

# CSE 444: Database Internals

## Section 6: Optimistic Concurrency Control

# Today

- Timestamp-based Concurrency Control
- Multiversion Concurrency Control
- Extra help on lab 3

# Problem 1: Timestamp-based Concurrency Control

# Timestamp-based Concurrency Control

- Some transaction,  $T$ .
- Some element (tuple/page),  $X$ .
- **TS**( $T$ ) - timestamp for transaction  $T$ 
  - Stays constant for all of  $T$ 's operations
- **WT**( $X$ ) – latest write timestamp for  $X$ 
  - Set  $WT(X) = TS(T)$
- **RT**( $X$ ) – latest read timestamp for  $X$ 
  - Set  $RT(X) = TS(T)$
- **C**( $X$ ) –  $X$ 's value has been committed
  - 1 if true, 0 if not

# Timestamp-based Concurrency Control

- **Actions for transaction T**
  - **Grant** a read/write request for a transaction
  - **Abort** (in case T violates physical reality – late actions)
  - **Delay** (make the Grant or Abort decision later)
    - When writing, the change is always tentative until we decide to commit. For this, we use a commit bit C to keep track if the transaction that last wrote X has committed
  - **Ignore *Thomas Write Rule*** – ignore outdated writes

# Timestamp-based Concurrency Control - Four Rules

- **Rule 1: Read** request on  $X$  by  $T$ 
  - $TS(T) < WT(X)$ , **abort**, (read too late)
  - $TS(T) \geq WT(X)$ , physically realizable
    - If  $C = 1$ , **grant**, update  $RT(X)$
    - If  $C = 0$ , **delay**  $T$

# Timestamp-based Concurrency Control - Four Rules

- **Rule 2: Write** request on **X** by **T**
  - $TS(T) < RT(X)$  (write too late)
    - **Abort**
  - $TS(T) \geq RT(X)$ , physically realizable
    - $TS(T) \geq WT(X)$ 
      - then **grant**, update  $WT(X)$ , set  $C = 0$  (as it's not committed yet)
    - $TS(T) < WT(X)$ 
      - If  $C = 1$ , **ignore** (*Thomas Write Rule* – ignore outdated writes)
      - If  $C = 0$ , **delay**

# Timestamp-based Concurrency Control - Four Rules

- **Rule 3: Commit** request by **T**
  - Set  $C = 1$  for all **X** written by **T**
  - Allow waiting transactions to proceed
- **Rule 4: Abort** transaction **T**
  - Check if the waiting transactions can proceed now.

# Timestamp-based Concurrency Control

Two transactions get started.

- $\text{Start}(T_1) \rightarrow \text{Start}(T_2)$

# Timestamp-based Concurrency Control

What will happen at the last request?

- $\text{Start}(T_1) \rightarrow \text{Start}(T_2) \rightarrow R_{T_1}(A) \rightarrow R_{T_2}(A) \rightarrow W_{T_1}(B) \rightarrow \mathbf{W_{T_2}(B)}$

# Timestamp-based Concurrency Control

What will happen at the last request?

- $\text{Start}(T_1) \rightarrow \text{Start}(T_2) \rightarrow R_{T_1}(A) \rightarrow R_{T_2}(A) \rightarrow W_{T_1}(B) \rightarrow \mathbf{W_{T_2}(B)}$ 
  - **ACCEPTED** [no need to check C(B)]

# Timestamp-based Concurrency Control

What will happen at the last request?

- $\text{Start}(T_1) \rightarrow \text{Start}(T_2) \rightarrow R_{T_1}(A) \rightarrow R_{T_2}(A) \rightarrow W_{T_1}(B) \rightarrow \mathbf{W_{T_2}(B)}$   
– **ACCEPTED** [no need to check C(B)]
  
- $\text{Start}(T_1) \rightarrow \text{Start}(T_2) \rightarrow R_{T_2}(A) \rightarrow \text{Commit}_{T_2} \rightarrow R_{T_1}(A) \rightarrow \mathbf{W_{T_1}(A)}$

# Timestamp-based Concurrency Control

What will happen at the last request?

