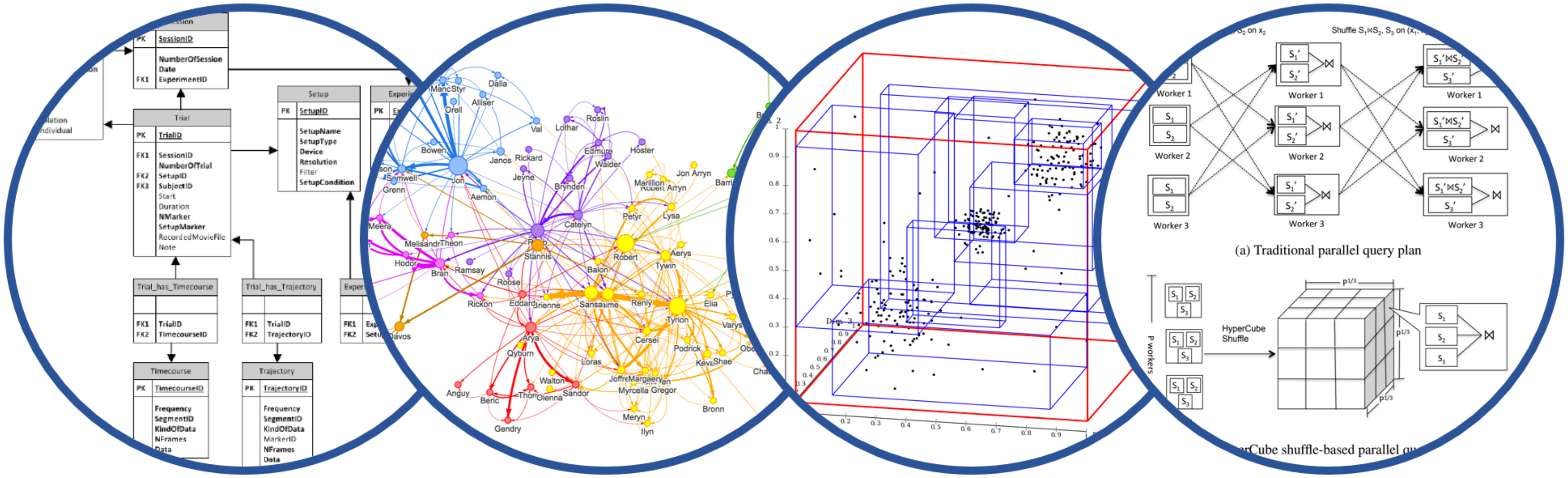

Database System Internals

Concurrency Control Intro

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Announcements

- HW 3 deadline extended to 2/16

- Lab 3 will be released tonight
  - Locking scheduler

- Quiz on Wednesday
  - Lab 1 material and how operators work
  - (No joins algorithms or cost estimations)

About Lab 3

- In lab 3, we implement transactions
- Focus on concurrency control
  - Want to run many transactions at the same time
  - Transactions want to read and write same pages
  - Will use locks to ensure conflict serializable execution
  - Use strict 2PL
- Build your own lock manager
  - Understand how locking works in depth
  - Ensure transactions rather than threads hold locks
    - Many threads can execute different pieces of the same transaction
    - Need to detect deadlocks and resolve them by aborting a transaction
  - But use Java synchronization to protect your data structures

Motivating Example

Client 1:
```
UPDATE Budget
SET money=money-100
WHERE pid = 1

UPDATE Budget
SET money=money+60
WHERE pid = 2

UPDATE Budget
SET money=money+40
WHERE pid = 3
```

Client 2:
```
SELECT sum(money)
FROM Budget
```

Would like to treat each group of instructions as a unit

Transaction

Definition: a transaction is a sequence of updates to the database with the property that either all complete, or none completes (all-or-nothing).

```
START TRANSACTION

[SQL statements]

COMMIT    or    ROLLBACK (=ABORT)
```

May be omitted if autocommit is off: first SQL query starts txn

In ad-hoc SQL: each statement = one transaction
This is referred to as autocommit

Motivating Example

```
START TRANSACTION
    UPDATE Budget
    SET money=money-100
    WHERE pid = 1

    UPDATE Budget
    SET money=money+60
    WHERE pid = 2

    UPDATE Budget
    SET money=money+40
    WHERE pid = 3
COMMIT  (or ROLLBACK)
```

```
SELECT sum(money)
FROM Budget
```

