

Database System Internals

Query Optimization (part 2)

Paul G. Allen School of Computer Science and Engineering
University of Washington, Seattle

Announcements

- Section AB online this week
- Other sections in-person

- HW 3 out now

Key Decisions

Search Space

Optimization rules

Optimization algorithm

Two Types of Plan Enumeration Algorithms

- Dynamic programming (in class)
 - Based on System R (aka Selinger) style optimizer[1979]
 - Limited to joins: *join reordering algorithm*
 - Bottom-up
- Rule-based algorithm (will not discuss)
 - Database of rules (=algebraic laws)
 - Usually: dynamic programming
 - Usually: top-down

Two Types of Optimizers

- **Rule-based (heuristic) optimizers:**
 - Apply greedily rules that always improve plan
 - Typically: push selections down
 - Very limited: no longer used today
- **Cost-based optimizers:**
 - Use a cost model to estimate the cost of each plan
 - Select the “cheapest” plan
 - We focus on cost-based optimizers

The Three Parts of an Optimizer

- Cost estimation
 - Based on cardinality estimation
- Search space
- Search algorithm

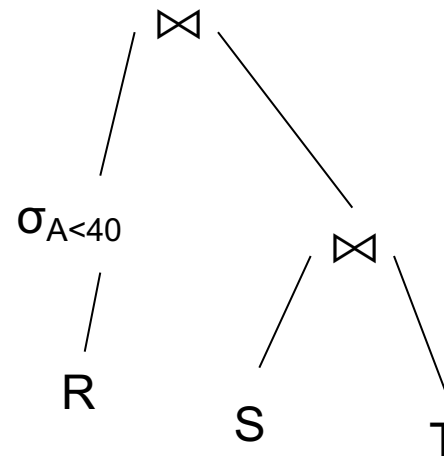
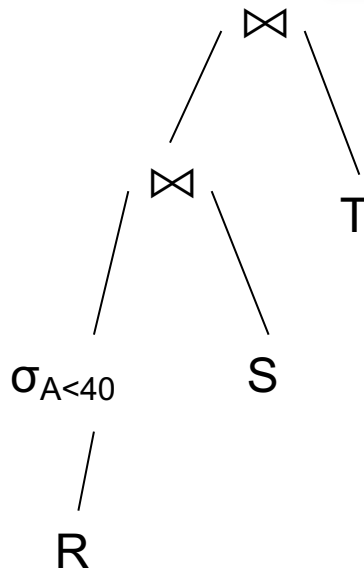
Three Approaches to Search Space

- Complete plans
- Bottom-up plans
- Top-down plans

Complete Plans

R(A,B)
S(B,C)
T(C,D)

```
SELECT *  
FROM R, S, T  
WHERE R.B=S.B and S.C=T.C and R.A<40
```



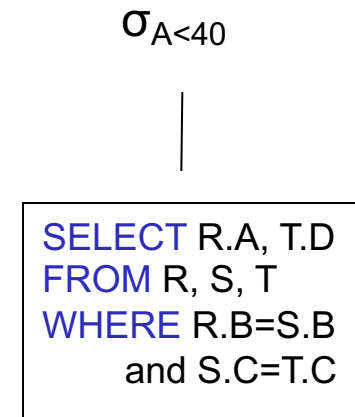
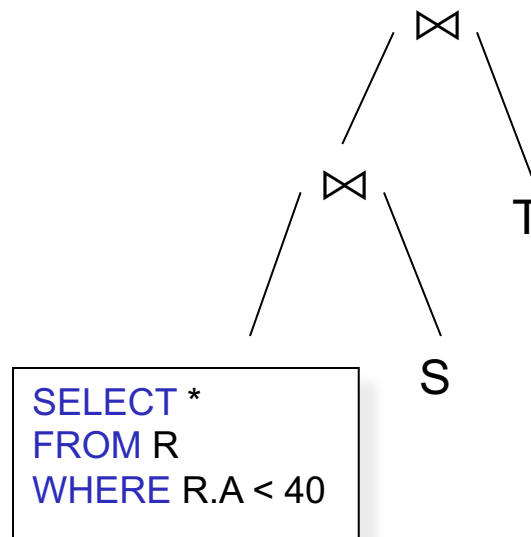
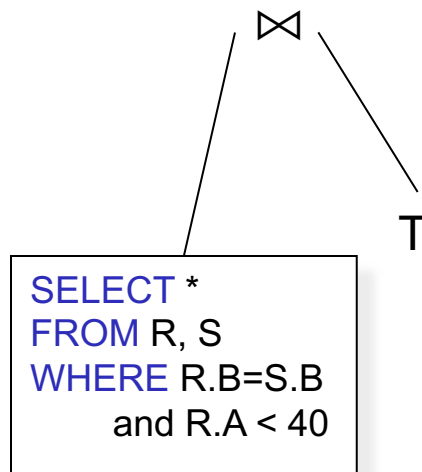
Why is this search space inefficient ?