- $\text{Start}(T_1) \rightarrow \text{Start}(T_2) \rightarrow R_{T_1}(A) \rightarrow R_{T_2}(A) \rightarrow W_{T_1}(B) \rightarrow \mathbf{W_{T_2}(B)}$ 
  - **ACCEPTED** [no need to check C(B)]
  
- $\text{Start}(T_1) \rightarrow \text{Start}(T_2) \rightarrow R_{T_2}(A) \rightarrow \text{Commit}_{T_2} \rightarrow R_{T_1}(A) \rightarrow \mathbf{W_{T_1}(A)}$ 
  - **ABORT**  $T_1$  because  $R_{T_2}(A)$  precedes

# Problem 2: Timestamp-based Concurrency Control

- $TS_1 \rightarrow TS_2 \rightarrow TS_3 \rightarrow TS_4 \rightarrow R_1(X) \rightarrow R_2(X) \rightarrow$   
 $W_2(X) \rightarrow W_1(X) \rightarrow W_3(Y) \rightarrow W_2(Y) \rightarrow C_3 \rightarrow$   
 $W_4(Z) \rightarrow C_4 \rightarrow R_2(Z)$
- Remember!
  - Note changes to RT, WT, A and C bit for each element
  - Apply four rules

























$ST_1 \rightarrow ST_2 \rightarrow ST_3 \rightarrow ST_4 \rightarrow R_1(X) \rightarrow R_2(X) \rightarrow W_2(X) \rightarrow W_1(X) \rightarrow W_3(Y) \rightarrow W_2(Y) \rightarrow C_3 \rightarrow W_4(Z) \rightarrow C_4 \rightarrow R_2(Z)$

T1	T2	T3	T4	X	Y	Z
1	2	3	4	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1
$R_1(X)$				RT=1		
	$R_2(X)$			RT=2		
	$W_2(X)$			WT=2, C=0		
$W_1(X)$ : abort						
		$W_3(Y)$			WT=3, C=0	
	$W_2(Y)$ : delay					

1. Physically realizable:

$TS(T_3) \geq RT(X)$  although  $TS(T_2) < WT(X)$

2. We could not apply Thomas' write rule (**ignore  $W_2(Y)$** ) since  $C=0$





ST<sub>1</sub> -> ST<sub>2</sub> -> ST<sub>3</sub> -> ST<sub>4</sub> -> R<sub>1</sub>(X) -> R<sub>2</sub>(X) -> W<sub>2</sub>(X) -> W<sub>1</sub>(X) -> W<sub>3</sub>(Y) -> W<sub>2</sub>(Y) -> C<sub>3</sub> -> W<sub>4</sub>(Z) -> C<sub>4</sub> -> R<sub>2</sub>(Z)

T1	T2	T3	T4	X	Y	Z
1	2	3	4	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1
R <sub>1</sub> (X)				RT=1		
	R <sub>2</sub> (X)			RT=2		
	W <sub>2</sub> (X)			WT=2, C=0		
W <sub>1</sub> (X): <b>abort</b>						
		W <sub>3</sub> (Y)			WT=3, C=0	
	W <sub>2</sub> (Y): <b>delay</b>					
		C <sub>3</sub>			C=1	

A later write by T<sub>3</sub> has been committed!

$ST_1 \rightarrow ST_2 \rightarrow ST_3 \rightarrow ST_4 \rightarrow R_1(X) \rightarrow R_2(X) \rightarrow W_2(X) \rightarrow W_1(X) \rightarrow W_3(Y) \rightarrow W_2(Y) \rightarrow C_3 \rightarrow W_4(Z) \rightarrow C_4 \rightarrow$   
 $R_2(Z)$

T1	T2	T3	T4	X	Y	Z
1	2	3	4	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1
$R_1(X)$				RT=1		
	$R_2(X)$			RT=2		
	$W_2(X)$			WT=2, C=0		
$W_1(X)$ : abort						
		$W_3(Y)$			WT=3, C=0	
	$W_2(Y)$ : delay					
		$C_3$			C=1	
	<b>Ignore <math>W_2(Y)</math> and proceed</b>					

$ST_1 \rightarrow ST_2 \rightarrow ST_3 \rightarrow ST_4 \rightarrow R_1(X) \rightarrow R_2(X) \rightarrow W_2(X) \rightarrow W_1(X) \rightarrow W_3(Y) \rightarrow W_2(Y) \rightarrow C_3 \rightarrow W_4(Z) \rightarrow C_4 \rightarrow R_2(Z)$