With autocommit and without **START TRANSACTION**, each SQL command is a transaction

ROLLBACK

- If the app gets to a place where it can't complete the transaction successfully, it can execute **ROLLBACK**

- This causes the system to "abort" the transaction
  - Database returns to a state without any of the changes made by the transaction

- Several reasons: user, application, system

Transactions

- Major component of database systems
- Critical for most applications; arguably more so than SQL

- Turing awards to database researchers:
  - Charles Bachman 1973
  - Edgar Codd 1981 for inventing relational dbs
  - **Jim Gray 1998 for inventing transactions**
  - Mike Stonebraker 2015 for INGRES and Postgres
    - And many other ideas after that

ACID Properties

ACID Properties

- **A**tomicity: Either all changes performed by transaction occur or none occurs
- **C**onsistency: A transaction as a whole does not violate integrity constraints
- **I**solation: Transactions appear to execute one after the other in sequence
- **D**urability: If a transaction commits, its changes will survive failures

What Could Go Wrong?

Why is it hard to provide ACID properties?

- **Concurrent** operations
  - Isolation problems
  - We saw one example earlier
- **Failures** can occur at any time
  - Atomicity and durability problems
  - Later lectures
- Transaction may need to **abort**

Terminology Needed For Lab 3

- **STEAL or NO-STEAL**
  - Can an update made by an uncommitted transaction overwrite the most recent committed value of a data item on disk?

- **FORCE or NO-FORCE**
  - Should all updates of a transaction be forced to disk before the transaction commits?

- Easiest for recovery: NO-STEAL/FORCE (lab 3)
- Highest performance: STEAL/NO-FORCE (lab 4)
- We will get back to this next week

Concurrent Execution Problems

- Write-read conflict: dirty read, inconsistent read
  - A transaction reads a value written by another transaction that has not yet committed
- Read-write conflict: unrepeatable read
  - A transaction reads the value of the same object twice. Another transaction modifies that value in between the two reads
- Write-write conflict: lost update
  - Two transactions update the value of the same object. The second one to write the value overwrites the first change

Schedules

A *schedule* is a sequence of interleaved actions from all transactions

Example

A and B are elements
in the database
t and s are variables
in tx source code

T1	T2
READ(A, t)	READ(A, s)
t := t+100	s := s*2
WRITE(A, t)	WRITE(A,s)
READ(B, t)	READ(B,s)
t := t+100	s := s*2
WRITE(B,t)	WRITE(B,s)

A Serial Schedule

T1	T2
READ(A, t)	
t := t+100	
WRITE(A, t)	
READ(B, t)	
t := t+100	
WRITE(B,t)	
	READ(A,s)
	s := s*2
	WRITE(A,s)
	READ(B,s)
	s := s*2
	WRITE(B,s)

A = 2
B = 2

A = 102
B = 102

A = 204
B = 204

A Serial Schedule

T1	T2	
		A = 2 B = 2
	READ(A,s)	
	s := s*2	
	WRITE(A,s)	
	READ(B,s)	
	s := s*2	
	WRITE(B,s)	A = 4 B = 4
READ(A, t)		
t := t+100		
WRITE(A, t)		
READ(B, t)		
t := t+100		
WRITE(B,t)		A = 104 B = 104

Serializable Schedule

A schedule is ***serializable*** if it is equivalent to a serial schedule

A Serializable Schedule

T1	T2	State
		A = 2 B = 2
READ(A, t) t := t+100 WRITE(A, t)		A = 102 B = 2
	READ(A,s) s := s*2 WRITE(A,s)	A = 204 B = 2
READ(B, t) t := t+100 WRITE(B,t)		A = 204 B = 102
	READ(B,s) s := s*2 WRITE(B,s)	A = 204 B = 204

This is a serializable schedule.
This is NOT a serial schedule

A Non-Serializable Schedule

T1	T2	
		A = 2 B = 2
READ(A, t) t := t+100 WRITE(A, t)		A = 102 B = 2
	READ(A,s) s := s*2 WRITE(A,s)	A = 204 B = 2
	READ(B,s) s := s*2 WRITE(B,s)	A = 204 B = 4
READ(B, t) t := t+100 WRITE(B,t)		A = 204 B = 104

Serializable Schedules

- The role of the scheduler is to ensure that the schedule is serializable

Q: Why not run only serial schedules ?
I.e. run one transaction after the other ?

Serializable Schedules

- The role of the scheduler is to ensure that the schedule is serializable

Q: Why not run only serial schedules ?
I.e. run one transaction after the other ?

