

# Database System Internals Architecture

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# Announcements

- Lab 1 part 1 is due on Monday at 11pm
  - Lab 1 in full is due on April 14th
  - “git pull upstream master” before building
  - Remember to git commit and git push often!
- HW1 is due next week on Friday
- 544 paper reviews can be handed in to me via email or Gradescope

# What we already know...

- **Database** = collection of related files
- **DBMS** = program that manages the database

# What we already know...

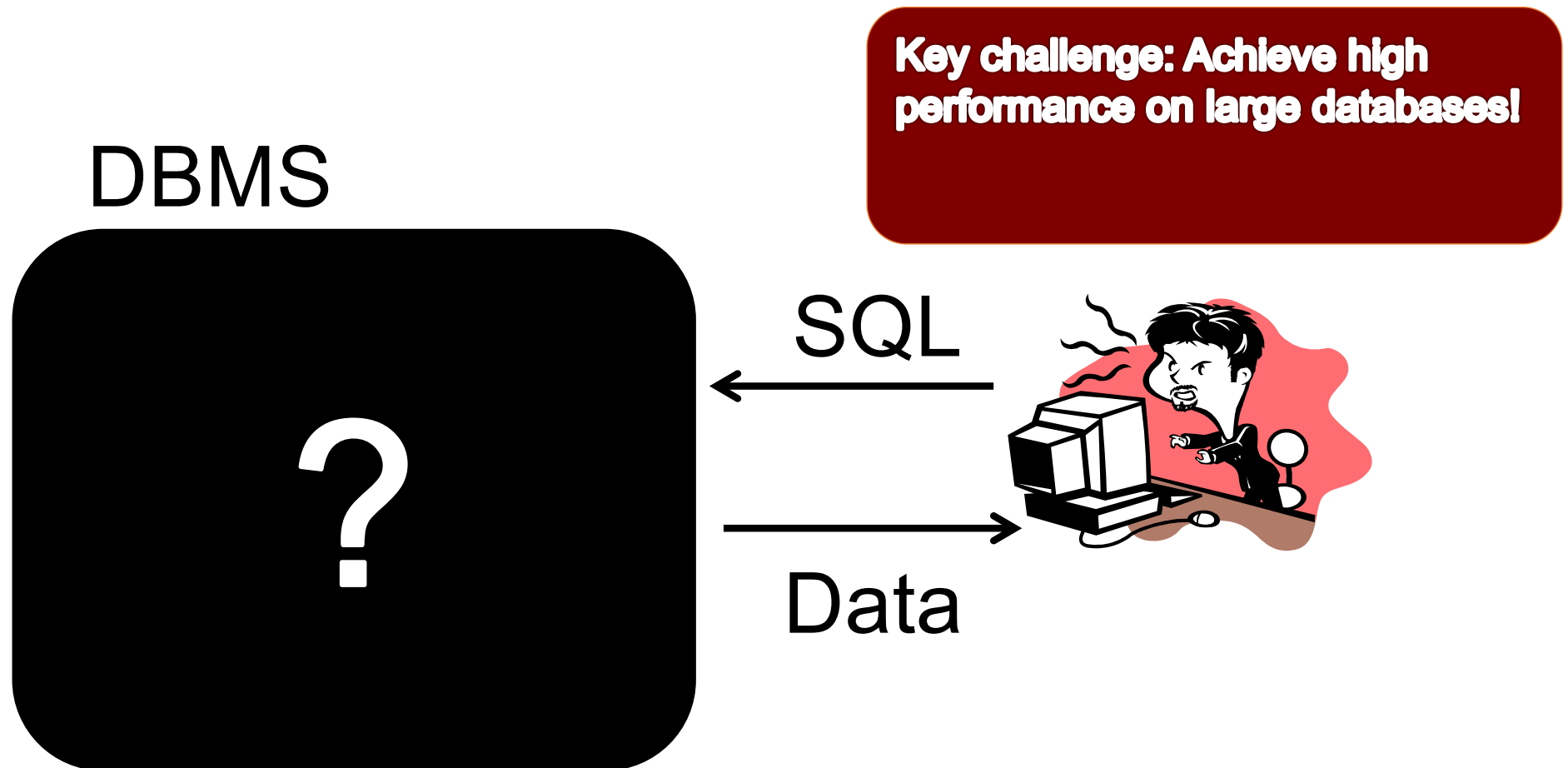
- **Data models**: relational, semi-structured (XML), graph (RDF), key-value pairs
- **Relational model**: defines only the logical model, and does not define a physical storage of the data

# What we already know...

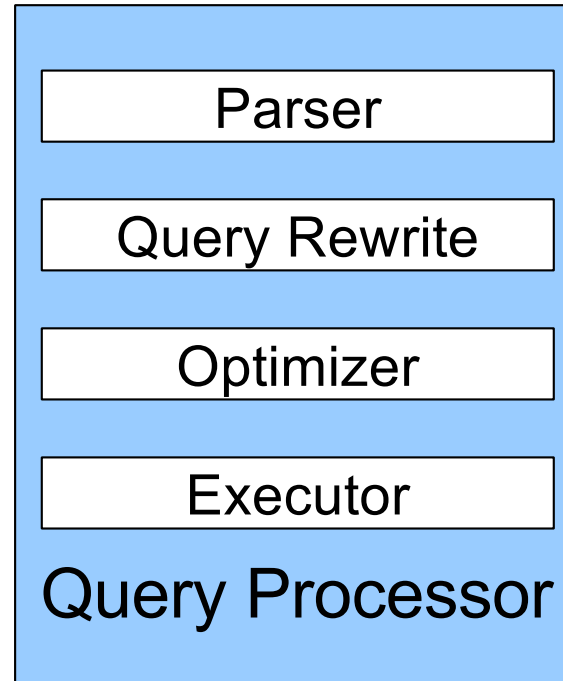
## Relational Query Language:

