

Principles of Relational Design

Chapter 12

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Things can go wrong!

- All the relational designs we've seen so far have been pretty good
 - fairly small, intuitive examples
 - E/R model not far away
- A carelessly designed schema can lead to big problems
- How do we evaluate a schema?
- How do we design a good one?

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Plan of Attack

- Study some informal principles (12.1)
- Functional dependencies: an important type of semantic constraint
 - define and illustrate (12.2)
 - use to define "normal forms" 2NF, 1NF, and BCNF (12.3-12.5)
- Decomposition algorithms (13.1)
- Multivalued dependencies and 4NF (13.2)
- Join dependencies and 5NF (13.3)

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Informal Guidelines by Elmasri and Navathe

- ¶ *Design a relation so that it is easy to explain its meaning.*
- ¶ *Design so that no insertion, deletion, or modification anomalies can occur.*
- ¶ *Avoid attributes whose values can be null.*
- ¶ *Design so that reasonable joins do not produce spurious tuples.*

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The "Universal Relation" Approach

- Assume we have identified all the individual pieces of data (attributes) of the problem.
The database design problem: group the attributes into relations.
- The informal guidelines are one way of evaluating the result.
- The theory of functional dependencies and normal forms gives a more precise way.

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Functional dependency defined

- Let X and Y be attributes
- **$X \rightarrow Y$ means that Y is a function of X.**
I.e., if you know the value of X, there's only one possible value of Y. We say that "Y is functionally dependent on X" or "X determines Y."
- Note: **$X \rightarrow Y$ does not imply $Y \rightarrow X$!**

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Examples

- If you know the SSN, there's only one possible name (is the reverse true?)
SSN \rightarrow NAME
- If you know the department number, you know the department name
DNO \rightarrow DNAME

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Dependencies between sets of attributes

- Given today's date and the date of birth, the age and the years until 65 are determined.
{TD, DB} \rightarrow {AGE, YTO65}
- If you know the file pathname, you can determine its size, owner, and date of last modification.
FP \rightarrow {SZ, O, MDT}
- Normally, when writing $X \rightarrow Y$, we assume that X and Y are sets of attributes.

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Facts about FD

- FDs are purely semantic in nature
- FDs are facts about the abstract relation, not just about a particular relation instance
 - They must hold for *all possible legal instances* of a relation
- All attributes of a relation are functionally dependent upon its key!
 - In fact, we can formally define keys in terms of FDs.

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Inference Rules for FDs

- Given a set of FDs, it may be possible to deduce others by purely syntactic means.
- Example: Given that $\{A\} \rightarrow \{B,C\}$, it follows that $\{A\} \rightarrow \{B\}$ also (and that $\{A\} \rightarrow \{C\}$)

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Armstrong's and other rules

- Armstrong's rules (provable from 1st principles):
 - IR1. Reflexive rule
 - IR2. Augmentation rule
 - IR3. Transitive rule
- Some other rules (provable from IR1-2):
 - IR4. Decomposition
 - IR5. Union (additive) rule
 - IR6. Pseudotransitivity

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Closures

- X^+ is the "closure" of X: the set of all attributes functionally determined by X (given a set of FDs)

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