A: Because of very poor throughput due to disk latency.

Lesson: main memory databases *may* schedule TXNs serially

Still Serializable, but…

T1	T2
READ(A, t)	
t := t+100	
WRITE(A, t)	
	READ(A,s)
	s := s + 200
	WRITE(A,s)
	READ(B,s)
	s := s + 200
	WRITE(B,s)
READ(B, t)	
t := t+100	
WRITE(B,t)	

Schedule is serializable because t=t+100 and s=s+200 commute

…we don't expect the scheduler to schedule this

To Be Practical

- Assume worst case updates:
  - Assume cannot commute actions done by transactions
- Therefore, we only care about reads and writes
  - Transaction = sequence of R(A)'s and W(A)'s

T_1: $r_1(A)$; $w_1(A)$; $r_1(B)$; $w_1(B)$
T_2: $r_2(A)$; $w_2(A)$; $r_2(B)$; $w_2(B)$

Conflicts

- Write-Read – WR
- Read-Write – RW
- Write-Write – WW

Conflict Serializability

Conflicts:

Two actions by same transaction T_i: $r_i(X); w_i(Y)$

Two writes by T_i, T_j to same element $w_i(X); w_j(X)$

Read/write by T_i, T_j to same element

$w_i(X); r_j(X)$

$r_i(X); w_j(X)$

Conflict Serializability

Definition A schedule is *conflict serializable* if it can be transformed into a serial schedule by a series of swappings of adjacent non-conflicting actions

- Every conflict-serializable schedule is serializable
- The converse is not true in general

Conflict Serializability

Example:

$r_1(A); w_1(A); r_2(A); w_2(A); r_1(B); w_1(B); r_2(B); w_2(B)$

Conflict Serializability

Example:

$r_1(A); w_1(A); r_2(A); w_2(A); r_1(B); w_1(B); r_2(B); w_2(B)$

$r_1(A); w_1(A); r_1(B); w_1(B); r_2(A); w_2(A); r_2(B); w_2(B)$

Conflict Serializability

Example:

$r_1(A); w_1(A); r_2(A); w_2(A); r_1(B); w_1(B); r_2(B); w_2(B)$

$r_1(A); w_1(A); r_1(B); w_1(B); r_2(A); w_2(A); r_2(B); w_2(B)$

Conflict Serializability

Example:

$r_1(A); w_1(A); r_2(A); w_2(A); r_1(B); w_1(B); r_2(B); w_2(B)$

$r_1(A); w_1(A); r_2(A); r_1(B); w_2(A); w_1(B); r_2(B); w_2(B)$

$r_1(A); w_1(A); r_1(B); w_1(B); r_2(A); w_2(A); r_2(B); w_2(B)$

Conflict Serializability

Example:

$r_1(A); w_1(A); r_2(A); w_2(A); r_1(B); w_1(B); r_2(B); w_2(B)$

$r_1(A); w_1(A); r_2(A); r_1(B); w_2(A); w_1(B); r_2(B); w_2(B)$

$r_1(A); w_1(A); r_1(B); w_1(B); r_2(A); w_2(A); r_2(B); w_2(B)$

Conflict Serializability

Example:

$r_1(A); w_1(A); r_2(A); w_2(A); r_1(B); w_1(B); r_2(B); w_2(B)$

$r_1(A); w_1(A); r_2(A); r_1(B); w_2(A); w_1(B); r_2(B); w_2(B)$

$r_1(A); w_1(A); r_1(B); r_2(A); w_2(A); w_1(B); r_2(B); w_2(B)$

....

$r_1(A); w_1(A); r_1(B); w_1(B); r_2(A); w_2(A); r_2(B); w_2(B)$

Testing for Conflict-Serializability

Precedence graph:

- A node for each transaction T_i,
- An edge from T_i to T_j whenever an action in T_i conflicts with, and comes before an action in T_j
- No edge for actions in the same transaction