- **Set-at-a-time**: instead of tuple-at-a-time
- **Declarative**: user says what they want and not how to get it
- **Query optimizer**: from *what* to *how*

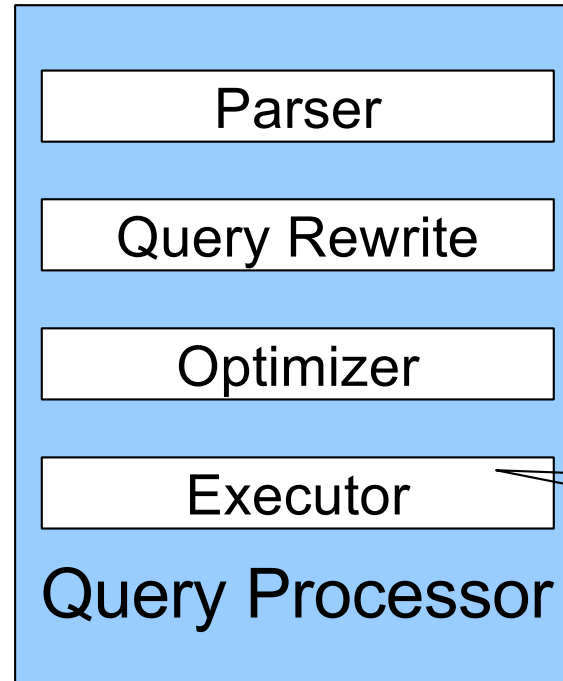
# How to Implement an RDBMS?