Answer: No way to do early pruning

Top-down Partial Plans

R(A,B)
S(B,C)
T(C,D)

```
SELECT *  
FROM R, S, T  
WHERE R.B=S.B and S.C=T.C and R.A<40
```

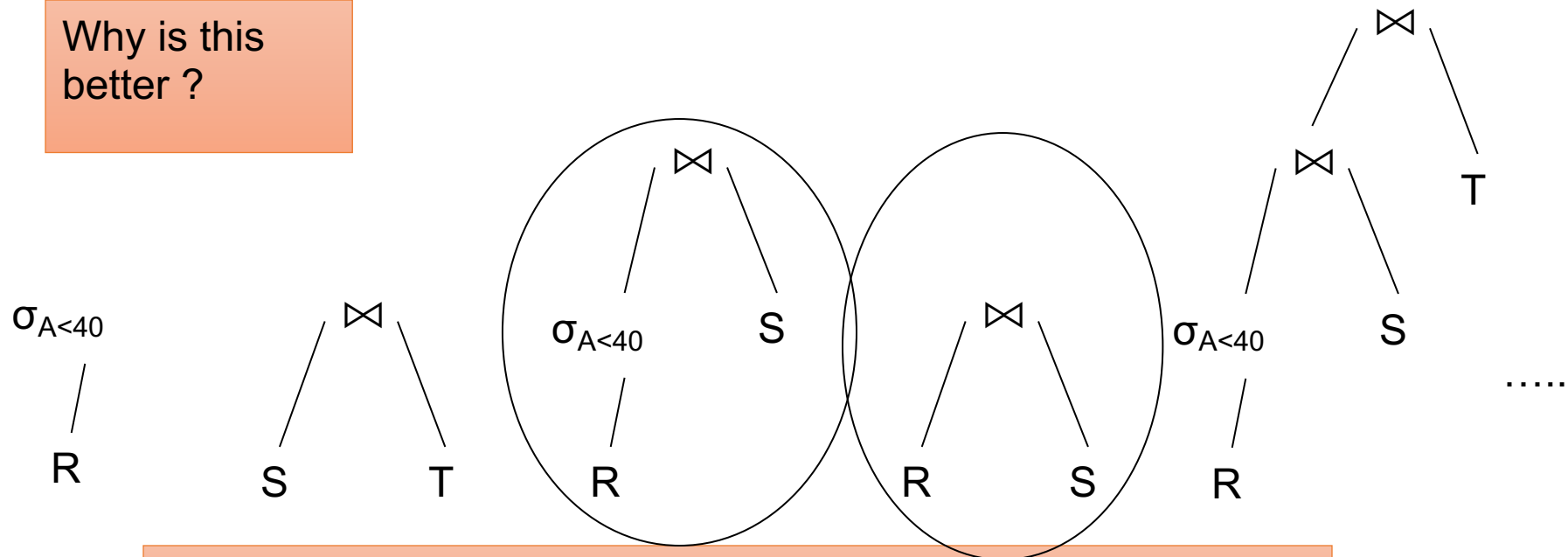


Bottom-up Partial Plans

R(A,B)
S(B,C)
T(C,D)

```
SELECT *  
FROM R, S, T  
WHERE R.B=S.B and S.C=T.C and R.A<40
```

Why is this better ?