- **The schedule is serializable iff the precedence graph is acyclic**

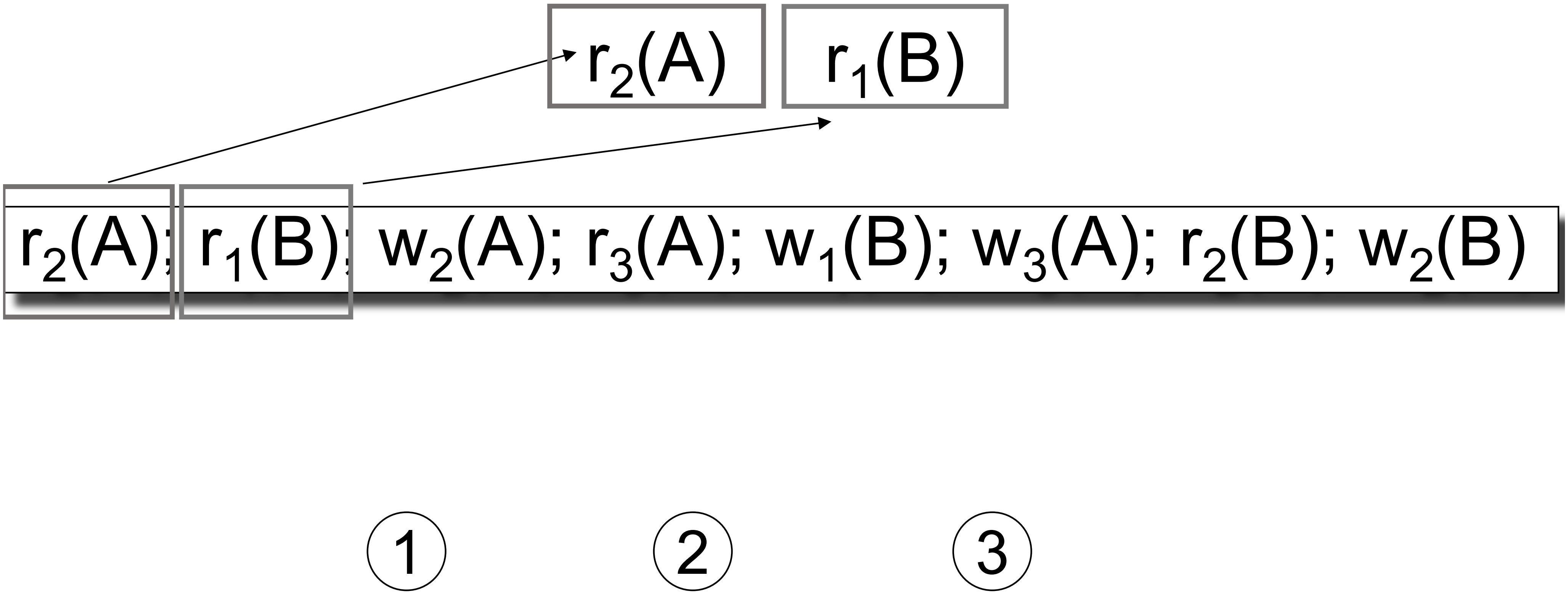
Testing for Conflict-Serializability

Important:

Always draw the full graph, unless ONLY asked if (yes or no) the schedule is conflict serializable

Example 1

$r_2(A); r_1(B); w_2(A); r_3(A); w_1(B); w_3(A); r_2(B); w_2(B)$


r2(A)
r1(B)
r2(A); r1(B); w2(A); r3(A); w1(B); w3(A); r2(B); w2(B)
1
2
3

$r_2(A)$ $r_1(B)$

$r_2(A)$; $r_1(B)$; $w_2(A)$; $r_3(A)$; $w_1(B)$; $w_3(A)$; $r_2(B)$; $w_2(B)$

1 2 3

$r_2(A)$ $r_1(B)$ No edge because no conflict (A != B)

$r_2(A); r_1(B); w_2(A); r_3(A); w_1(B); w_3(A); r_2(B); w_2(B)$

$r_2(A)$ $w_2(A)$

$r_2(A); r_1(B); w_2(A); r_3(A); w_1(B); w_3(A); r_2(B); w_2(B)$

1 2 3

$r_2(A)$ $w_2(A)$

No edge because same txn (2)

$r_2(A)$; $r_1(B)$; $w_2(A)$; $r_3(A)$; $w_1(B)$; $w_3(A)$; $r_2(B)$; $w_2(B)$

$r_2(A)$ $r_3(A)$?

$r_2(A); r_1(B); w_2(A); r_3(A); w_1(B); w_3(A); r_2(B); w_2(B)$

$r_2(A)$ $w_1(B)$?

$r_2(A); r_1(B); w_2(A); r_3(A); w_1(B); w_3(A); r_2(B); w_2(B)$

1 2 3

$r_2(A)$ $w_3(A)$?