# DBMS Architecture



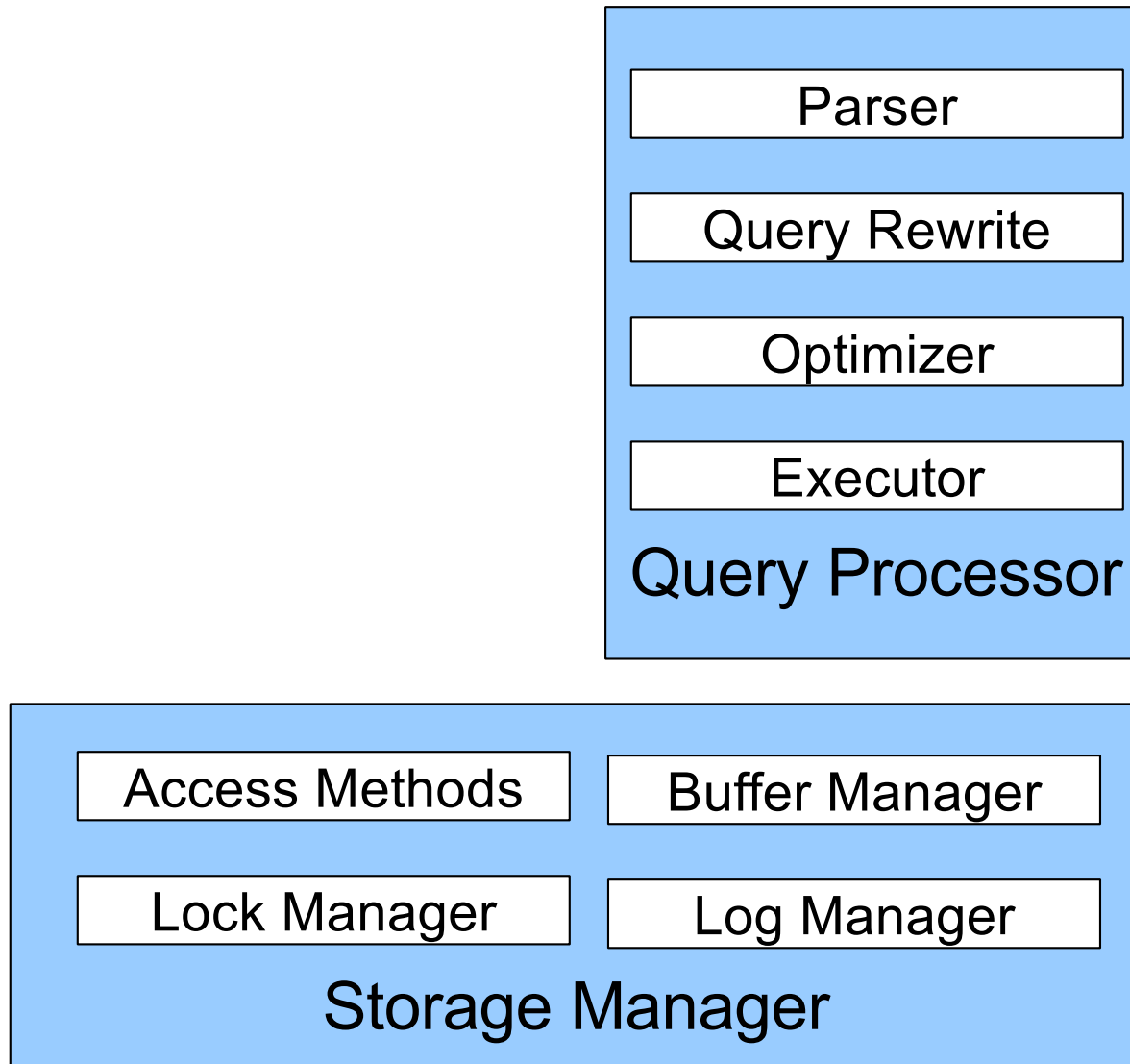
# DBMS Architecture



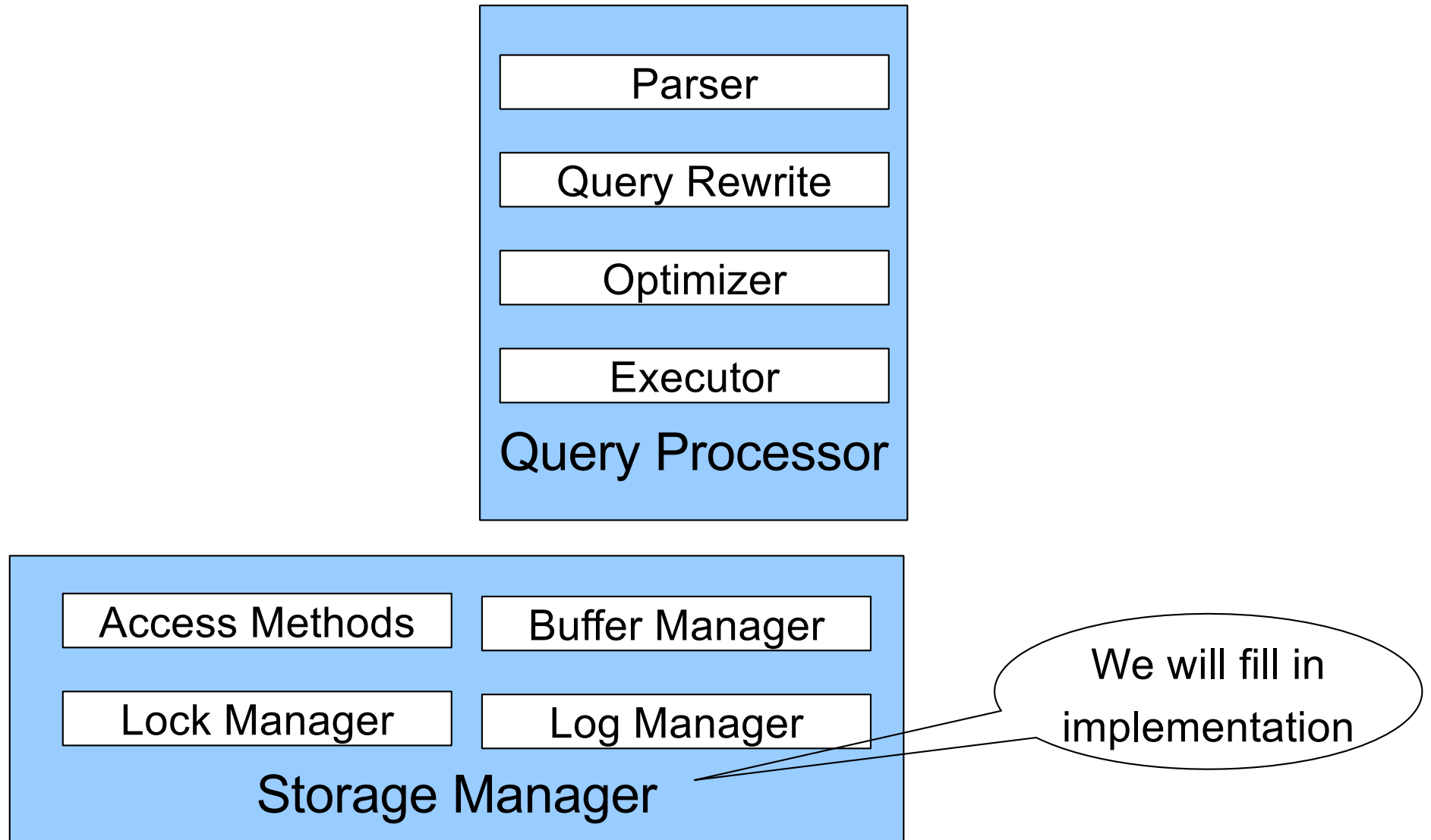
We will fill in  
implementation



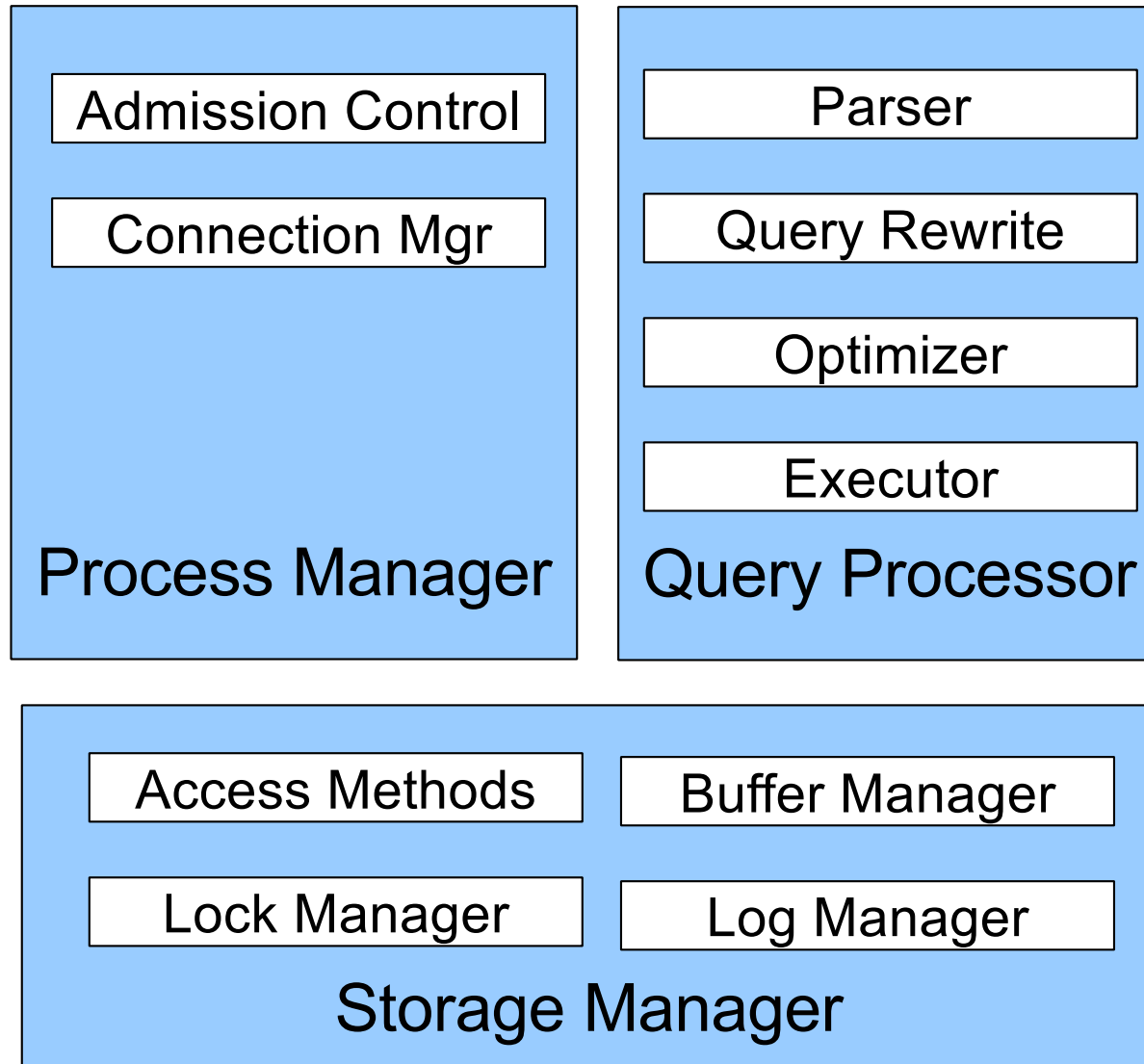
