CSE 451: Operating Systems

Section 4

Project 2 Intro; Threads

Project 1

- *Congratulations, you're all kernel hackers now!
- *We're going to give you a break and have you do some userspace work ©

Project 2: user-level threads

- * Part A: due Wednesday, May 2 at 11:59pm
 - * Implement part of a user thread library
 - * Add synchronization primitives
 - * Solve a synchronization problem
- * Part B: due Friday, May 18 at 11:59pm
 - * Implement a multithreaded web server
 - * Add preemption
 - * Get some results and write a (small) report

Project 2 notes

- *Start EARLY!
 - * It's looooooong
 - * Read the assignment carefully
 - * Read it again
 - * Understand the skeleton code
- *Use the same groups as for project 1

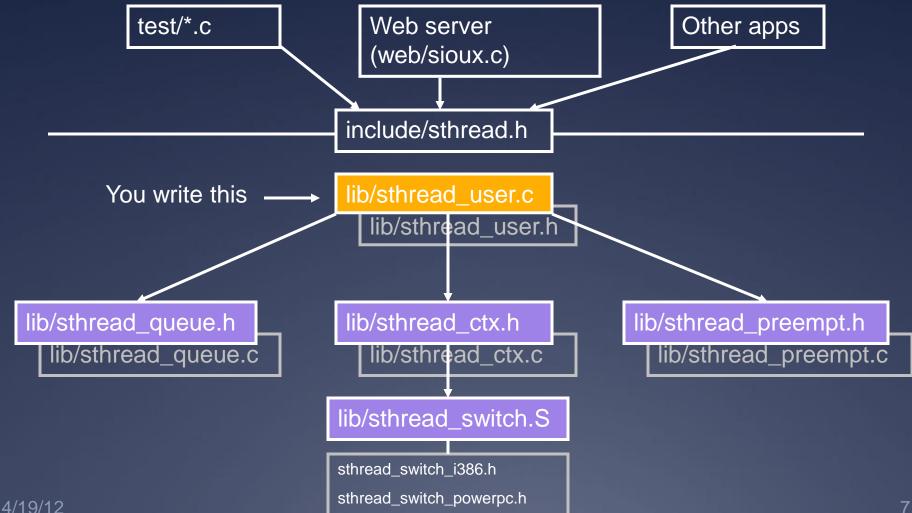
Project 2 tips

- *Understand what the provided code does for you
- * Division of work
 - * Part 3 can be completed without parts 1 and 2
- * More tools
 - * ddd
 - * (Or just gdb if you're not a fan of GUIs)

Simplethreads

- *We give you:
 - * Skeleton functions for thread interface
 - * Machine-specific code (x86, i386, PowerPC)
 - * Support for creating new stacks
 - * Support for saving regs/switching stacks
 - * A queue data structure (why?)
 - * Very simple test programs
 - * You should write more, and include them in the turnin
 - * A single-threaded web server

Simplethreads code structure



Pthreads

- *Pthreads (POSIX threads) is a preemptive, kernel-level thread library
- *Simplethreads is similar to Pthreads
- Project 2: compare your implementation against Pthreads
 - * ./configure --with-pthreads

Thread operations

*What functions do we need for a userspace thread library?

Simplethreads API

```
void sthread init()
    * Initialize the whole system
sthread t sthread create (func start func,
   void *arg)
    * Create a new thread and make it runnable
void sthread yield()
    * Give up the CPU
void sthread exit(void *ret)
    * Exit current thread
void* sthread join(sthread t t)
    * Wait for specified thread to exit
```

Simplethreads internals

*Structure of the TCB:

```
struct _sthread {
   sthread_ctx_t *saved_ctx;
   /**
     * Add your fields to the thread
     * data structure here.
     */
};
```

Sample multithreaded program

* (this slide and next – see test-create.c)

Sample multithreaded program

```
int main(int argc, char *argv[]) {
  sthread init();
   for (i = 0; i < 3; i++) {
      if (sthread create(thread start,
                           (\text{void }^*) \& i) == \text{NULL})  {
         printf("sthread create failed\n");
         exit(1);
   // needs to be called multiple times
   sthread yield();
   printf("back in main\n");
  return 0;
```

