CSE 451: Operating Systems Spring 2006

Module 13 Secondary Storage

John Zahorjan zahorjan@cs.washington.edu Allen Center 534

Secondary storage

- · Secondary storage typically:
 - is anything that is outside of "primary memory"
 - does not permit direct execution of instructions or data retrieval via machine load/store instructions
- · Characteristics:
 - it's large: 50-1000GB
 - it's cheap: \$0.45/GB
 - it's persistent: data survives power loss
 - it's slow: milliseconds to access
 - · why is this slow??
 - it does fail, if rarely

5/15/2006 © 2006 Gribble, Lazowska, Levy, Zahorjan

Another trip down memory lane ... IBM 2314 About the size of 6 refrigerators 8 x 29MB (M!)

Disk trends

2

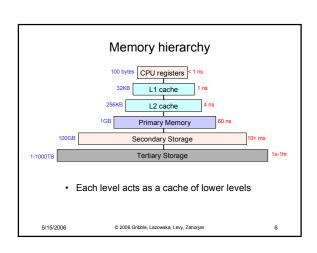
- · Disk capacity, 1975-1989
 - doubled every 3+ years
 - 25% improvement each year
 - factor of 10 every decade
 - Still exponential, but far less rapid than processor performance
- · Disk capacity since 1990
 - doubling every 12 months
 - 100% improvement each year
 - factor of 1000 every decade
 - 10x as fast as processor performance!

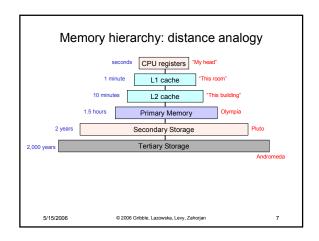
5/15/2006 © 2006 Gribble, Lazowska, Levy 4

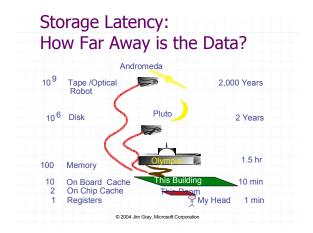
- Only a few years ago, we purchased disks by the megabyte (and it hurt!)

 The second seco
- Today, 1 GB (a billion bytes) costs \$1 \$0.50 from Dell (except you have to buy in increments of 40 80 GB)
- => 1 TB costs \$1K \$500, 1 PB costs \$1M \$500K
- In 3.2 years, 1 GB will cost \$.10
 => 1 TB for \$100, 1 PB for \$100K

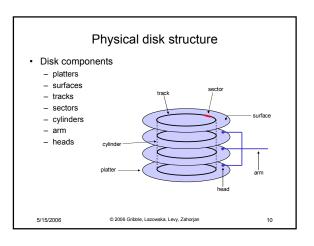
5/15/2006 © 2006 Gribble, Lazowska, Levy







Disks and the OS Disks are messy, messy devices errors, bad blocks, missed seeks, etc. Job of OS is to hide this mess from higher-level software low-level device drivers (initiate a disk read, etc.) higher-level abstractions (files, databases, etc.) OS may provide different levels of disk access to different clients physical disk block (surface, cylinder, sector) disk logical block (disk block #) file logical (filename, block or record or byte #)



Disk performance Performance depends on a number of steps seek: moving the disk arm to the correct cylinder · depends on how fast disk arm can move seek times aren't diminishing very quickly (why?) - rotation (latency): waiting for the sector to rotate under head · depends on rotation rate of disk rates are increasing, but slowly (why?) transfer: transferring data from surface into disk controller, and from there sending it back to host · depends on density of bytes on disk increasing, and very quickly · When the OS uses the disk, it tries to minimize the cost of all of these steps - particularly seeks and rotation © 2006 Gribble, Lazowska, Levy, Zahorjan 5/15/2006

Disk scheduling Seeks are very expensive, so the OS attempts to schedule disk requests that are queued waiting for the disk - FCFS (do nothing) · reasonable when load is low long waiting time for long request queues - SSTF (shortest seek time first) minimize arm movement (seek time), maximize request rate · unfairly favors middle blocks SCAN (elevator algorithm) service requests in one direction until done, then reverse · skews wait times non-uniformly (why?) - C-SCAN · like scan, but only go in one direction (typewriter) · uniform wait time: © 2006 Gribble, Lazowska, Levy, Zahorjan 5/15/2006

Interacting with disks

- In the old days...
 - OS would have to specify cylinder #, sector #, surface #, transfer size
 - · i.e., OS needs to know all of the disk parameters
- · Modern disks are even more complicated
 - $\,-\,$ not all sectors are the same size, sectors are remapped, \dots
 - disk provides a higher-level interface, e.g., SCSI
 - exports data as a logical array of blocks [0 ... N]
 - maps logical blocks to cylinder/surface/sector
 - OS only needs to name logical block #, disk maps this to cylinder/surface/sector
 - on-board cach
 - as a result, physical parameters are hidden from OS
 - both good and bad

5/15/2006

© 2006 Gribble, Lazowska, Levy, Zahorjan

13