T1	T2	T3	T4	X	Y	Z
1	2	3	4	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1
	Ignore $W_2(Y)$ and <b>proceed</b>					
			$W_4(Z)$			

ST<sub>1</sub> -> ST<sub>2</sub> -> ST<sub>3</sub> -> ST<sub>4</sub> -> R<sub>1</sub>(X) -> R<sub>2</sub>(X) -> W<sub>2</sub>(X) -> W<sub>1</sub>(X) -> W<sub>3</sub>(Y) -> W<sub>2</sub>(Y) -> C<sub>3</sub> -> W<sub>4</sub>(Z) -> C<sub>4</sub> ->

1. Physically realizable:

$TS(T_4) \geq RT(X)$  and  $TS(T_4) \geq WT(X)$

2. Update WT and C (not committed yet)

Y	Z
RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1
	WT=4, C = 0

ignore W<sub>2</sub>(Y)  
and proceed

W<sub>4</sub>(Z)

$ST_1 \rightarrow ST_2 \rightarrow ST_3 \rightarrow ST_4 \rightarrow R_1(X) \rightarrow R_2(X) \rightarrow W_2(X) \rightarrow W_1(X) \rightarrow W_3(Y) \rightarrow W_2(Y) \rightarrow C_3 \rightarrow W_4(Z) \rightarrow C_4 \rightarrow$   
 $R_2(Z)$

T1	T2	T3	T4	X	Y	Z
1	2	3	4	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1
	Ignore $W_2(Y)$ and <b>proceed</b>					
			$W_4(Z)$			WT=4, C = 0
			$C_4$			C=1

ST<sub>1</sub> -> ST<sub>2</sub> -> ST<sub>3</sub> -> ST<sub>4</sub> -> R<sub>1</sub>(X) -> R<sub>2</sub>(X) -> W<sub>2</sub>(X) -> W<sub>1</sub>(X) -> W<sub>3</sub>(Y) -> W<sub>2</sub>(Y) -> C<sub>3</sub> -> W<sub>4</sub>(Z) -> C<sub>4</sub> -> R<sub>2</sub>(Z)

T1	T2	T3	T4	X	Y	Z
1	2	3	4	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1
	Ignore W <sub>2</sub> (Y) and <b>proceed</b>					
			W <sub>4</sub> (Z)			WT=4, C = 0
			C <sub>4</sub>			C=1
	R <sub>2</sub> (Z)					

$ST_1 \rightarrow ST_2 \rightarrow ST_3 \rightarrow ST_4 \rightarrow R_1(X) \rightarrow R_2(X) \rightarrow W_2(X) \rightarrow W_1(X) \rightarrow W_3(Y) \rightarrow W_2(Y) \rightarrow C_3 \rightarrow W_4(Z) \rightarrow C_4$

1. **NOT** Physically realizable:  
 $TS(T_2) < WT(Z)$   
 Abort/rollback

T4	X	Y	Z
4	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1	RT = 0, WT = 0, C = 1
			WT=4, C = 0
			C=1

and proceed

$R_2(Z)$ : abort

# Timestamp-based Concurrency Control

Questions?

# Multiversion Concurrency Control

- Maintains **old** versions of database elements in addition the current version in the database itself.
- The idea is to allow reads that would otherwise result in an abort (as the current version was written by future transaction)

# Problem with Timestamp-Based Scheduling

T1	T2	T3	T4	A
150	200	175	225	RT = 0 WT = 0
$R_1(A)$				RT = 150
$W_1(A)$				WT = 150
	$R_2(A)$			RT = 200
	$W_2(A)$			WT = 200
		$R_3(A)$		
		<b>Abort</b>		
			$R_4(A)$	RT = 225

Had to abort because  
WT(A) is greater than  
my own timestamp

Would have been useful if I  
had access to an old version  
of A (from 150)...