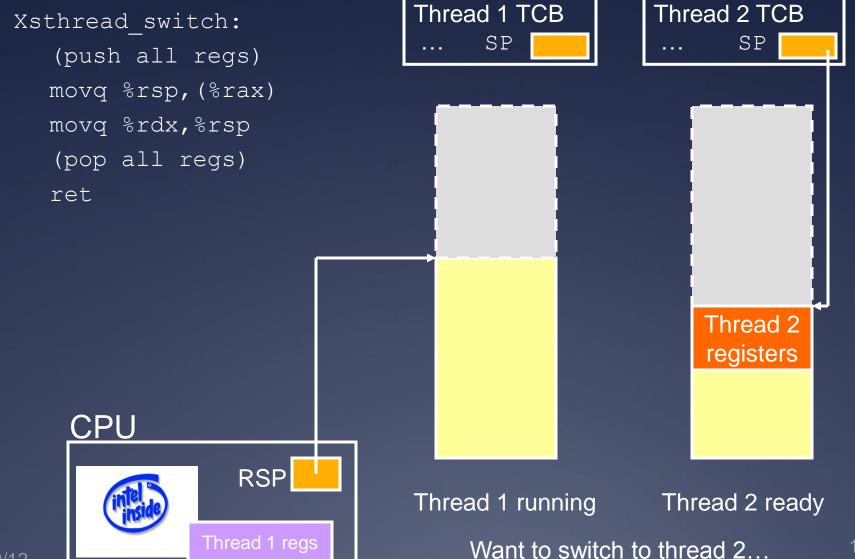
Managing contexts

- *(Provided for you in project 2)
- *Thread context = thread stack + stack pointer

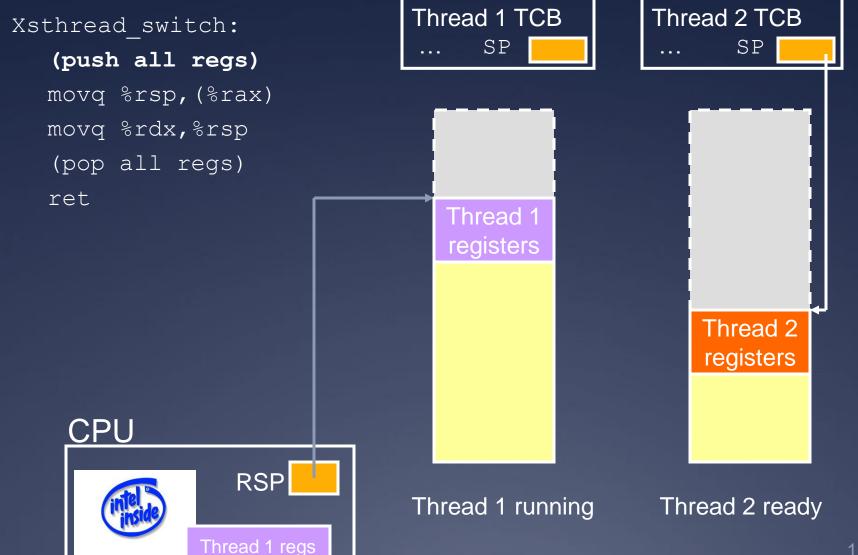
```
sthread_new_ctx(func_to_run)
```

- * creates a new thread context that can be switched to sthread_free_ctx(some_old_ctx)
- * Deletes the supplied context
 sthread_switch(oldctx, newctx)
 - * Puts current context into oldctx
 - * Takes newctx and makes it current

