Principles of Relational Design

Chapter 12

Things can go wrong!

- All the relational designs we've seen so far have been pretty good
  - fairly small, intuitive examples
  - E/R model not far away
- A carelessly designed schema can lead to big problems
- How do we evaluate a schema?
- How do we design a good one?

Plan of Attack

- Study some informal principles (12.1)
- Functional dependencies: an important type of semantic constraint
  - define and illustrate (12.2)
  - use to define "normal forms" 2NF, 1NF, and BCNF (12.3-12.5)
- Decomposition algorithms (13.1)
- Multivalued dependencies and 4NF (13.2)
- Join dependencies and 5NF (13.3)

Informal Guidelines by Elmasri and Navathe

¶ ***Design a relation so that it is easy to explain its meaning.***

¶ ***Design so that no insertion, deletion, or modification anomalies can occur.***

¶ ***Avoid attributes whose values can be null.***

¶ ***Design so that reasonable joins do not produce spurious tuples.***

The "Universal Relation" Approach

- Assume we have identified all the individual pieces of data (attributes) of the problem.

The database design problem: group the attributes into relations.

- The informal guidelines are one way of evaluating the result.
- The theory of functional dependencies and normal forms gives a more precise way.

Functional dependency defined

- Let X and Y be attributes
- **X→ Y means that Y is a function of X**. I.e., if you know the value of X, there's only one possible value of Y. We say that "Y is functionally dependent on X" or "X determines Y."
- Note: **X→ Y does not imply Y→ X !**

Examples

- If you know the SSN, there's only one possible name (is the reverse true?)

SSN→NAME

- If you know the department number, you know the department name

DNO → DNAME

Dependencies between sets of attributes

- Given today's date and the date of birth, the age and the years until 65 are determined.

{TD, DB} → {AGE, YTO65}

- If you know the file pathname, you can determine its size, owner, and date of last modification.

FP →{SZ, O, MDT}

- Normally, when writing X → Y, we assume that X and Y are sets of attributes.

Facts about FD

- FDs are purely semantic in nature
- FDs are facts about the abstract relation, not just about a particular relation instance
  - They must hold for *all possible legal instances* of a relation
- All attributes of a relation are functionally dependent upon its key!
  - In fact, we can formally define keys in terms of FDs.

Inference Rules for FDs

- Given a set of FDs, it may be possible to deduce others by purely syntactic means.
- Example: Given that {A}→{B,C}, it follows that {A}→{B} also (and that {A} →{C})

Armstrong's and other rules

- Armstrong's rules (provable from 1st principles):
  - IR1. Reflexive rule
  - IR2. Augmentation rule
  - IR3. Transitive rule
- Some other rules (provable from IR1-2):
  - IR4. Decomposition
  - IR5. Union (additive) rule
  - IR6. Pseudotransitivity

Closures

- X+ is the "closure" of X: the set of all attributes functionally determined by X (given a set of FDs)