$r_2(A); r_1(B); w_2(A); r_3(A); w_1(B); w_3(A); r_2(B); w_2(B)$

1 2 3

$r_2(A)$ $w_3(A)$

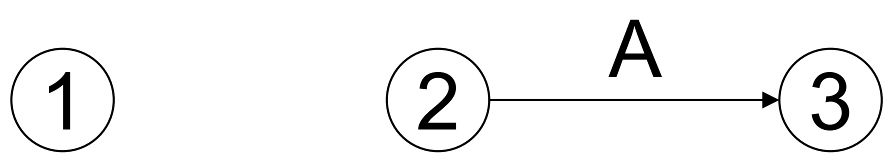
Edge! Conflict from T2 to T3

$r_2(A); r_1(B); w_2(A); r_3(A); w_1(B); w_3(A); r_2(B); w_2(B)$

$r_2(A)$ $w_3(A)$

Edge! Conflict from T2 to T3

$r_2(A); r_1(B); w_2(A); r_3(A); w_1(B); w_3(A); r_2(B); w_2(B)$



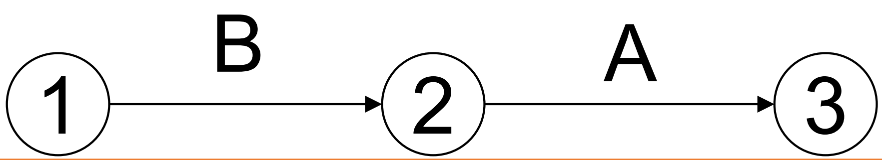
$r_2(A)$ $r_2(B)$?

$r_2(A); r_1(B); w_2(A); r_3(A); w_1(B); w_3(A); r_2(B); w_2(B)$

And so on until compared every pair of actions…

Example 1

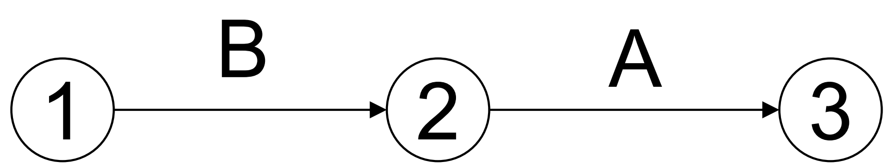
$r_2(A); r_1(B); w_2(A); r_3(A); w_1(B); w_3(A); r_2(B); w_2(B)$



More edges, but repeats of the same directed edge not necessary

Example 1

$r_2(A);\ r_1(B);\ w_2(A);\ r_3(A);\ w_1(B);\ w_3(A);\ r_2(B);\ w_2(B)$



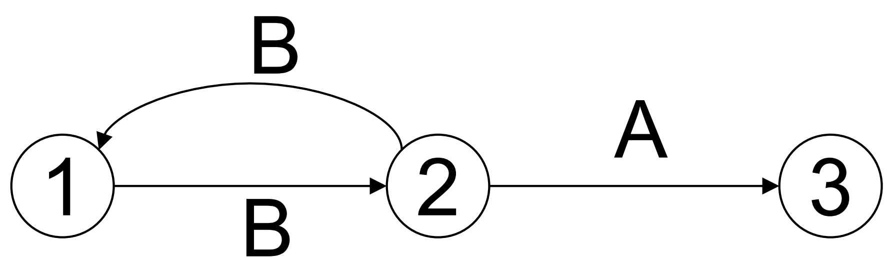
This schedule is **conflict-serializable**

Example 2

$r_2(A); r_1(B); w_2(A); r_2(B); r_3(A); w_1(B); w_3(A); w_2(B)$

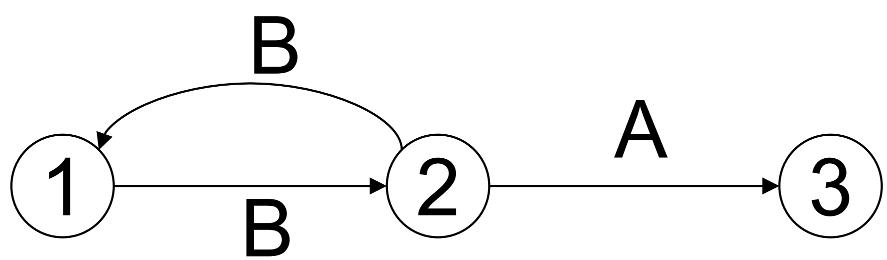
Example 2

$r_2(A); r_1(B); w_2(A); r_2(B); r_3(A); w_1(B); w_3(A); w_2(B)$



Example 2

$r_2(A); r_1(B); w_2(A); r_2(B); r_3(A); w_1(B); w_3(A); w_2(B)$



This schedule **is NOT conflict-serializable**

View Equivalence

- A serializable schedule need not be conflict serializable, even under the “worst case update” assumption

$$w_1(X); w_2(X); w_2(Y); w_1(Y); w_3(Y);$$

Is this schedule conflict-serializable ?

