct mapped caches, not an issue!

- Policy goal: reduce cache misses
  - Improve expected case performance
  - Also: reduce likelihood of very poor performance

A Simple Policy

- Random?
  - Replace a random entry
- FIFO?
  - Replace the entry that has been in the cache the longest time
  - What could go wrong?

FIFO in Action



Reference	A	B	C	D	E	A	B	C	D	E	A	B	C	D	E
1	A				E				D				C		
2		B				A				E				D	
3			C				B				A				E
4				D				C				B			

Worst case for FIFO is if program strides through memory that is larger than the cache

MIN, LRU, LFU

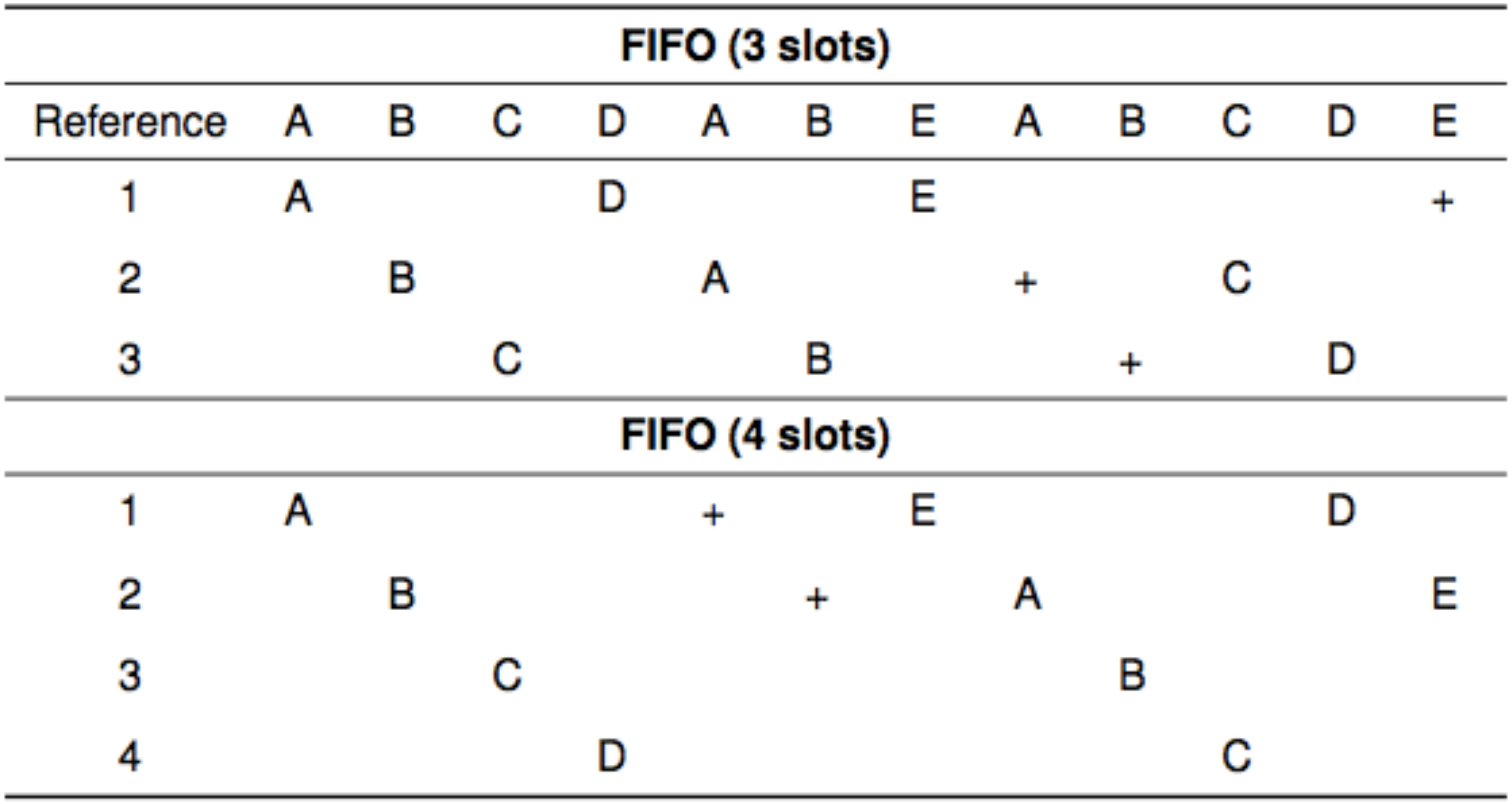
- MIN
  - Replace the cache entry that will not be used for the longest time into the future
  - Optimality proof based on exchange: if evict an entry used sooner, that will trigger an earlier cache miss
- Least Recently Used (LRU)
  - Replace the cache entry that has not been used for the longest time in the past
  - Approximation of MIN
- Least Frequently Used (LFU)
  - Replace the cache entry used the least often (in the recent past)

LRU/MIN for Sequential Scan

							LRU								
Reference	A	B	C	D	E	A	B	C	D	E	A	B	C	D	E
1	A				E				D				C		
2		B				A				E				D	
3			C				B				A				E
4				D				C				B			
							MIN								
1	A					+					+			+	
2		B					+					+	C		
3			C					+	D					+	
4				D	E					+					+