# Multiversion Timestamps

T1	T2	T3	T4	A <sub>0</sub>	A <sub>150</sub>	A <sub>200</sub>
150	200	175	225	RT = 0 WT = 0		
R <sub>1</sub> (A)				Read		
W <sub>1</sub> (A)					Create	
	R <sub>2</sub> (A)				Read	
	W <sub>2</sub> (A)					Create
		R <sub>3</sub> (A)			Read	
			R <sub>4</sub> (A)			Read

Don't have to abort

Just read a previous value of  
A

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
			$W_4(A)$							

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
			$W_4(A)$						Create	

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
$W_1(A)$			$W_4(A)$						Create	

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
$W_1(A)$			$W_4(A)$							Create
						Create				

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
$W_1(A)$	$R_2(A)$		$W_4(A)$			Create				Create

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
$W_1(A)$	$R_2(A)$		$W_4(A)$			Create				Create
						RT=2				

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
$W_1(A)$	$R_2(A)$	$R_3(A)$	$W_4(A)$			Create RT=2				Create

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
$W_1(A)$	$R_2(A)$	$R_3(A)$	$W_4(A)$			Create RT=2 RT=3				Create

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
$W_1(A)$	$R_2(A)$ $W_2(A)$	$R_3(A)$	$W_4(A)$			Create RT=2 RT=3			Create	

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
$W_1(A)$	$R_2(A)$ $W_2(A)$ <b>abort</b>	$R_3(A)$	$W_4(A)$			Create				Create
						RT=2				
						RT=3				

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
$W_1(A)$	$R_2(A)$ $W_2(A)$ <b>abort</b>	$R_3(A)$	$W_4(A)$	$R_5(A)$		Create	RT=2	RT=3		Create

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
$W_1(A)$	$R_2(A)$ $W_2(A)$ <b>abort</b>	$R_3(A)$	$W_4(A)$	$R_5(A)$		Create RT=2 RT=3			Create RT=5	

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
$W_1(A)$	$R_2(A)$ $W_2(A)$ <b>abort</b>	$R_3(A)$	$W_4(A)$	$R_5(A)$ $W_5(A)$		Create RT=2 RT=3			Create RT=5	

# Second Example w/ Multiversion

$T_1$	$T_2$	$T_3$	$T_4$	$T_5$	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	3	4	5						
$W_1(A)$	$R_2(A)$ $W_2(A)$ <b>abort</b>	$R_3(A)$	$W_4(A)$	$R_5(A)$ $W_5(A)$		Create RT=2 RT=3			Create RT=5	Create

# Second Example w/ Multiversion

T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	A <sub>0</sub>	A <sub>1</sub>	A <sub>2</sub>	A <sub>3</sub>	A <sub>4</sub>	A <sub>5</sub>
1	2	3	4	5						
W <sub>1</sub> (A)	R <sub>2</sub> (A) W <sub>2</sub> (A) <b>abort</b>	R <sub>3</sub> (A)	W <sub>4</sub> (A)  R <sub>4</sub> (A)	R <sub>5</sub> (A) W <sub>5</sub> (A)		Create RT=2 RT=3			Create RT=5	Create

# Second Example w/ Multiversion

T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	A <sub>0</sub>	A <sub>1</sub>	A <sub>2</sub>	A <sub>3</sub>	A <sub>4</sub>	A <sub>5</sub>
1	2	3	4	5						
W <sub>1</sub> (A)	R <sub>2</sub> (A) W <sub>2</sub> (A) <b>abort</b>	R <sub>3</sub> (A)	W <sub>4</sub> (A)  R <sub>4</sub> (A)	R <sub>5</sub> (A) W <sub>5</sub> (A)		Create RT=2 RT=3			Create RT=5	Create RT=5



# Second Example w/ Multiversion

T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	A <sub>0</sub>	A <sub>1</sub>	A <sub>2</sub>	A <sub>3</sub>	A <sub>4</sub>	A <sub>5</sub>
1	2	3	4	5						
W <sub>1</sub> (A)	R <sub>2</sub> (A)	R <sub>3</sub> (A)	W <sub>4</sub> (A)			Create			Create	
	W <sub>2</sub> (A) <b>abort</b>			R <sub>5</sub> (A)		RT=2				RT=5
				W <sub>5</sub> (A)						RT=5 Create
R <sub>1</sub> (A)			R <sub>4</sub> (A)							
						RT=3				

