Advanced Memory Management Introduction to File Systems

Last Time

- Cache Replacement Policies
 FIFO, MIN, LRU, LFU, Clock
- Memory-mapped files
- Demand-paged virtual memory

Main Points

- Applications of memory management
 - What can we do with ability to trap on memory references to individual pages?
- File systems and persistent storage
 - Goals
 - Abstractions
 - Interfaces

Address Translation Uses

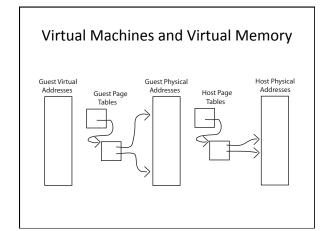
- · Process isolation
 - Keep a process from touching anyone else's memory, or the kernel's
- $\bullet \ \ \textit{Efficient interprocess communication}$
 - Shared regions of memory between processes
- Shared code segments
 - E.g., common libraries used by many different programs
- Program initialization
 - Start running a program before it is entirely in memory
- Dynamic memory allocation
 - Allocate and initialize stack/heap pages on demand

Address Translation (more)

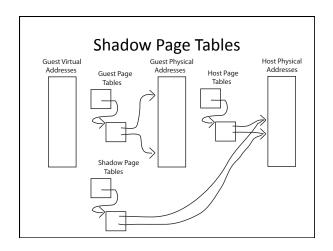
- Program debugging
 - Data breakpoints when address is accessed
- Zero-copy I/O
 - Directly from I/O device into/out of user memory
- Memory mapped files
 - Access file data using load/store instructions
- Demand-paged virtual memory
 - Illusion of near-infinite memory, backed by disk or memory on other machines

Address Translation (even more)

- Checkpoint/restart
 - Transparently save a copy of a process, without stopping the program while the save happens
- · Persistent data structures
 - Implement data structures that can survive system reboots
- · Process migration
 - Transparently move processes between machines
- Information flow control
 - Track what data is being shared externally
- · Distributed shared memory
 - Illusion of memory that is shared between machines

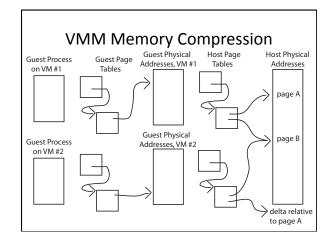


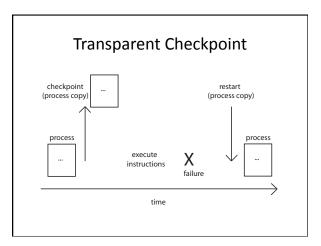
S	Segment Table				age	Table A	Page Table B	
0	Page	Tab	le A	0		0002	0	0001
1	Page Table B			1		0006	1	0004
X	(rest	invalid)		2		0000	2	0003
				3		0005	X	(rest invalid)
				x	(re	est invalid)		
	Segment			nt Tabl	e	Pag	e Table K	
		0		e Table		0	BEEF	
		X	(res	t invalid	d)	1	F000	
						2	CAFE	
						3	3333	
						4	(invalid)	
						5	BA11	
						6	DEAD	
						7	5555	
						х (rest invali	d)



Hardware Support for Virtual Machines

- x86 recently added hardware support for running virtual machines at user level
- Operating system kernel initializes two sets of translation tables
 - One for the guest OS
 - One for the host OS
- Hardware translates address in two steps
 - First using guest OS tables, then host OS tables
 - TLB holds composition





Question

- At what point can we resume the execution of a checkpointed program?
 - When the checkpoint starts?
 - When the checkpoint is entirely on disk?

Incremental Checkpoint Memory Checkpoint Incremental Incremental Memory Checkpoint Checkpoint Checkpoint В C R R D Q Q S

Deterministic Debugging

- Can we precisely replay the execution of a multi-threaded process?
 - If process does not have a memory race
- From a checkpoint, record:
 - All inputs and return values from system calls
 - All scheduling decisions
 - All synchronization operations
 - Ex: which thread acquired lock in which order

Process Migration

- What if we checkpoint a process and then restart it on a different machine?
 - Process migration: move a process from one machine to another
 - Special handling needed if any system calls are in progress
 - Where does the system call return to?

Cooperative Caching

- Can we demand page to memory on a different machine?
 - Remote memory over LAN much faster than disk
 - On page fault, look in remote memory first before fetching from disk

Distributed Virtual Memory

- Can we make a network of computers appear to be a shared-memory multiprocessor?
 - Read-write: if page is cached only on one machine
 - Read-only: if page is cached on several machines
 - Invalid: if page is cached read-write on a different machine
- On read page fault:
 - Change remote copy to read-only
- Copy remote version to local machine
- On write page fault (if cached):
 - Change remote copy to invalid
 - Change local copy to read-write

Recoverable Virtual Memory

- Data structures that survive failures
 - Want a consistent version of the data structure
 - User marks region of code as needing to be atomic
 - Begin transaction, end transaction
 - If crash, restore state before or after transaction

Recoverable Virtual Memory

- On begin transaction:
 - Snapshot data structure to disk
 - Change page table permission to read-only
- On page fault:
 - Mark page as modified by transaction
- Change page table permission to read-write
- On end transaction:
 - Log changed pages to disk
 - Commit transaction when all mods are on disk
- Recovery
 - Read last snapshot + logged changes, if committed

File Systems

- Abstraction on top of persistent storage
 - Magnetic disk
 - Flash memory (e.g., USB thumb drive)
- Devices provide
 - Storage that (usually) survives across machine crashes
 - Block level (random) access
 - Large capacity at low cost
 - Relatively slow performance
 - Magnetic disk read takes 10-20M processor instructions

File System as Illusionist: Hide Limitations of Physical Storage

- Persistence of data stored in file system:
 - Even if crash happens during an update
 - Even if disk block becomes corrupted
 - Even if flash memory wears out
- Naming:
 - Named data instead of disk block numbers
 - Directories instead of flat storage
 - Byte addressable data even though devices are blockoriented
- · Performance:
 - Cached data
 - Data placement and data structure organization
- · Controlled access to shared data

File System Abstraction

- File system: persistent, named data
- File: named collection of data in file system
 - UNIX: linear sequence of bytes
 - Windows/MacOS: collection of linear sequences
 - Example: a Word file might contain text, pictures, spreadsheets, formatting templates, ...
- · File consists of
 - Metadata
 - Access permissions, creation date
 - Data
 - File contents

File System Abstraction

- Directory
 - Group of named files or subdirectories
 - Mapping from file name to file metadata location
- Path
 - String that uniquely identifies file or directory
 - Ex: /cse/www/education/courses/cse451/12au
- Links
 - Hard link: link from name to metadata location
 - Soft link: link from name to alternate name
- Mount
 - Mapping from name in one file system to root of another

UNIX File System API

- create, link, unlink, createdir, rmdir
 - Create file, link to file, remove link
 - Create directory, remove directory
- open, close, read, write, seek
 - Open/close a file for reading/writing
 - Seek resets current position
- fsync
 - File modifications can be cached
 - fsync forces modifications to disk (like a memory barrier)