LRU															
Reference	A	B	A	C	B	D	A	D	E	D	A	E	B	A	C
1	A		+				+				+			+	
2		B			+								+		
3				C					E			+			
4						D		+		+					C
FIFO															
1	A		+				+		E						
2		B			+						A			+	
3				C								+	B		
4						D		+		+					C
MIN															
1	A		+				+				+			+	
2		B			+								+		C
3				C					E			+			
4						D		+		+					

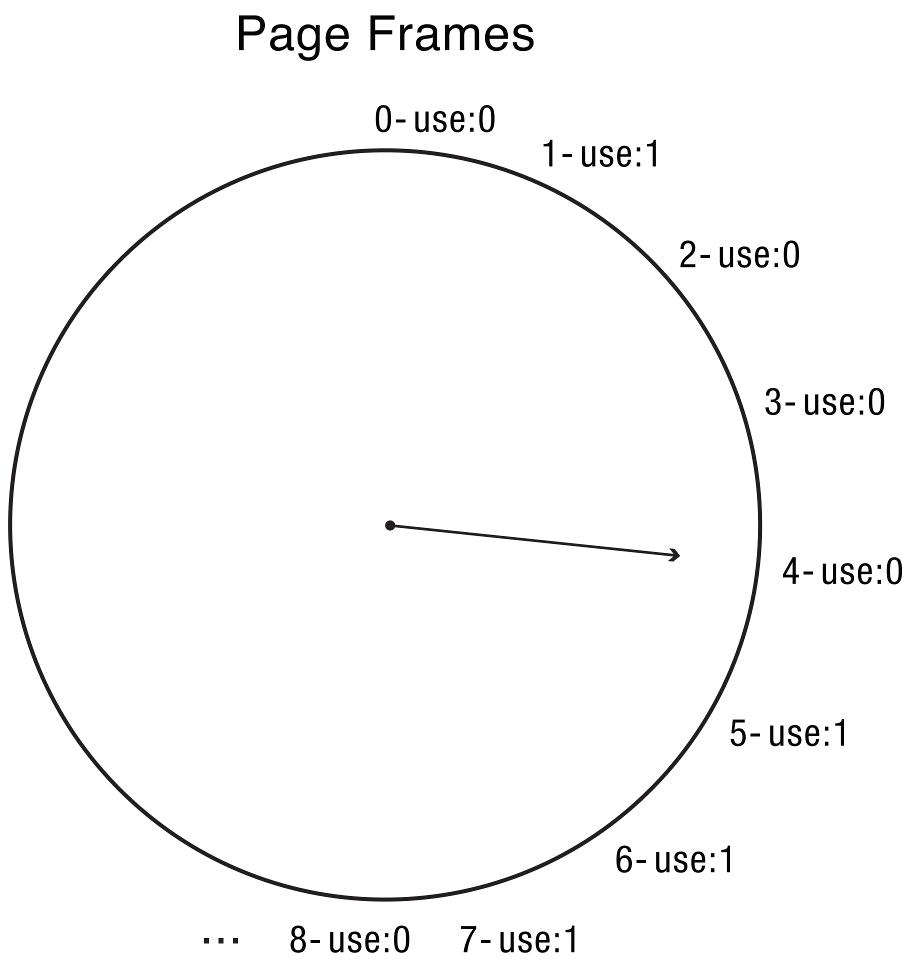
Belady’s Anomaly



						FIFO (3 slots)						
Reference	A	B	C	D	A	B	E	A	B	C	D	E
1	A			D			E					+
2		B			A			+		C		
3			C			B			+		D	
						FIFO (4 slots)						
1	A				+		E				D	
2		B				+		A				E
3			C						B			
4				D						C		

Clock Algorithm: Estimating LRU

- Periodically, sweep through all pages
- If page is unused, reclaim
- If page is used, mark as unused



Nth Chance: Not Recently Used

- Instead of one bit per page, keep an integer
  - notInUseSince: number of sweeps since last use
- Periodically sweep through all page frames

```
if (page is used) {
   notInUseSince = 0;
} else if (notInUseSince < N) {
   notInUseSince++;
} else {
    reclaim page;
}
```

Implementation Note

- Clock and Nth Chance can run synchronously
  - In page fault handler, run algorithm to find next page to evict
  - Might require writing changes back to disk first
- Or asynchronously
  - Create a thread to maintain a pool of recently unused, clean pages
  - Find recently unused dirty pages, write mods back to disk
  - Find recently unused clean pages, mark as invalid and move to pool
  - On page fault, check if requested page is in pool!
  - If not, evict that page