We will prune bad plans for sub-expressions

Query Optimizer Overview

- **Input:** A logical query plan
- **Output:** A good physical query plan
- **Basic query optimization algorithm**
 - Enumerate alternative plans (logical and physical)
 - Compute estimated cost of each plan
 - Compute number of I/Os
 - Optionally take into account other resources
 - Choose plan with lowest cost
 - This is called cost-based optimization

Search Algorithm

- Dynamic programming (in class)
 - Based on System R (aka Selinger) style optimizer[1979]
 - Limited to joins: *join reordering algorithm*
 - Bottom-up
- Rule-based algorithm (will not discuss)
 - Database of rules (=algebraic laws)
 - Usually: dynamic programming
 - Usually: top-down

Dynamic Programming

Originally proposed in System R [1979]

- Only handles single block queries:

```
SELECT list  
FROM R1, ..., Rn  
WHERE cond1 AND cond2 AND . . . AND condk
```

- Some heuristics for search space enumeration:
 - Selections down
 - Projections up
 - Avoid cartesian products

Dynamic Programming

```
SELECT list  
FROM R1, ..., Rn  
WHERE cond1 AND cond2 AND . . . AND condk
```

- For each subquery $Q \subseteq \{R_1, \dots, R_n\}$ compute the following:
 - $T(Q)$ = the estimated size of Q
 - $\text{Plan}(Q)$ = a best plan for Q
 - $\text{Cost}(Q)$ = the estimated cost of that plan

Dynamic Programming

```
SELECT list  
FROM R1, ..., Rn  
WHERE cond1 AND cond2 AND . . . AND condk
```

- **Step 1:** For each $\{R_i\}$ do:
 - $T(\{R_i\}) = T(R_i)$
 - $\text{Plan}(\{R_i\}) =$ access method for R_i
 - $\text{Cost}(\{R_i\}) =$ cost of access method for R_i

- **Step 2:** For each $Q \subseteq \{R_1, \dots, R_n\}$ of size k do:
 - $T(Q)$ = use estimator
 - Consider all partitions $Q = Q' \cup Q''$
compute $\text{cost}(\text{Plan}(Q') \bowtie \text{Plan}(Q''))$
 - $\text{Cost}(Q)$ = the smallest such cost
 - $\text{Plan}(Q)$ = the corresponding plan

- **Note**
 - If we restrict to left-linear trees: Q'' = single relation
 - May want to avoid cartesian products


```
SELECT list  
FROM R1, ..., Rn  
WHERE cond1 AND cond2 AND . . . AND condk
```

- **Step 3:** Return Plan($\{R_1, \dots, R_n\}$)

Example

```
SELECT *  
FROM R, S, T, U  
WHERE cond1 AND cond2 AND . . .
```

- $R \bowtie S \bowtie T \bowtie U$
- Assumptions:

```
T(R) = 2000  
T(S) = 5000  
T(T) = 3000  
T(U) = 1000
```

- Every join selectivity is 0.001

Example

$$T(R) = 2000$$

$$T(S) = 5000$$

$$T(T) = 3000$$

$$T(U) = 1000$$

Assume

$$B(..) = T(..)/10$$

Join selectivity
is 0.001

Subquery	T	Plan	Cost
R	2000		
S	5000		
T	3000		
U	1000		
RS			
RT			
RU			
ST			
SU			
TU			
RST			
RSU			
RTU			
STU			
RSTU			

Example

$T(R) = 2000$
 $T(S) = 5000$
 $T(T) = 3000$
 $T(U) = 1000$

Assume
 $B(..) = T(..)/10$

Join selectivity
is 0.001

Subquery	T	Plan	Cost
R	2000		
S	5000		
T	3000		
U	1000		
RS	10000		
RT	6000		
RU	2000		
ST	15000		
SU	5000		
TU	3000		
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

Example

$T(R) = 2000$
 $T(S) = 5000$
 $T(T) = 3000$
 $T(U) = 1000$

Assume
 $B(..) = T(..)/10$

Join selectivity
is 0.001

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000		
T	3000		
U	1000		
RS	10000		
RT	6000		
RU	2000		
ST	15000		
SU	5000		
TU	3000		
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

Example

$$T(R) = 2000$$

$$T(S) = 5000$$

$$T(T) = 3000$$

$$T(U) = 1000$$

Assume

$$B(..) = T(..)/10$$

Join selectivity
is 0.001

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000		
U	1000		
RS	10000		
RT	6000		
RU	2000		
ST	15000		
SU	5000		
TU	3000		
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

Example

$T(R) = 2000$
 $T(S) = 5000$
 $T(T) = 3000$
 $T(U) = 1000$

Assume
 $B(..) = T(..)/10$

Join selectivity
 is 0.001

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000	Table scan	300
U	1000	Clustered index scan U.F	100
RS	10000		
RT	6000		
RU	2000		
ST	15000		
SU	5000		
TU	3000		
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