# DBMS Architecture



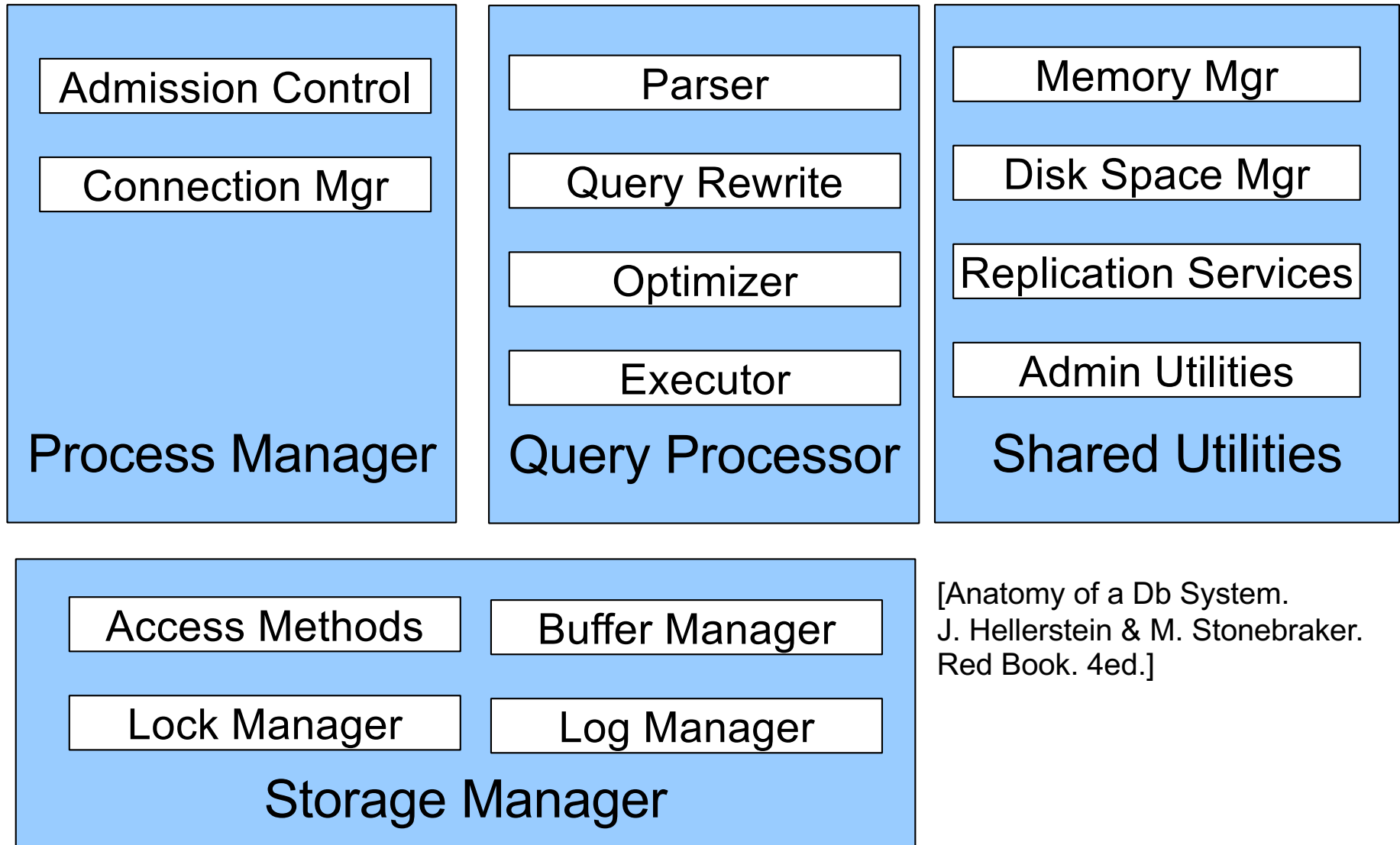
# DBMS Architecture



# DBMS Architecture



# DBMS Architecture



[Anatomy of a Db System.  
J. Hellerstein & M. Stonebraker.  
Red Book. 4ed.]