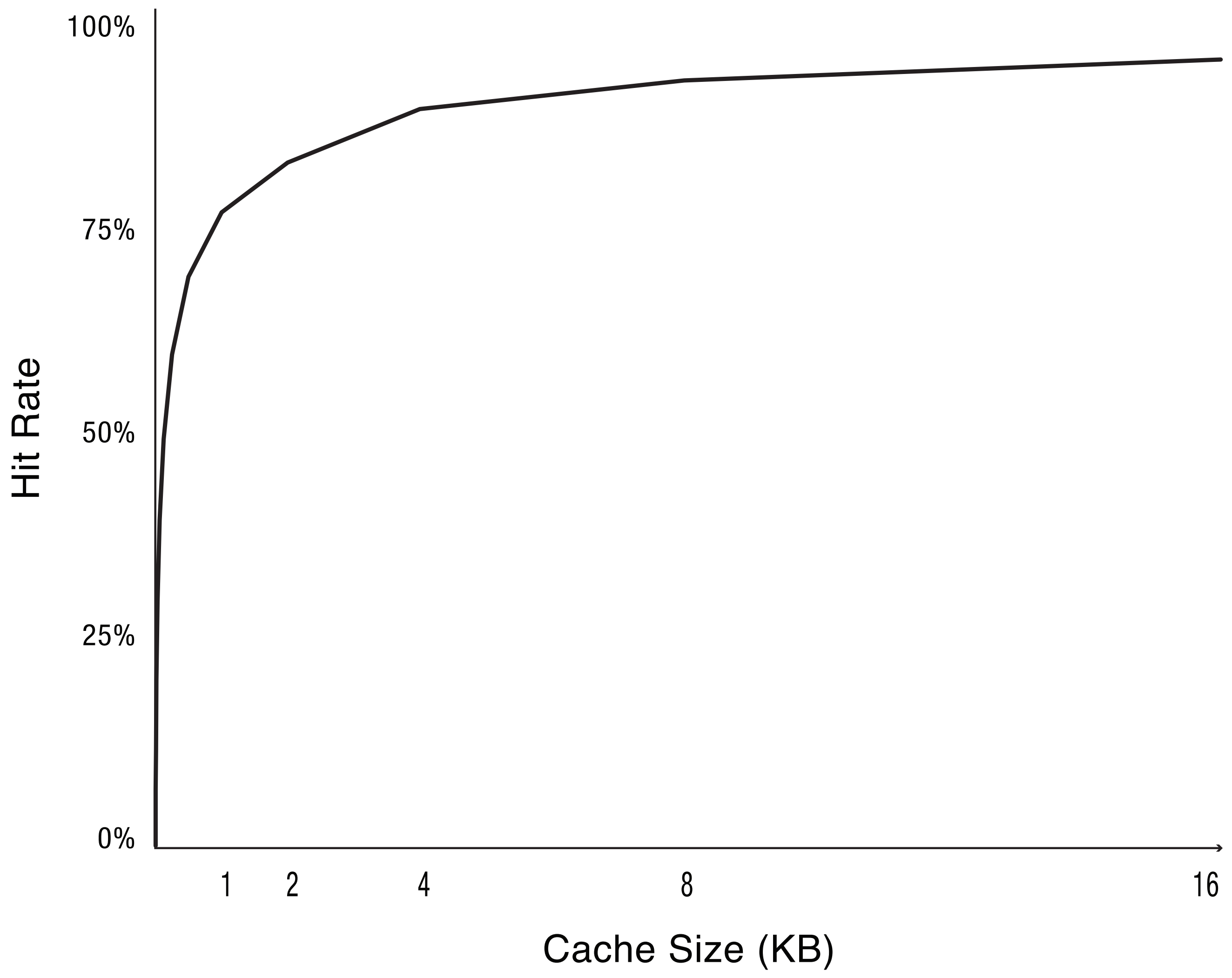
Recap

- MIN is optimal
  - replace the page or cache entry that will be used farthest into the future
- LRU is an approximation of MIN
  - For programs that exhibit spatial and temporal locality
- Clock/Nth Chance is an approximation of LRU
  - Bin pages into sets of “not recently used”