# Second Example w/ Multiversion

T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	A <sub>0</sub>	A <sub>1</sub>	A <sub>2</sub>	A <sub>3</sub>	A <sub>4</sub>	A <sub>5</sub>
1	2	3	4	5						
W <sub>1</sub> (A)	R <sub>2</sub> (A)	R <sub>3</sub> (A)	W <sub>4</sub> (A)			Create			Create	
	W <sub>2</sub> (A) <b>abort</b>			R <sub>5</sub> (A)		RT=2			RT=5	
				W <sub>5</sub> (A)		RT=3			RT=5	Create
R <sub>1</sub> (A)			R <sub>4</sub> (A)							
C						RT=3				

# Second Example w/ Multiversion

T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	A <sub>0</sub>	A <sub>1</sub>	A <sub>2</sub>	A <sub>3</sub>	A <sub>4</sub>	A <sub>5</sub>
1	2	3	4	5						
W <sub>1</sub> (A)	R <sub>2</sub> (A)	R <sub>3</sub> (A)	W <sub>4</sub> (A)			Create			Create	
	W <sub>2</sub> (A)					RT=2				
	<b>abort</b>					RT=3				
				R <sub>5</sub> (A)					RT=5	
				W <sub>5</sub> (A)					RT=5	Create
R <sub>1</sub> (A)			R <sub>4</sub> (A)							
C					X	RT=3				

# Second Example w/ Multiversion

T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	A <sub>0</sub>	A <sub>1</sub>	A <sub>2</sub>	A <sub>3</sub>	A <sub>4</sub>	A <sub>5</sub>
1	2	3	4	5						
W <sub>1</sub> (A)	R <sub>2</sub> (A)	R <sub>3</sub> (A)	W <sub>4</sub> (A)			Create			Create	
	W <sub>2</sub> (A) <b>abort</b>					RT=2				
				R <sub>5</sub> (A)		RT=3			RT=5	
				W <sub>5</sub> (A)					RT=5	Create
R <sub>1</sub> (A)			R <sub>4</sub> (A)							
C					X	RT=3				

X means that we can delete this version

# Second Example w/ Multiversion

T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	A <sub>0</sub>	A <sub>1</sub>	A <sub>2</sub>	A <sub>3</sub>	A <sub>4</sub>	A <sub>5</sub>
1	2	3	4	5						
W <sub>1</sub> (A)	R <sub>2</sub> (A)	R <sub>3</sub> (A)	W <sub>4</sub> (A)			Create			Create	
	W <sub>2</sub> (A) <b>abort</b>			R <sub>5</sub> (A)		RT=2				RT=5
			R <sub>4</sub> (A)	W <sub>5</sub> (A)		RT=3				RT=5 Create
R <sub>1</sub> (A)										
C		C			X					

X means that we can delete this version

# Second Example w/ Multiversion

T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	A <sub>0</sub>	A <sub>1</sub>	A <sub>2</sub>	A <sub>3</sub>	A <sub>4</sub>	A <sub>5</sub>
1	2	3	4	5						
W <sub>1</sub> (A)	R <sub>2</sub> (A)	R <sub>3</sub> (A)	W <sub>4</sub> (A)			Create			Create	
	W <sub>2</sub> (A) <b>abort</b>					RT=2				
				R <sub>5</sub> (A)		RT=3			RT=5	
			R <sub>4</sub> (A)	W <sub>5</sub> (A)					RT=5	Create
R <sub>1</sub> (A)										
C		C			X	RT=3				
							X			

X means that we can delete this version

# Second Example w/ Multiversion

T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	A <sub>0</sub>	A <sub>1</sub>	A <sub>2</sub>	A <sub>3</sub>	A <sub>4</sub>	A <sub>5</sub>
1	2	3	4	5						
W <sub>1</sub> (A)	R <sub>2</sub> (A)	R <sub>3</sub> (A)	W <sub>4</sub> (A)			Create			Create	
	W <sub>2</sub> (A) <b>abort</b>					RT=2				
				R <sub>5</sub> (A)		RT=3			RT=5	
			R <sub>4</sub> (A)	W <sub>5</sub> (A)					RT=5	Create
R <sub>1</sub> (A)							RT=3			
C		C			X					
						X				

X means that we can delete this version