Midterm 2 solutions

```
/* Question 1 - mutexes */
typedef struct _mutex {
  int lock; //initialized to 0
  queue t queue; //initialized to a valid empty queue
  thread towner; //initialized to NULL
} *mutex t;
void mutex lock(mutex t mutex) {
  assert(mutex);
   while(mutex->lock) {
//{\rm the} exact details here are unimportant, the idea is
//to put current thread onto mutex->queue and yield
//here is one possible solution
     thread t \text{ temp} = \text{current} thread;
     current thread = thread dequeue(ready queue);
      thread enqueue(mutex->queue, temp);
      thread yield();
   }
  mutex->lock = 1;
  mutex->owner = current thread;
}
void mutex_unlock(mutex_t mutex) {
   assert(mutex);
  if(mutex->owner!= current thread) {
      fprintf(stderr, "Error: unlocking thread doesn't own mutex");
     assert(0);
  if(! thread queue empty(mutex->queue)) {
//here, pull a waiting thread off the queue and schedule it to run
      thread t th = thread dequeue(mutex->queue);
      assert(th);
      thread enqueue(ready queue, th);
   }
  mutex->owner = NULL;
  \text{mutex-}>\text{lock}=0;
```

```
/* Question 2 - spinlocks */
   typedef\ struct\ \_spinlock\ \{
      int lock; //initialized to 0
      thread_t owner; //initialized to NULL
   } *spinlock t;
   void\ spinlock\_acquire(spinlock\_t\ s)\ \{
   //check for errors
      assert(s);
      while(test_and_set(&s->lock)); //spin until the lock is atomically taken
once free
         assert(s->owner == NULL);
         s->owner = current\_thread;
      }
   void \ spinlock\_release(spinlock\_t \ s) \ \{
   //check for errors ...
      assert(s);
      if(s->owner != current thread) {
         fprintf(stderr, "Error: unlocking thread doesn't own spinlock");
         assert(0);
      }
      s->owner = NULL;
      s->lock = 0; // ... and release the lock
```

Question 3.

- 1. Give two compelling reasons why a program might use threads. Two primary reasons are multiprocessors and I/O concurrency. A secondary reason is using threads as a program structuring tool.
- 2. Consider the following three program fragments found within the same multithreaded program.

```
fragment 1:
  P(s1);
  P(s2);
  do stuff();
  V(s2);
  V(s1);
fragment 2:
  P(s2);
                    SWAP THE ORDER OF P(s1) AND P(s2)
  P(s1);
  do stuff();
  V(s1);
  V(s2);
fragment 3:
  P(s1);
  P(s2);
  do stuff();
  V(s2);
  V(s1);
```

When this program runs, it deadlocks. While still providing for mutual exclusion using the two semaphores s1 and s2, eliminate the deadlock condition by making the fewest possible changes to the code above.

- 3. In a priority-based scheduling system, explain why we assign $\rm I/O$ bound tasks higher priority than CPU bound tasks?
- To increase CPU and device utilization and avoid the CONVOY EFFECT. The convoy effect is when all processes wait for one big process to get off the CPU.
- 4. Give two important reasons why programs can't simply disable interrupts to achieve mutual exclusion.

Multiprocessors

Can't disable interrupts and run a user program

5. Show the proper way to wait on a condition variable having Hoare semantics?

if (condition) wait

6. From an implementation perspective, what is the principle advantage of Mesa monitor semantics?

They are less strict, programmer can be lazier, sloppier.

INCORRECT ANSWER: EASIER TO IMPLEMENT

7. Monitors are often preferable to semaphores for two reasons. What are the reasons?

Separation of scheduling and mutual exclusion

Language level support