# Goal for Today

Overview of query execution

Overview of storage manager

# Query Processor

# Example Database Schema

Supplier (sno, sname, scity, sstate)

Part (pno, pname, psize, pcolor)

Supplies (sno, pno, price)

## View: Suppliers in Seattle

```
CREATE VIEW NearbySupp AS
SELECT sno, sname
FROM Supplier
WHERE scity='Seattle' AND sstate='WA'
```

# Example Query

```
Supplier (sno, sname, scity, sstate)  
Part (pno, pname, psize, pcolor)  
Supplies (sno, pno, price)
```

- Find the names of all suppliers in Seattle who supply part number 2

```
SELECT sname  
FROM NearbySupp  
WHERE sno IN ( SELECT sno  
                FROM Supplies  
                WHERE pno = 2 )
```



# Query Processor

- **Step 1: Parser**

- Parses query into an internal format
- Performs various checks using **catalog**

- **Step 2: Query rewrite**

- View rewriting, flattening, etc.

# Rewritten Version of Our Query

```
Supplier (sno, sname, scity, sstate)  
Part (pno, pname, psize, pcolor)  
Supplies (sno, pno, price)
```

Original query:

```
SELECT sname  
FROM NearbySupp  
WHERE sno IN ( SELECT sno  
                FROM Supplies  
                WHERE pno = 2 )
```

Rewritten query (expanding NearbySupp view):

```
SELECT S.sname  
FROM Supplier S, Supplies U  
WHERE S.scity='Seattle' AND S.sstate='WA'  
AND S.sno = U.sno  
AND U.pno = 2;
```

# Query Processor

## ▪ Step 3: Optimizer

- Find an efficient query plan for executing the query
- **A query plan is**
  - **Logical:** An extended relational algebra tree
  - **Physical:** With additional annotations at each node
    - Access method to use for each relation
    - Implementation to use for each relational operator

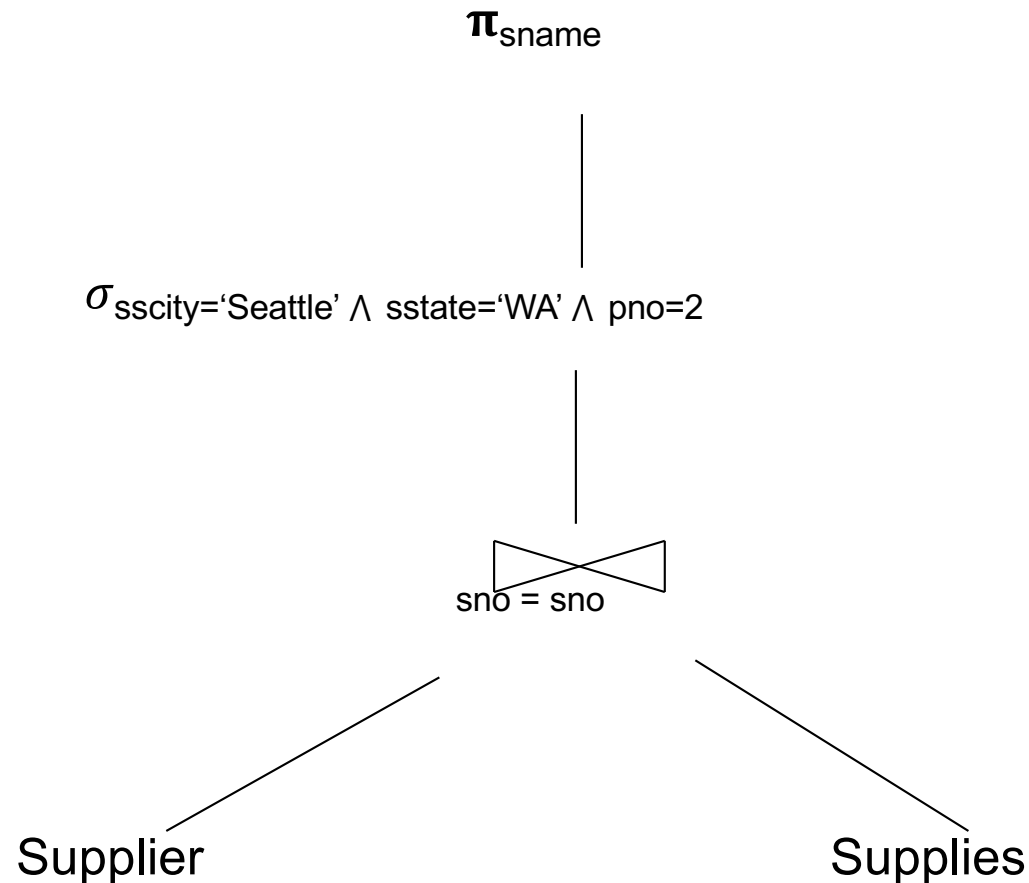
## ▪ Step 4: Executor

- Actually executes the physical plan

# Logical Query Plan

```
SELECT S.sname
FROM Supplier S, Supplies U
WHERE
S.scity='Seattle'
AND S.sstate='WA'
AND S.sno = U.sno
AND U.pno = 2;
```