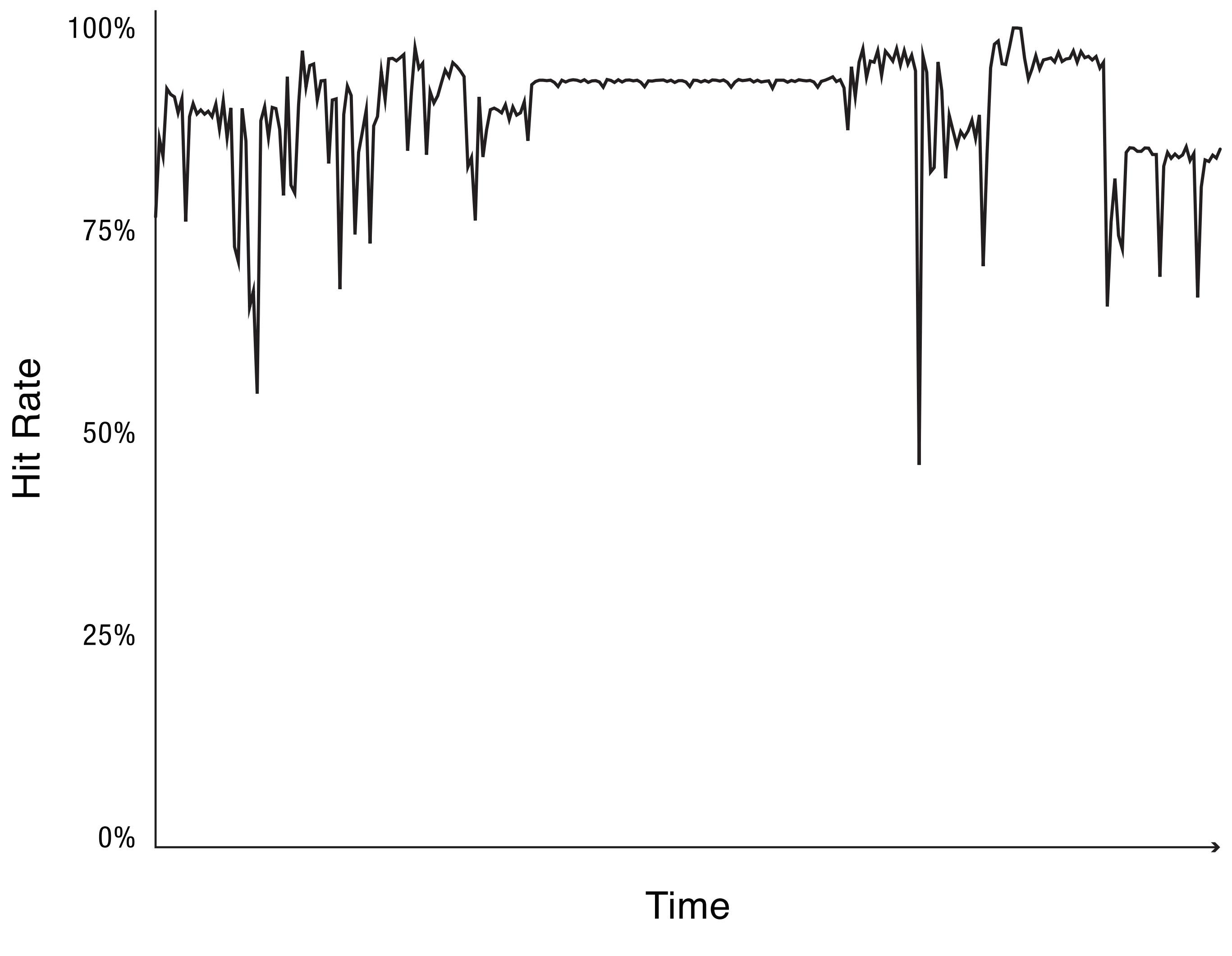
Working Set Model

- Working Set: set of memory locations that need to be cached for reasonable cache hit rate
- Thrashing: when system has too small a cache