Example

$T(R) = 2000$
 $T(S) = 5000$
 $T(T) = 3000$
 $T(U) = 1000$

Assume
 $B(..) = T(..)/10$

Join selectivity
 is 0.001

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000	Table scan	300
U	1000	Clustered index scan U.F	100
RS	10000	$R \bowtie S$ nested loop join	...
RT	6000		
RU	2000		
ST	15000		
SU	5000		
TU	3000		
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

Example

$$T(R) = 2000$$

$$T(S) = 5000$$

$$T(T) = 3000$$

$$T(U) = 1000$$

Assume
 $B(..) = T(..)/10$

Join selectivity
 is 0.001

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000	Table scan	300
U	1000	Clustered index scan U.F	100
RS	10000	$R \bowtie S$ nested loop join	...
RT	6000	$R \bowtie T$ index join	...
RU	2000		
ST	15000		
SU	5000		
TU	3000		
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

Example

$$T(R) = 2000$$

$$T(S) = 5000$$

$$T(T) = 3000$$

$$T(U) = 1000$$

Assume

$$B(..) = T(..)/10$$

Join selectivity
is 0.001

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000	Table scan	300
U	1000	Clustered index scan U.F	100
RS	10000	R ⋈ S nested loop join	...
RT	6000	R ⋈ T index join	...
RU	2000	R ⋈ U index join	
ST	15000	S ⋈ T hash join	
SU	5000	...	
TU	3000	...	
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

Example

$$T(R) = 2000$$

$$T(S) = 5000$$

$$T(T) = 3000$$

$$T(U) = 1000$$

Assume

$$B(..) = T(..)/10$$

Join selectivity
is 0.001

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000	Table scan	300
U	1000	Clustered index scan U.F	100
RS	10000	R ⋈ S nested loop join	...
RT	6000	R ⋈ T index join	...
RU	2000	R ⋈ U index join	
ST	15000	S ⋈ T hash join	
SU	5000	...	
TU	3000	...	
RST	30000	(RT) ⋈ S hash join	...
RSU	10000	(SU) ⋈ R merge join	
RTU	6000	...	
STU	15000		
RSTU	30000		

Example

$$T(R) = 2000$$

$$T(S) = 5000$$

$$T(T) = 3000$$

$$T(U) = 1000$$

Assume

$$B(..) = T(..)/10$$

Join selectivity
is 0.001

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000	Table scan	300
U	1000	Clustered index scan U.F	100
RS	10000	R ⋈ S nested loop join	...
RT	6000	R ⋈ T index join	...
RU	2000	R ⋈ U index join	
ST	15000	S ⋈ T hash join	
SU	5000	...	
TU	3000	...	
RST	30000	(RT) ⋈ S hash join	...
RSU	10000	(SU) ⋈ R merge join	
RTU	6000	...	
STU	15000		
RSTU	30000	(RT) ⋈ (SU) hash join	...

Discussion

- For the subset $\{RS\}$, need to consider both $R \bowtie S$ and $S \bowtie R$
 - Because the cost may be different!
- When computing the cheapest plan for $(Q) \bowtie R$, we may consider new access methods for R , e.g. an index look-up that makes sense only in the context of the join

Discussion

```
SELECT list  
FROM R1, ..., Rn  
WHERE cond1 AND cond2 AND . . . AND condk
```

Given a query with n relations R_1, \dots, R_n

- How many entries do we have in the dynamic programming table?
- For each entry, how many alternative plans do we need to inspect?

```
SELECT list  
FROM R1, ..., Rn  
WHERE cond1 AND cond2 AND . . . AND condk
```