Supplier (sno, sname, scity, sstate)  
Part (pno, pname, psize, pcolor)  
Supplies (sno, pno, price)



# Physical Query Plan

- Logical query plan with extra annotations
- **Implementation choice** for each operator
- **Access path selection** for each relation
  - Bottom of tree = read from disk
  - Use a file scan or use an index

# Physical Query Plan

Supplier (sno, sname, scity, sstate)  
Part (pno, pname, psize, pcolor)  
Supplies (sno, pno, price)

(On the fly)

$\pi_{\text{sname}}$

(On the fly)

$\sigma_{\text{scity}='Seattle' \wedge \text{sstate}='WA' \wedge \text{pno}=2}$

(Nested loop)

sno = sno

Suppliers

(File scan)

Supplies

(File scan)

# Query Executor

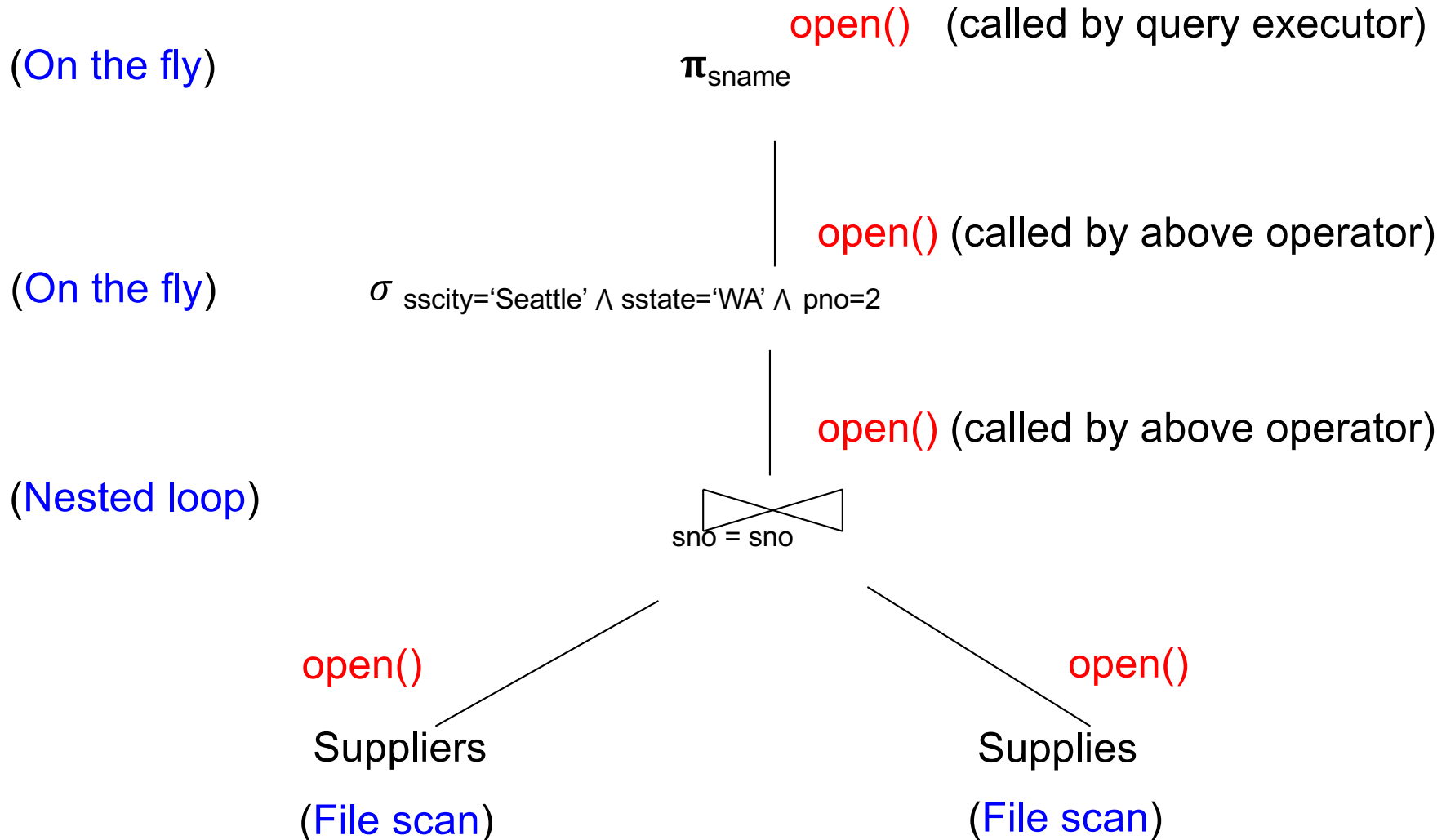
# Iterator Interface

- Each **operator implements Operator.java**
- **open()**
  - Initializes operator state
  - Sets parameters such as selection predicate
- **next()**
  - **Returns a Tuple!**
  - Operator invokes next() recursively on its inputs
  - Performs processing and produces an output tuple
- **close()**: clean-up state
- Operators also have reference to their **child** operator in the query plan



# Query Execution

Supplier (sno, sname, scity, sstate)  
Part (pno, pname, psize, pcolor)  
Supplies (sno, pno, price)