Cache Working Set



Phase Change Behavior



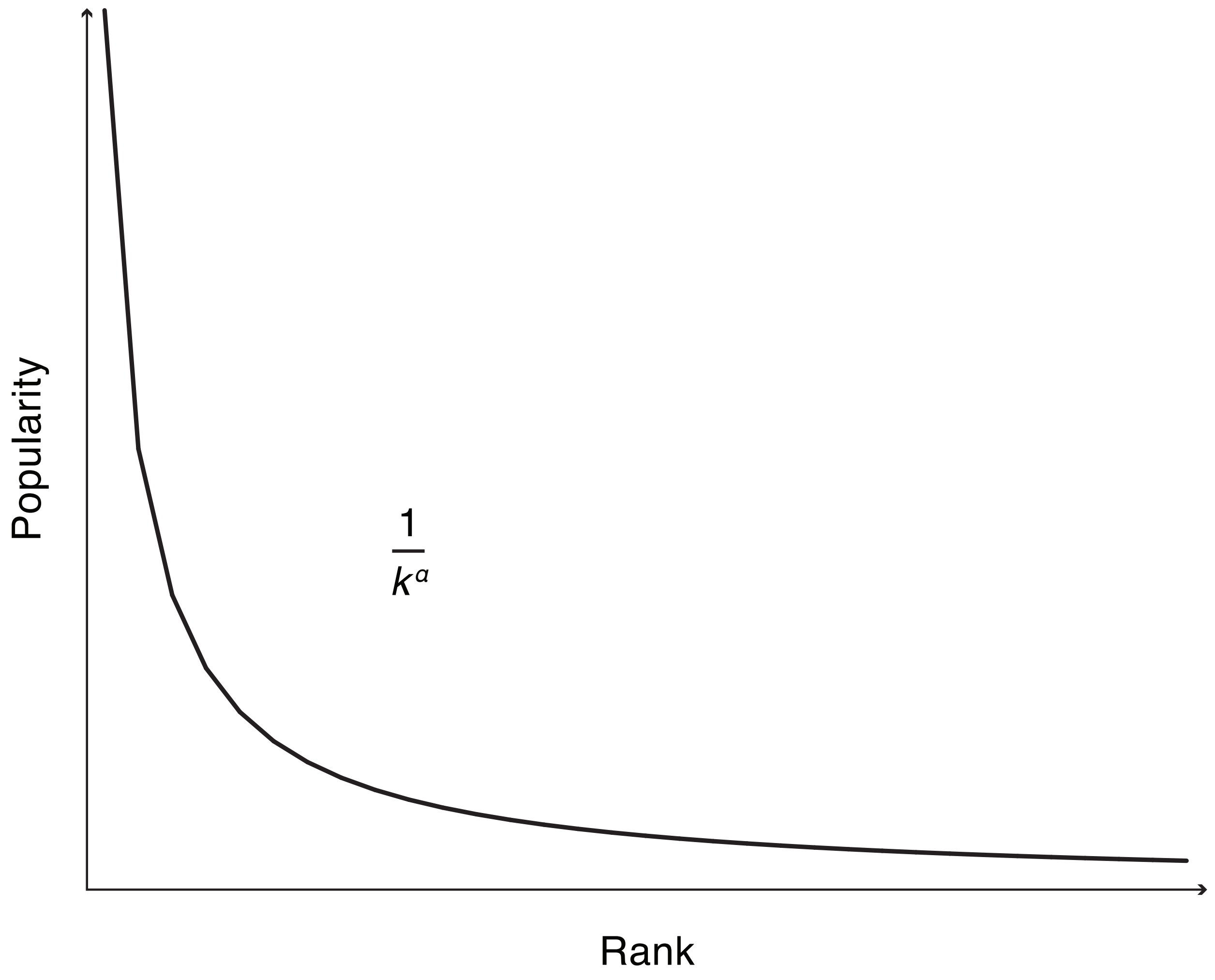
Question

- What happens to system performance as we increase the number of processes?
  - If the sum of the working sets > physical memory?

Zipf Distribution

- Caching behavior of many systems are not well characterized by the working set model
- An alternative is the Zipf distribution
  - Popularity ~ 1/k^c, for kth most popular item, 1 < c < 2