Given a query with n relations R_1, \dots, R_n

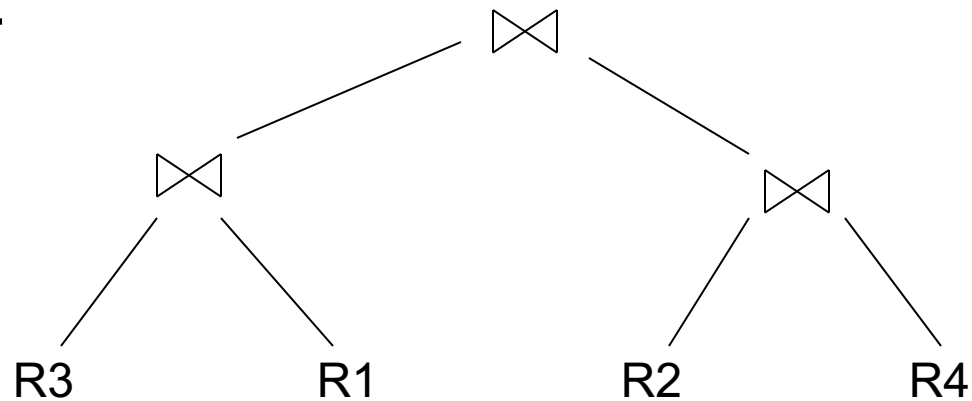
- How many entries do we have in the dynamic programming table?
 - A: $2^n - 1$
- For each entry, how many alternative plans do we need to inspect?
 - A: for each entry with k tables, examine $2^k - 2$ plans

Reducing the Search Space

- Left-linear trees
- No cartesian products

Join Trees

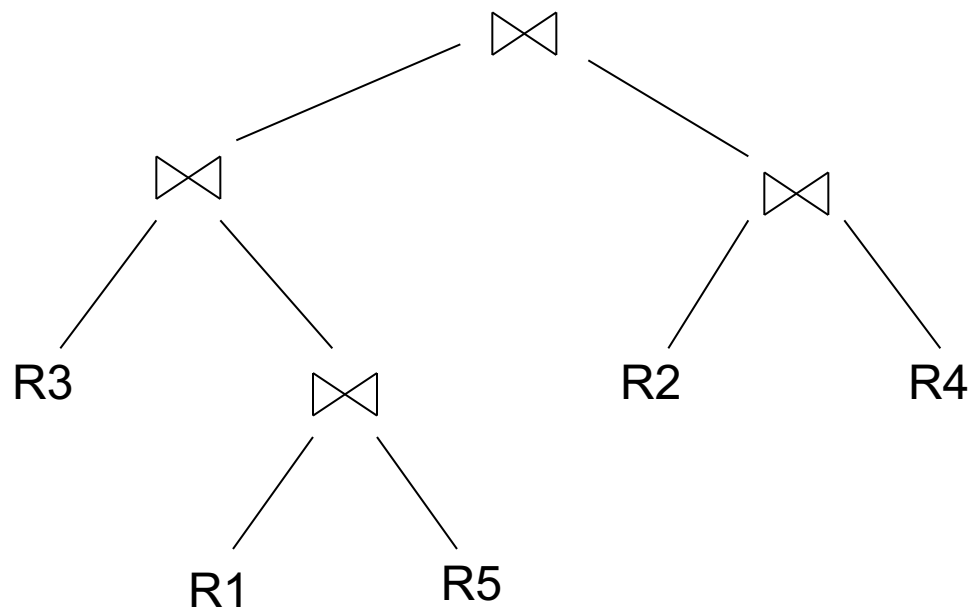
- $R_1 \bowtie R_2 \bowtie \dots \bowtie R_n$
- Join tree:



- A plan = a join tree
- A partial plan = a subtree of a join tree

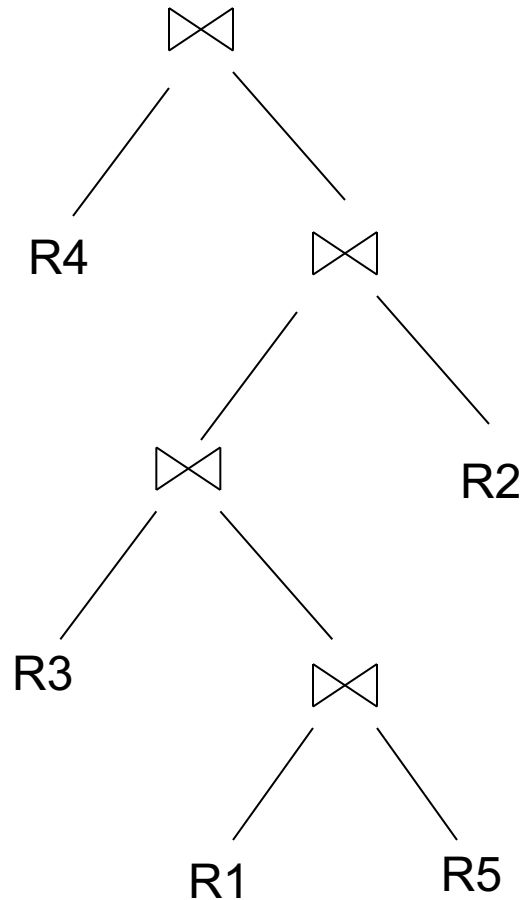
Types of Join Trees

- Bushy:



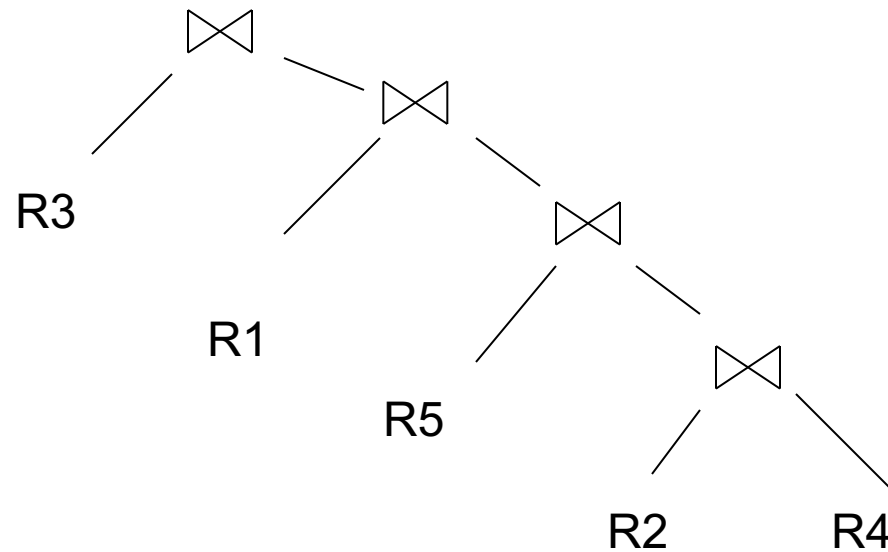
Types of Join Trees

- Linear :



Types of Join Trees

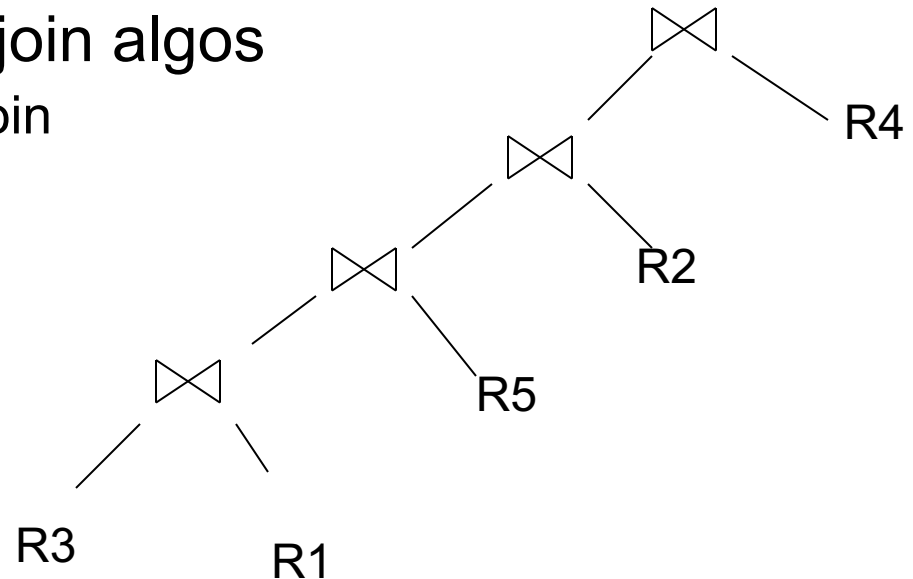
- Right deep:



Types of Join Trees

- Left deep:

- Work well with existing join algos
 - Nested-loop and hash-join
- Facilitate pipelining



- Dynamic programming can be used with all trees

No Cartesian Products

$$R(A,B) \bowtie S(B,C) \bowtie T(C,D)$$

Plan: $(R(A,B) \bowtie T(C,D)) \bowtie S(B,C)$
has a cartesian product.

Most query optimizers will not consider it