# Query Execution

Supplier (sno, sname, scity, sstate)  
Part (pno, pname, psize, pcolor)  
Supplies (sno, pno, price)

(On the fly)

$\pi_{\text{sname}}$  next()

pull-based execution

(On the fly)

$\sigma_{\text{scity}='Seattle' \wedge \text{sstate}='WA' \wedge \text{pno}=2}$  next()

(Nested loop)

next()  
sno = sno

next()

Suppliers

(File scan)

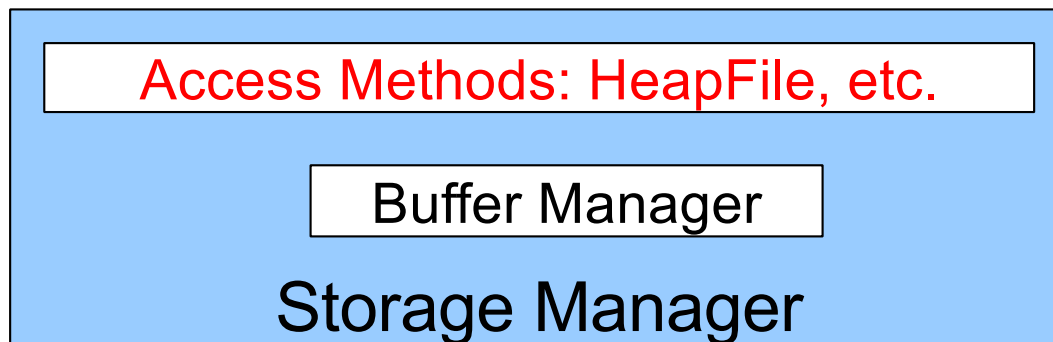
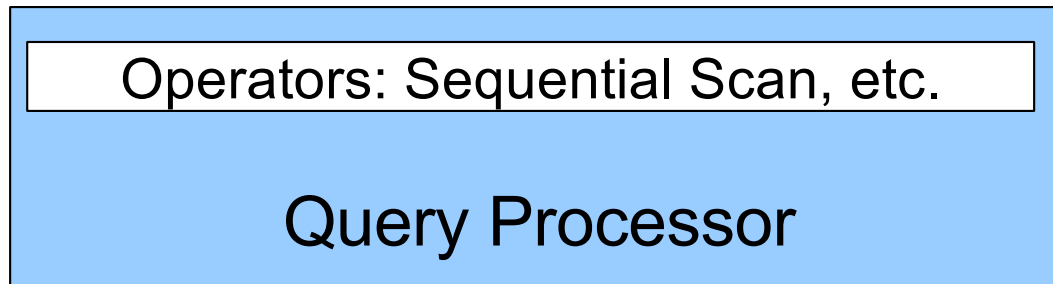
next()

Supplies

(File scan)