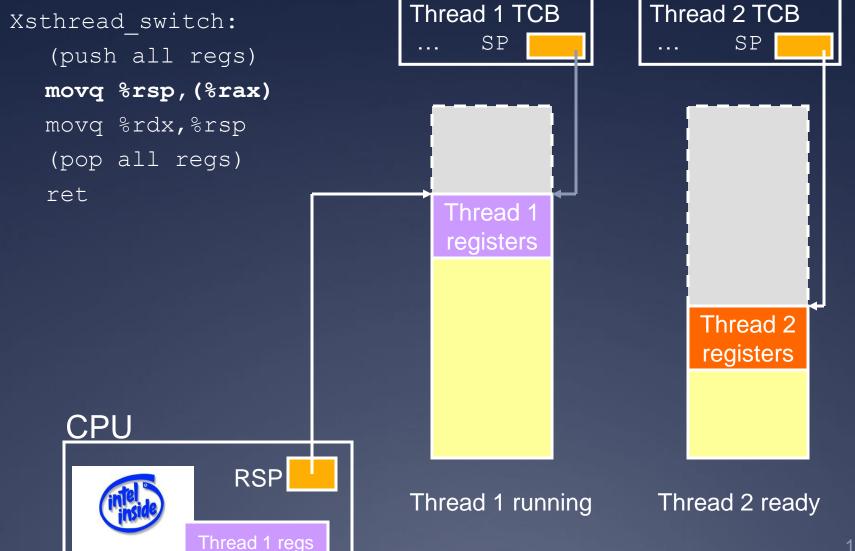
How sthread_switch works



Push old context

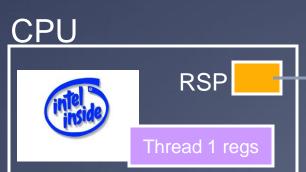


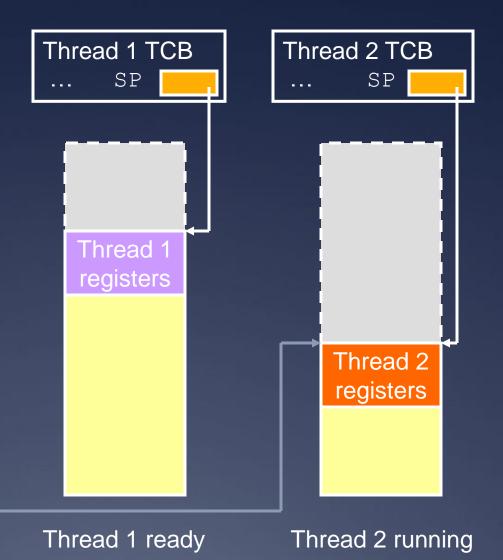
Save old stack pointer



Change stack pointers

```
Xsthread_switch:
   (push all regs)
   movq %rsp,(%rax)
   movq %rdx,%rsp
   (pop all regs)
   ret
```



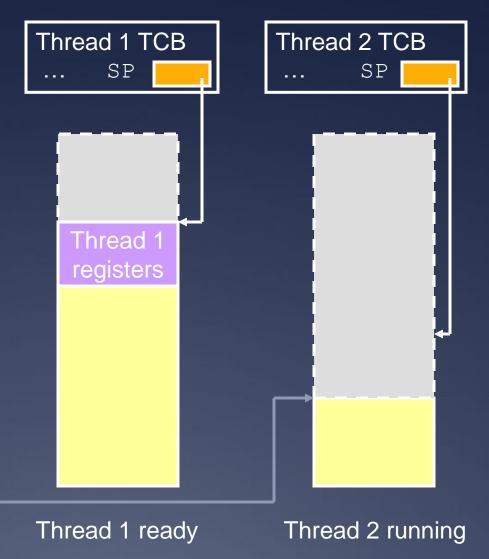


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Pop off new context

```
Xsthread_switch:
   (push all regs)
   movq %rsp, (%rax)
   movq %rdx, %rsp
   (pop all regs)
   ret
```



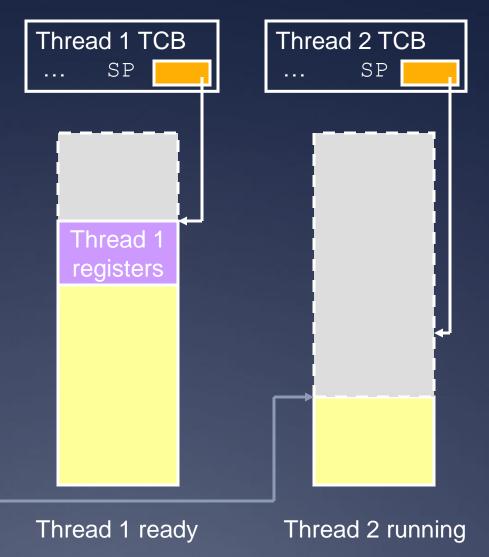


Done; return

```
Xsthread_switch:
   (push all regs)
   movq %rsp,(%rax)
   movq %rdx,%rsp
   (pop all regs)
   ret
```