Zipf Distribution

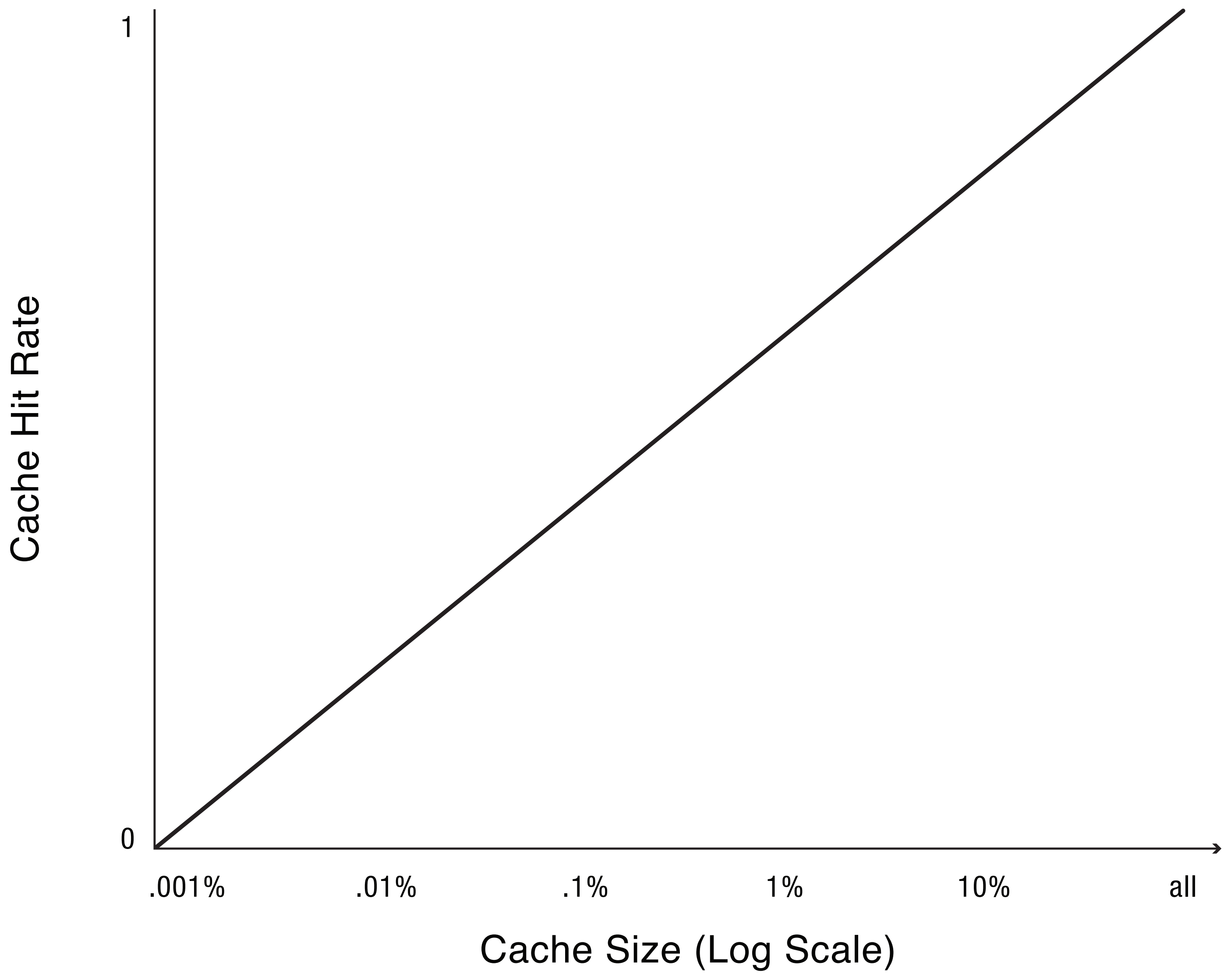


Zipf Examples

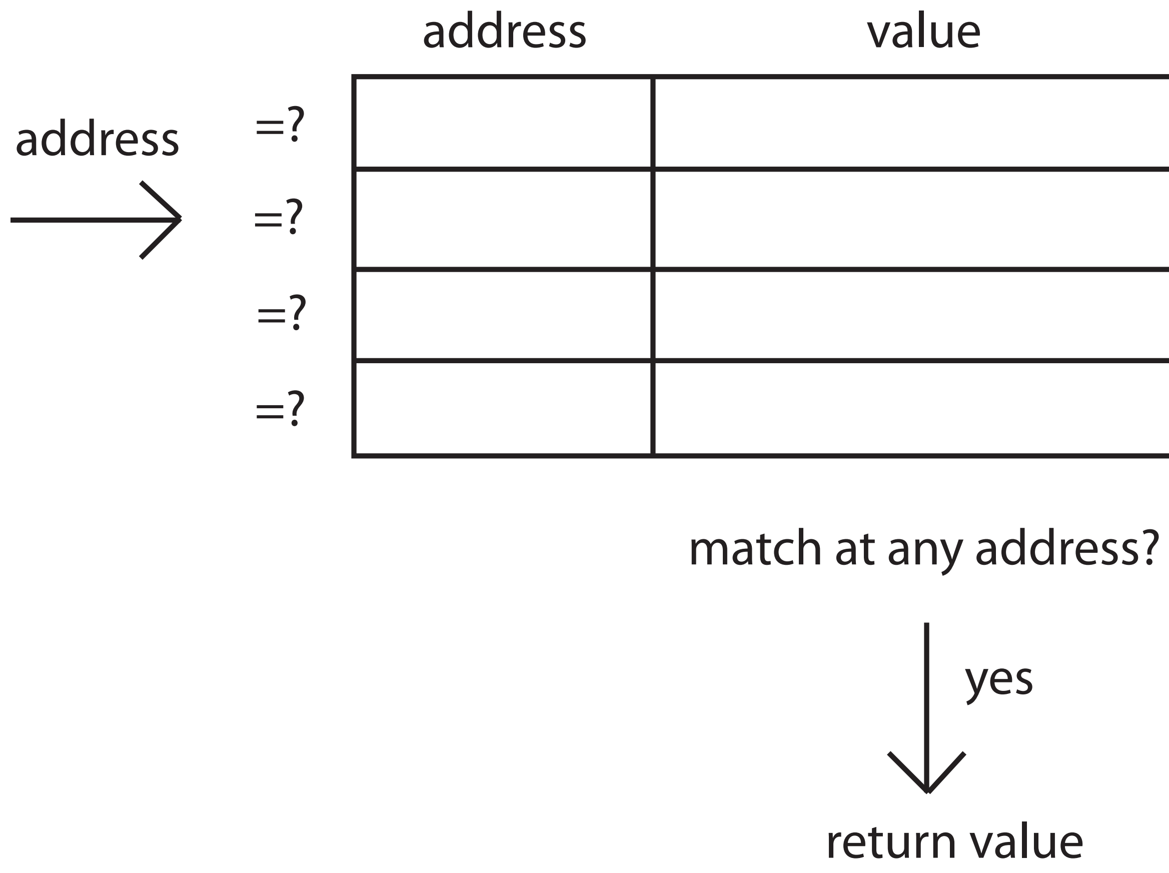
- Web pages
- Movies
- Library books
- Words in text
- Salaries
- City population
- …

