CPU Scheduling

Main Points

- Scheduling policy: what to do next, when there are multiple threads ready to run
 - Or multiple packets to send, or web requests to serve, or ...
- Definitions
 - response time, throughput, predictability
- Uniprocessor policies
 - FIFO, round robin, optimal

Example

- You manage a web site, that suddenly becomes wildly popular. Do you?
 - Buy more hardware?
 - Implement a different scheduling policy?
 - Turn away some users? Which ones?
- How much worse will performance get if the web site becomes even more popular?

Definitions

- Task/Job
- User request: e.g., mouse click, web request, shell command, ...
- · Latency/response time
- How long does a task take to complete?
- Throughput
 - How many tasks can be done per unit of time?
- Overhead
 - How much extra work is done by the scheduler?
- Fairness
 - How equal is the performance received by different users?
- Predictability
 - How consistent is the performance over time?

More Definitions

- Workload
 - Set of tasks for system to perform
- Preemptive scheduler
 - If we can take resources away from a running task
- Work-conserving
 - Resource is used whenever there is a task to run
 - For non-preemptive schedulers, work-conserving is not always
- · Scheduling algorithm
 - takes a workload as input
 - decides which tasks to do first

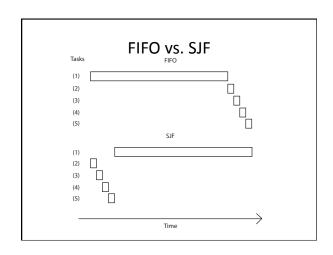
 - Performance metric (throughput, latency) as output
 Only preemptive, work-conserving schedulers to be considered

First In First Out (FIFO)

- · Schedule tasks in the order they arrive
 - Continue running them until they complete or give up the processor
- · Example: memcached
 - Facebook cache of friend lists, ...
- On what workloads is FIFO particularly bad?

Shortest Job First (SJF)

- Always do the task that has the shortest remaining amount of work to do
 - Often called Shortest Remaining Time First (SRTF)
- Suppose we have five tasks arrive one right after each other, but the first one is much longer than the others
 - Which completes first in FIFO? Next?
 - Which completes first in SJF? Next?



Shortest Job First

- Claim: SJF is optimal for average response time

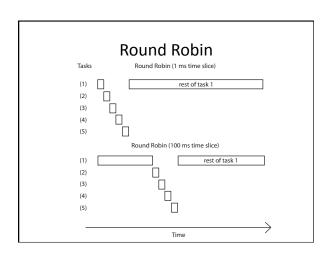
 Why?
- For what workloads is FIFO optimal?
- Pessimal?
- Does SJF have any downsides?

Starvation and Sample Bias

- Suppose you want to compare FIFO and SJF on some sequence of arriving tasks
 - Compute average response time as the average for tasks that start/end in some window
- Is this valid or invalid?

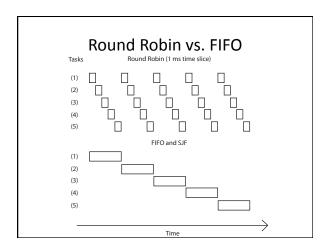
Round Robin

- Each task gets resource for a fixed period of time (time quantum)
 - If task doesn't complete, it goes back in line
- Need to pick a time quantum
 - What if time quantum is too long?
 - Infinite?
 - What if time quantum is too short?
 - One instruction?



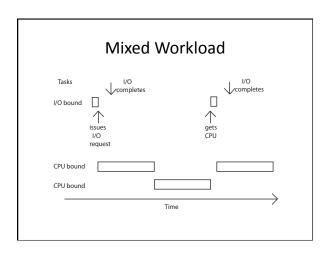
Round Robin vs. FIFO

• Assuming zero-cost time slice, is Round Robin always better than FIFO?



Round Robin vs. Fairness

• Is Round Robin always fair?



Max-Min Fairness

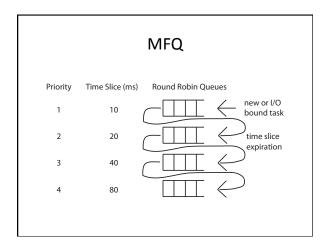
- How do we balance a mixture of repeating tasks:
 - Some I/O bound, need only a little CPU
 - Some compute bound, can use as much CPU as they are assigned
- One approach: maximize the minimum allocation given to a task
 - Schedule the smallest task first, then split the remaining time using max-min

Multi-level Feedback Queue (MFQ)

- · Goals:
 - Responsiveness
 - Low overhead
 - Starvation freedom
 - Some tasks are high/low priority
 - Fairness (among equal priority tasks)
- Not perfect at any of them!
 - Used in Linux (and probably Windows, MacOS)

MFQ

- Set of Round Robin queues
 - Each queue has a separate priority
- High priority queues have short time slices
 - Low priority queues have long time slices
- Scheduler picks first thread in highest priority queue
- Tasks start in highest priority queue
 - If time slice expires, task drops one level



Uniprocessor Summary

- FIFO is simple and minimizes overhead.
- If tasks are variable in size, then FIFO can have very poor average response time.
- If tasks are equal in size, FIFO is optimal in terms of average response time.
- Considering only the processor, SJF is optimal in terms of average response time.
- SJF is pessimal in terms of variance in response time.

Uniprocessor Summary

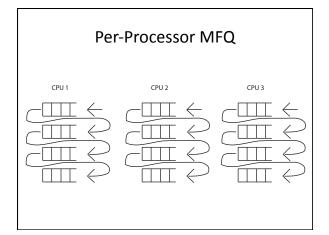
- If tasks are variable in size, Round Robin approximates SJF.
- If tasks are equal in size, Round Robin will have very poor average response time.
- Tasks that intermix processor and I/O benefit from SJF and can do poorly under Round Robin.
- Max-min fairness can improve response time for I/Obound tasks.
- Round Robin and Max-min fairness both avoid starvation.
- By manipulating the assignment of tasks to priority queues, an MFQ scheduler can achieve a balance between responsiveness, low overhead, and fairness.

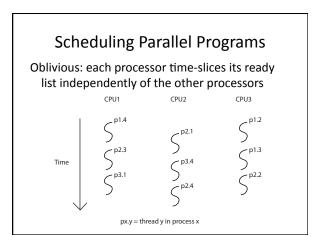
Multiprocessor Scheduling

- What would happen if we used MFQ on a multiprocessor?
 - Contention for scheduler spinlock
 - Programs will have more threads to take advantage of multiprocessor, so more contention
- · Amdahl's Law
 - Speedup on a multiprocessor limited by whatever runs sequentially
 - Runtime >= Sequential portion + parallel/# procs

Multiprocessor Scheduling

- Modern processor is 100x slower without a cache
- Cache effects of a single ready list:
 - Cache coherence overhead
 - MFQ data structure would ping between caches
 - Fetching data from other caches can be even slower than re-fetching from DRAM
 - Cache reuse
 - Thread's data from last time it ran is often still in its old cache





Scheduling Parallel Programs

 What happens if one thread gets time-sliced while other threads from the same program are still running?