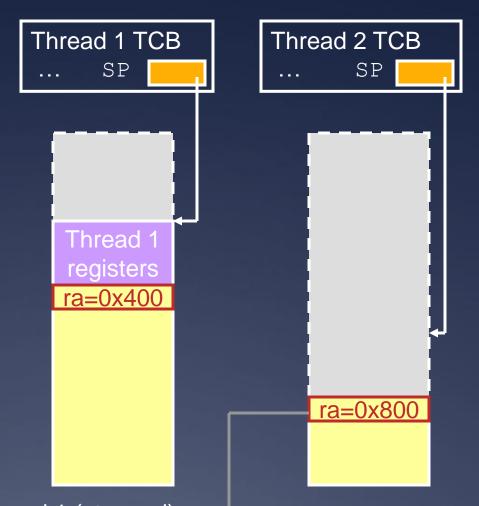
- What got switched?
 - RSP
 - PC (how?)
 - Other registers





Adjusting the PC

ret pops off the new return address!



CPU

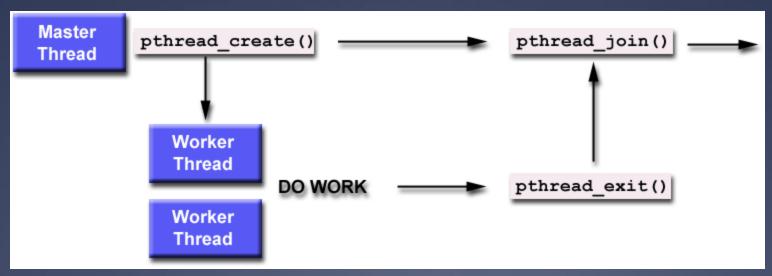


Thread 1 (stopped): sthread_switch(t1,t2); 0x400: printf("test 1");

Thread 2 (running): sthread_switch(t2,...); 0x800: printf("test 2");

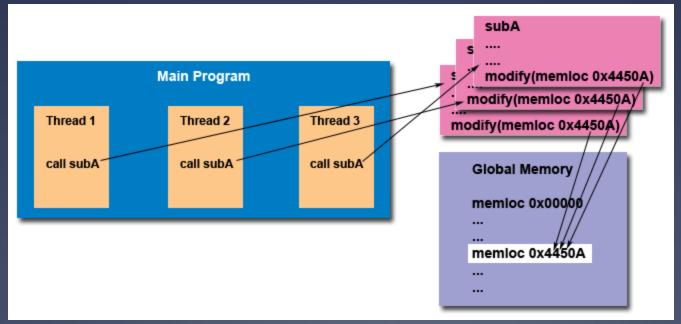
Thread joining

- *With Pthreads (and Sthreads):
 - * Master thread calls join on worker thread
 - * Join blocks until worker thread exits.
 - * Join returns the return value of the worker thread.



The need for synchronization

- *Thread safety:
 - * An application's ability to execute multiple threads simultaneously without "clobbering" shared data or creating "race" conditions



Synchronization primitives: mutexes

```
sthread_mutex_t sthread_mutex_init()
void sthread_mutex_free(sthread_mutex_t lock)

void sthread_mutex_lock(sthread_mutex_t lock)

* When returns, thread is guaranteed to acquire lock
void sthread_mutex_unlock(
    sthread_mutex_t lock)
```

Synchronization primitives: condition variables

```
sthread cond t sthread cond init()
void sthread cond free (sthread cond t cond)
void sthread cond signal (sthread cond t cond)
    * Wake-up one waiting thread, if any
void sthread cond broadcast(
   sthread cond t cond)
    * Wake-up all waiting threads, if any
void sthread cond wait (sthread cond t cond,
   sthread mutex t lock)
    * Wait for given condition variable
```

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Returning thread is guaranteed to hold the lock

Things to think about

- *How do you create a thread?
 - * How do you pass arguments to the thread's start function?
 - * Function pointer passed to sthread_new_ctx() doesn't take any arguments
- *How do you deal with the initial (main) thread?
- *How do you block a thread?

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Things to think about

- *When and how do you reclaim resources for a terminated thread?
 - * Can a thread free its stack itself?
- *Where does sthread_switch return?
- *Who and when should call sthread_switch?
- *What should be in struct _sthread_mutex, struct _sthread_cond?

Things to think about

- *Working with synchronization: When does it make sense to disable interrupts?
 - * Which actions are atomic at the application level versus at the thread level?
- *When using forkbomb, run "ulimit -Su 64" to limit the number of processes/threads
 - * Allows you to log in from another session even if you hit the above limit
 - * Add it to your .bash_profile so it happens automatically

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