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Topics for Today

- Project 1
- Processes
- Parallelism and Concurrency
- Permissions and Privileges

Project 1

Things to consider and remind you:

- NT functions are made from user mode, but are only executed in the kernel
- You need to monitor both the number of times NT functions are called and returned.

- These may not necessarily be the same!

Project 1

- Questions?
- Concerns?

Recap from lecture

• What is a process?

Recap from lecture

- What is a process?
 - An execution entity
 - A running instance of a program

How does an OS on a single processor hardware run multiple processes?

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How about multiple processors?

- The answer depends on implementation...
 - One master processor
 - Master processor schedules processes to itself and secondary processors

What does the OS do when there are no processes to run?

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- Run an idle process!
 - Periodically checks for any new tasks to run
 - Loops the HLT instruction to save CPU time

- Executing multiple tasks simultaneously
- This will be a focus for Project 2!

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```
NTSTATUS
NTReadFile(...) {
```

}

```
...
CSE451Info.readcalls++;
return status;
```

This is broken on a multiprocessor. Why?

```
NTSTATUS
NTReadFile(...) {
```

```
...
int tmp = CSE451Info.readcalls;
CSE451Info.readcalls = tmp+1;
return status;
}
```

This is broken on a multiprocessor. Why? Very subtle...

```
Thread 1: read "hi.txt"
```

//cse451Info.readcalls == 4

NTSTATUS NTReadFile(...) {

}

int tmp = CSE451Info.readcalls;

CSE451Info.readcalls = tmp+1; return status; Thread 2: read "foo.c"

//cse451Info.readcalls == 4

NTSTATUS NTReadFile(...) {

• • •

int tmp = CSE451Info.readcalls; CSE451Info.readcalls = tmp+1;

return status;

Thread 1: read "hi.txt"

//cse451Info.readcalls == 4

NTSTATUS NTReadFile(...) { ... int tmp = CSE451Info.readcalls;

```
CSE451Info.readcalls = tmp+1;
return status;
```

}

// CSE451Info.readcalls == 5
Should be CSE451Info.readcalls == 6!

Thread 2: read "foo.c"

//cse451Info.readcalls == 4

NTSTATUS NTReadFile(...) {

•••

int tmp = CSE451Info.readcalls; CSE451Info.readcalls = tmp+1;

return status;

}

// CSE451Info.readcalls == 5

How do we solve this consistency issue?

- Use a mutex
 - Denote start (lock) and end (unlock) of a critical section.
 - Ensures critical section only gets accessed by one thread at a time.

Thread 1: read "hi.txt"

NTSTATUS NTReadFile(...) {

•••

}

... acquire(&some_mutex); int tmp = CSE451Info.readcalls; CSE451Info.readcalls = tmp+1; release(&some_mutex); Thread 2: read "foo.c"

NTSTATUS NTReadFile(...) {

•••

acquire(&some_mutex); int tmp = CSE451Info.readcalls; CSE451Info.readcalls = tmp+1; release(&some_mutex);

}

•••

- All hardware resources must be protected.
- User access to these resources must be restricted.
- Recap: How does the OS issue protection?

Things to consider:

• Why can only the OS create processes?

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- Can the OS ever deny a user program from issuing a system call? If so, when? If not, why?

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- Why can only the OS create processes?
- Can the OS ever deny a user program from issuing a system call? If so, when? If not, why?
- Files can only be accessed by file descriptors, and not addresses. Why?

Practice Problem (homework prep!)

Which of the following are privileged instructions?

- 1. Set value of the timer.
- 2. Read the clock.
- 3. Clear memory.
- 4. Issue a trap instruction.
- 5. Switch from user to kernel mode.
- 6. Access I/O device.