Common thread: popularity is self-reinforcing

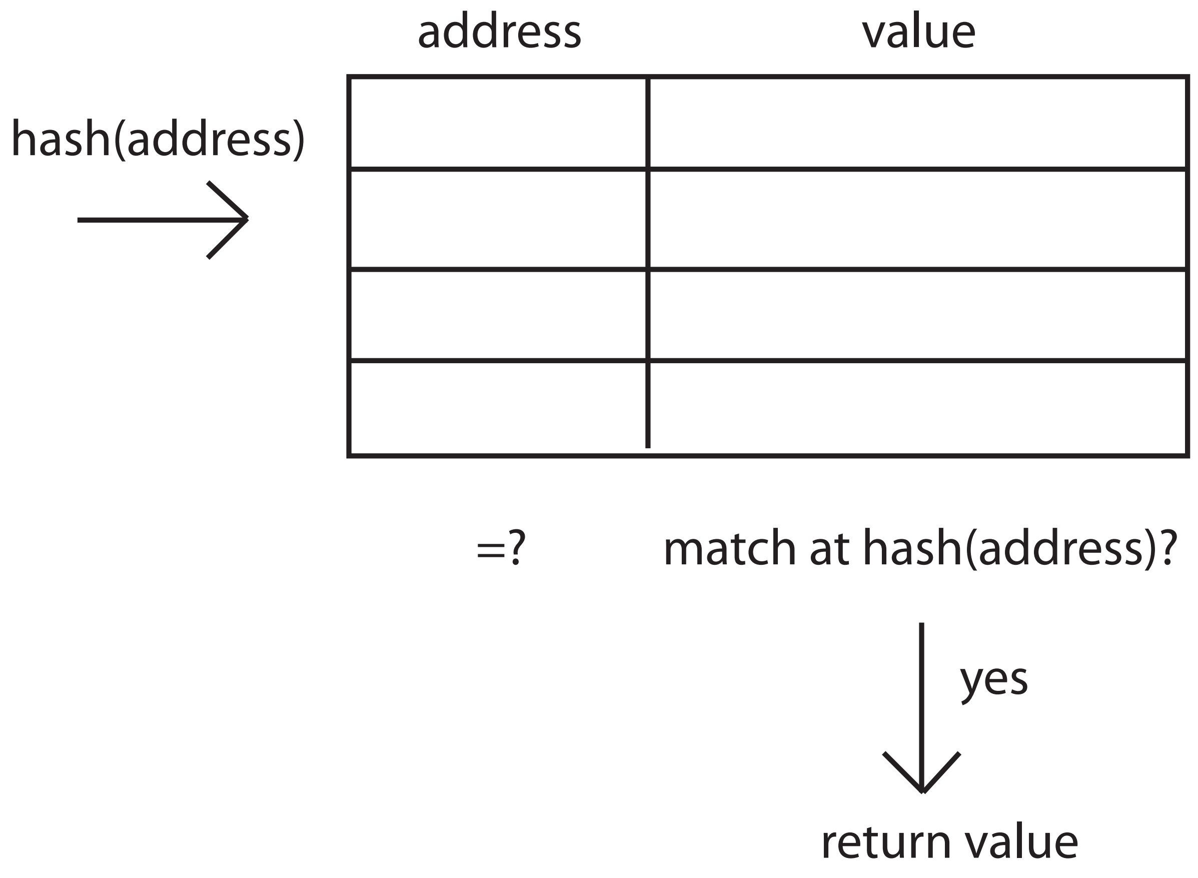
Zipf and Caching



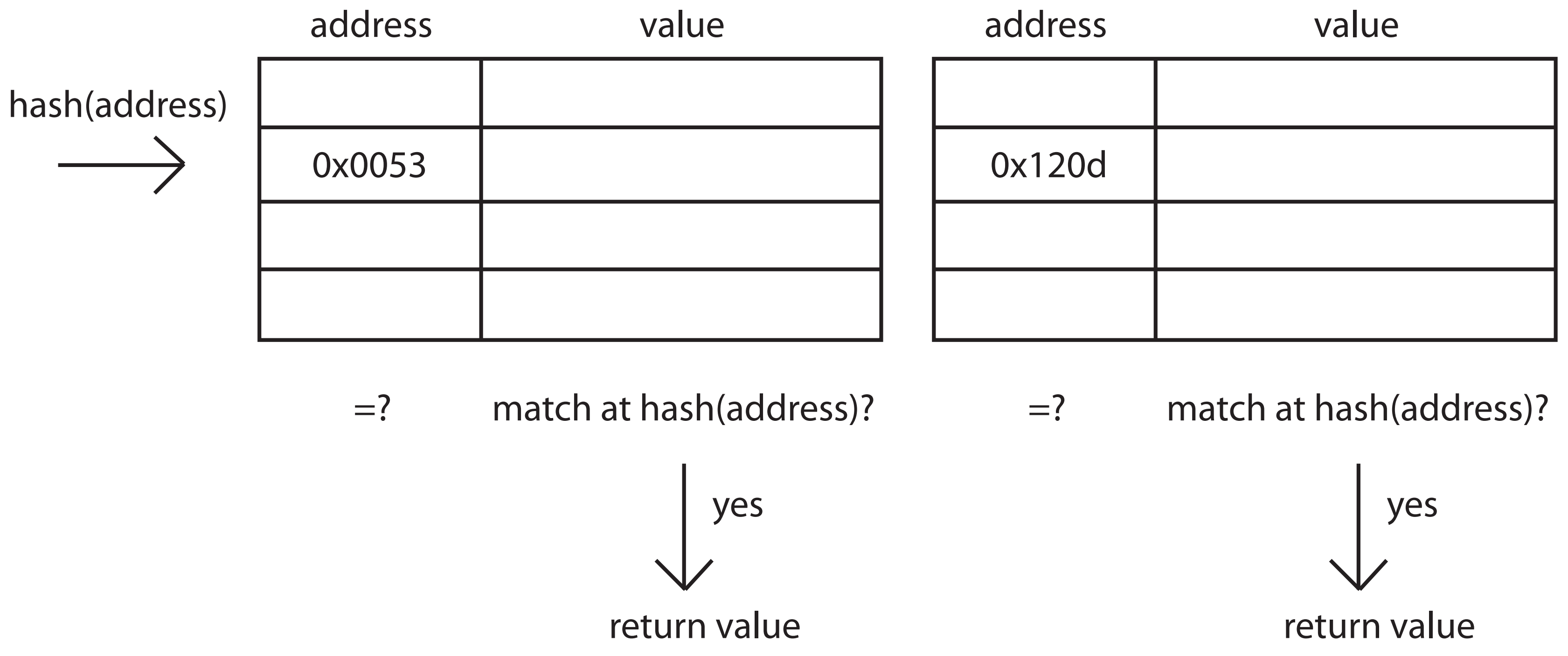
Cache Lookup: Fully Associative



Cache Lookup: Direct Mapped



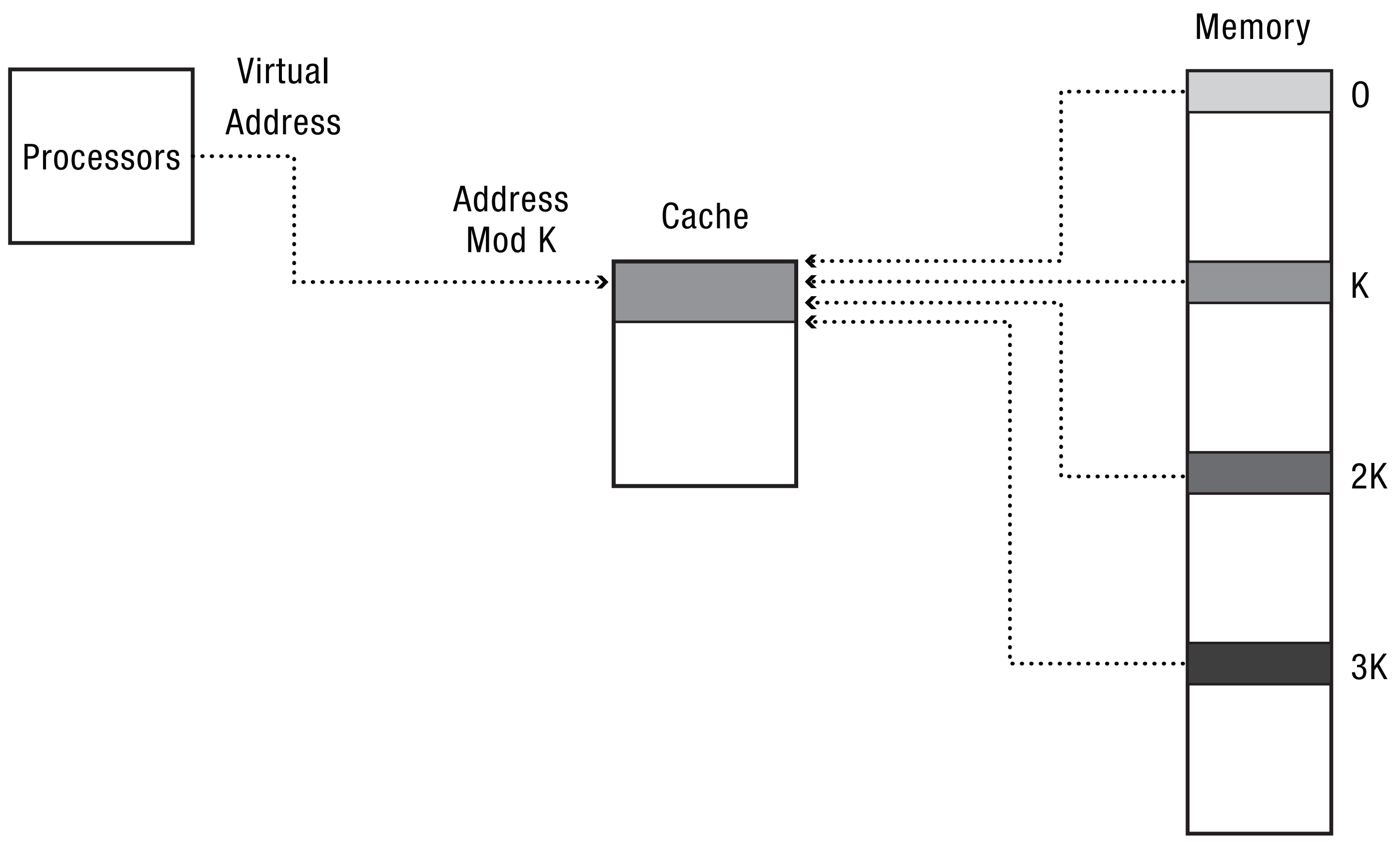
Cache Lookup: Set Associative



Page Coloring

- What happens when cache size >> page size?
  - Direct mapped or set associative
  - Multiple pages map to the same cache line
- OS page assignment matters!
  - Example: 8MB cache, 4KB pages
  - 1 of every 2K pages lands in same place in cache
- What should the OS do?

Page Coloring


Processors
Virtual
Address
Address
Mod K
Cache
Memory
0
K
2K
3K