View Equivalence

- A serializable schedule need not be conflict serializable, even under the “worst case update” assumption

$w_1(X); w_2(X); w_2(Y); w_1(Y); w_3(Y);$

Is this schedule conflict-serializable ?

No…

View Equivalence

- A serializable schedule need not be conflict serializable, even under the “worst case update” assumption

$$w_1(X); w_2(X); w_2(Y); w_1(Y); w_3(Y);$$

Lost write

$$w_1(X); w_1(Y); w_2(X); w_2(Y); w_3(Y);$$

Equivalent, but not conflict-equivalent

View Equivalence

T1	T2	T3
W1(X)		
	W2(X)	
	W2(Y)	
	CO2	
W1(Y)		
CO1		
		W3(Y)
		CO3

Lost

T1	T2	T3
W1(X)		
W1(Y)		
CO1		
	W2(X)	
	W2(Y)	
	CO2	
		W3(Y)
		CO3

Serializable, but not conflict serializable

View Equivalence

Two schedules S, S’ are *view equivalent* if:

- If T reads an initial value of A in S,
  then T reads the initial value of A in S’

- If T reads a value of A written by T’ in S,
  then T reads a value of A written by T’ in S’

- If T writes the final value of A in S,
  then T writes the final value of A in S’

View-Serializability

A schedule is *view serializable* if it is view equivalent to a serial schedule

Remark:

- If a schedule is *conflict serializable*, then it is also *view serializable*
- But not vice versa

Schedules with Aborted Transactions

- When a transaction aborts, the recovery manager undoes its updates
- But some of its updates may have affected other transactions !

Schedules with Aborted Transactions

T1	T2
R(A)	
W(A)	
	R(A)
	W(A)
	R(B)
	W(B)
	Commit
Abort	

What's wrong?

Schedules with Aborted Transactions

T1	T2
R(A)	
W(A)	
	R(A)
	W(A)
	R(B)
	W(B)
	Commit
Abort	

What's wrong?

Cannot abort T1 because cannot undo T2

Recoverable Schedules

A schedule is *recoverable* if:

- It is conflict-serializable, and
- Whenever a transaction T commits, all transactions that have written elements read by T have already committed

Recoverable Schedules

A schedule is *recoverable* if:

- It is conflict-serializable, and
- Whenever a transaction T commits, all transactions that have written elements read by T have already committed

Recoverable Schedules

T1	T2
R(A)	
W(A)	
	R(A)
	W(A)
	R(B)
	W(B)
	Commit
?	

Nonrecoverable

T1	T2
R(A)	
W(A)	
	R(A)
	W(A)
	R(B)
	W(B)
Commit	
	Commit

Recoverable

Recoverable Schedules

T1	T2	T3	T4
R(A)			
W(A)			
	R(A)		
	W(A)		
	R(B)		
	W(B)		
		R(B)	
		W(B)	
		R(C)	
		W(C)	
			R(C)
			W(C)
			R(D)
			W(D)
Abort			

How do we recover ?

Cascading Aborts

- If a transaction T aborts, then we need to abort any other transaction T' that has read an element written by T

- A schedule *avoids cascading aborts* if whenever a transaction reads an element, the transaction that has last written it has already committed.

We base our locking scheme on this rule!

Avoiding Cascading Aborts

T1	T2
R(A)	
W(A)	
	R(A)
	W(A)
	R(B)
	W(B)
. . .	
	. . .

With cascading aborts

T1	T2
R(A)	
W(A)	
Commit	
	R(A)
	W(A)
	R(B)
	W(B)
	. . .

Without cascading aborts

Serializability

- Serial
- Serializable
- Conflict serializable
- View serializable

Recoverability

- Recoverable
- Avoids cascading deletes

Scheduler

- The scheduler:
- Module that schedules the transaction's actions, ensuring serializability

- Two main approaches
  - Pessimistic: locks
  - Optimistic: timestamps, multi-version, validation