





/hat We Already Know
Supplier (sno, sname, scity, sstate) Part (pno, pname, psize, pcolor) Supply (sno, pno, price)
For each SQL query
SELECT S.sname FROM Supplier S, Supply U WHERE S.scity='Seattle' AND S.sstate='WA' AND S.sno = U.sno AND U.pno = 2
There exist many logical query plans
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# What We Also Know

- For each logical plan...
- There exist many physical plans









# Query Optimizer Overview

- Input: A logical query plan
- Output: A good physical query plan
- Basic query optimization algorithm
  - Enumerate alternative plans (logical and physical)

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- Compute estimated cost of each plan
  - Compute number of I/Os
  - · Optionally take into account other resources
- · Choose plan with lowest cost
- · This is called cost-based optimization

# Query Optimizer Overview • Input: A logical query plan • Output: A good physical query plan

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# Observations



# In order to make the right choice

- Need to have statistics over the data
- The B's, the T's, the V's
- Commonly: histograms over base data
  - In SimpleDB as well... lab 5.

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Commutativity, Associativity, Distributivity
$R \cup S = S \cup R, \ R \cup (S \cup T) = (R \cup S) \cup T$
$R \bowtie S = S \bowtie R, R \bowtie (S \bowtie T) = (R \bowtie S) \bowtie T$
$R\bowtie(S\cupT)\ =\ (R\bowtieS)\cup(R\bowtieT)$
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# Laws Involving Projections

 $\Pi_{\mathsf{M}}(\mathsf{R}\bowtie\mathsf{S})=\Pi_{\mathsf{M}}(\Pi_{\mathsf{P}}(\mathsf{R})\bowtie\Pi_{\mathsf{Q}}(\mathsf{S}))$ 

$$\label{eq:matrix} \begin{split} \Pi_{M}(\Pi_{N}(\mathsf{R})) &= \Pi_{M}(\mathsf{R}) \\ /^{*} \text{ note that } \mathsf{M} \subseteq \mathsf{N} \ ^{*} / \end{split}$$

 Example R(A,B,C,D), S(E, F, G) Π<sub>A,B,G</sub>(R ⋈ <sub>D=E</sub> S) = Π<sub>A,B,G</sub> (Π<sub>A,B,D</sub>(R) ⋈ <sub>D=E</sub> Π<sub>E,G</sub>(S))

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Laws for grouping and aggregation  $\gamma_{A, agg(D)}(R(A,B) \bowtie_{B=C} S(C,D)) = \gamma_{A, agg(D)}(R(A,B) \bowtie_{B=C} (\gamma_{C, agg(D)}S(C,D)))$ 

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# Search Space Challenges

### Search space is huge!

- · Many possible equivalent trees
- · Many implementations for each operator
- Many access paths for each relation
  - File scan or index + matching selection condition
- Cannot consider ALL plans
  - Heuristics: only partial plans with "low" cost

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# Key Decisions

### Logical plan

- What logical plans do we consider (left-deep, bushy?) Search Space
- Which algebraic laws do we apply, and in which context(s)? Optimization rules
- In what order do we explore the search space? Optimization algorithm

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# Even More Key Decisions!

# Physical plan

- What physical operators to use?
- What access paths to use (file scan or index)?
- Pipeline or materialize intermediate results?
- These decisions also affect the search space

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