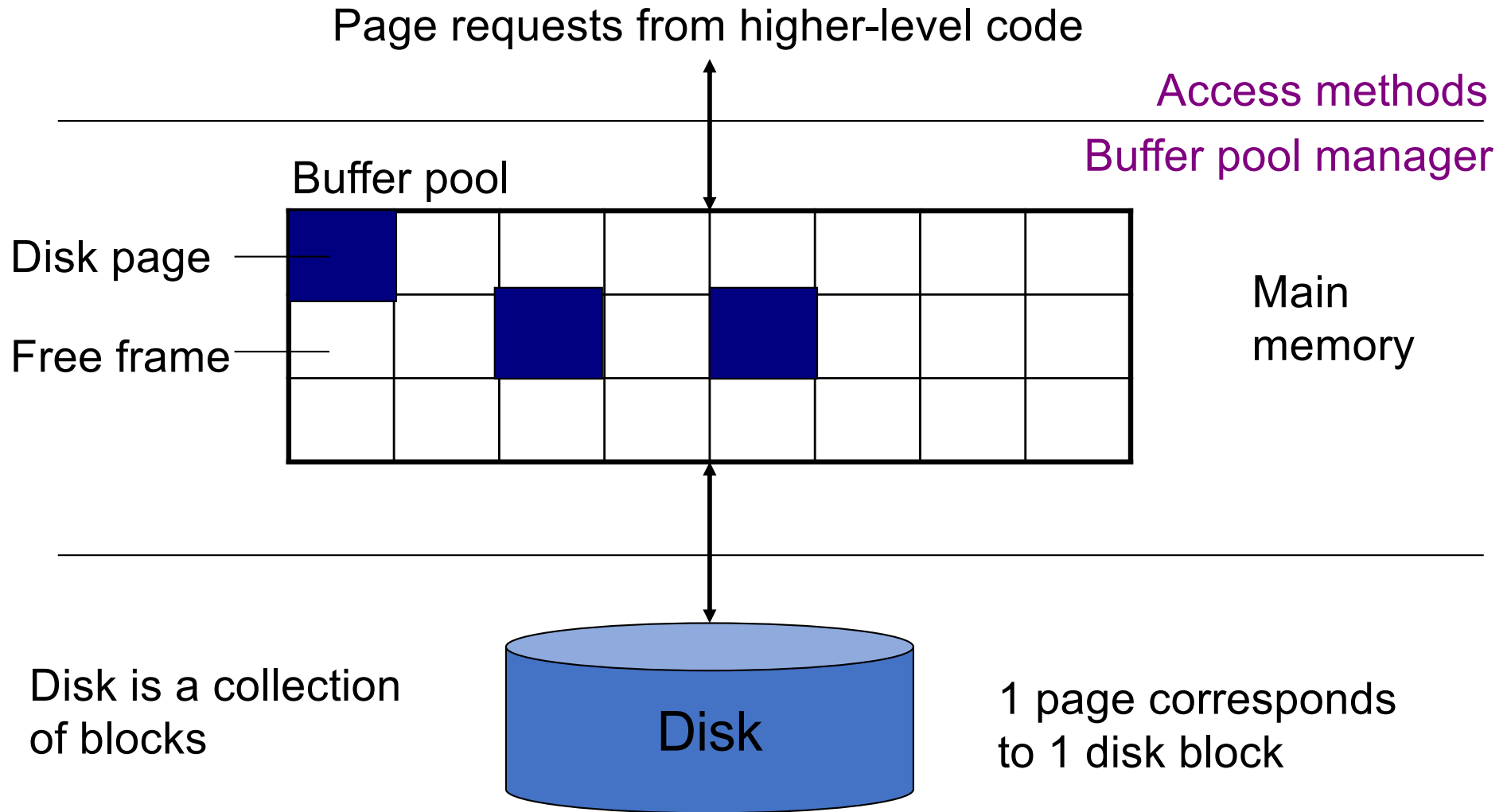
# Storage Manager

# Access Methods



- **Operators:** Process data
- **Access methods:** Organize data to support fast access to desired subsets of records
- **Buffer manager:** Caches data in memory. Reads/writes data to/from disk as needed
- **Disk-space manager:** Allocates space on disk for files/access methods

# Buffer Manager (BufferPool in SimpleDB)



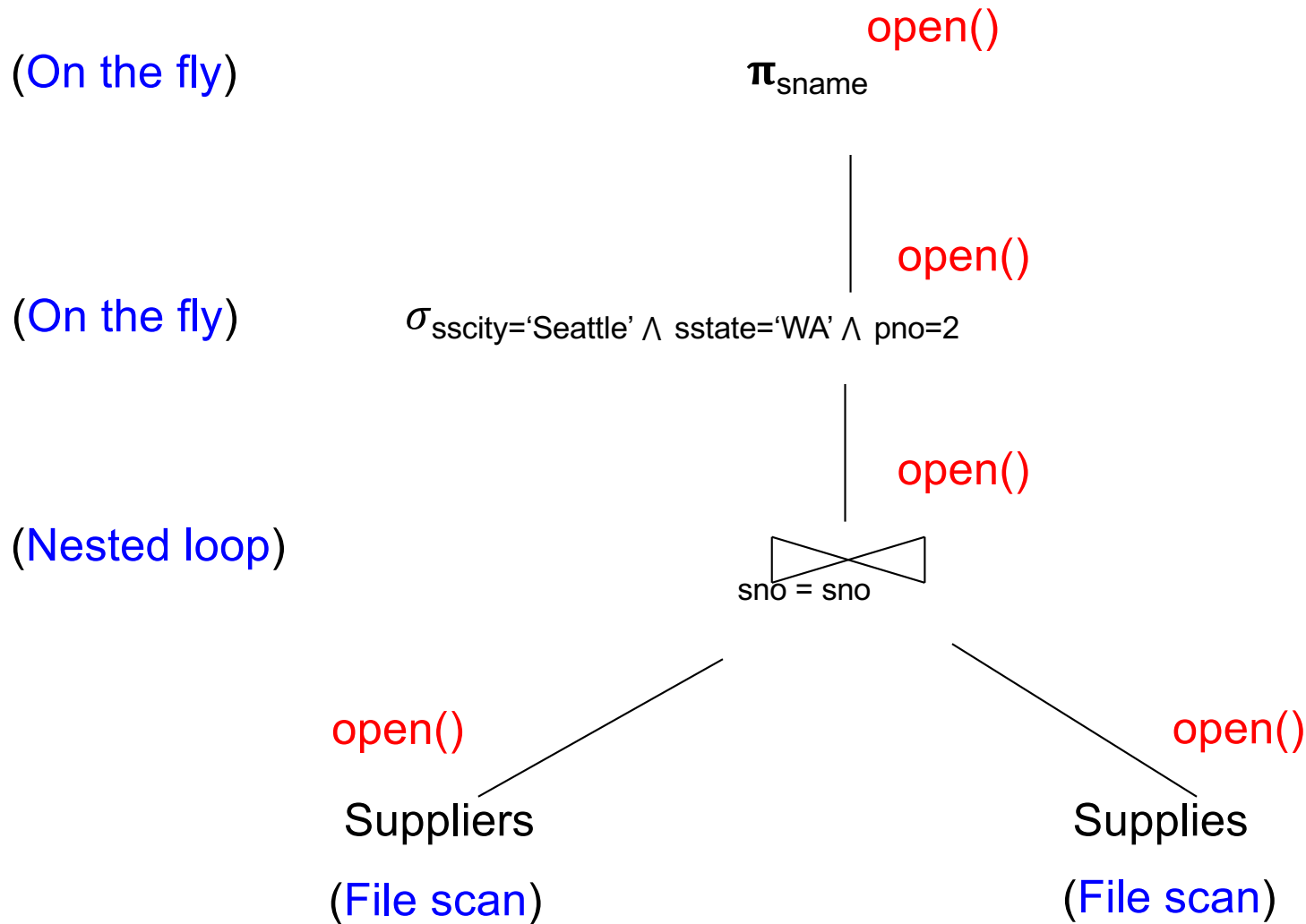
# Buffer Manager

- Brings pages in from memory and caches them
- Eviction policies
  - Random page (ok for SimpleDB)
  - Least-recently used
  - The “clock” algorithm (see book)
- Keeps track of which **pages are dirty**
  - A dirty page has changes not reflected on disk
  - Implementation: Each page includes a dirty bit

# Access Methods

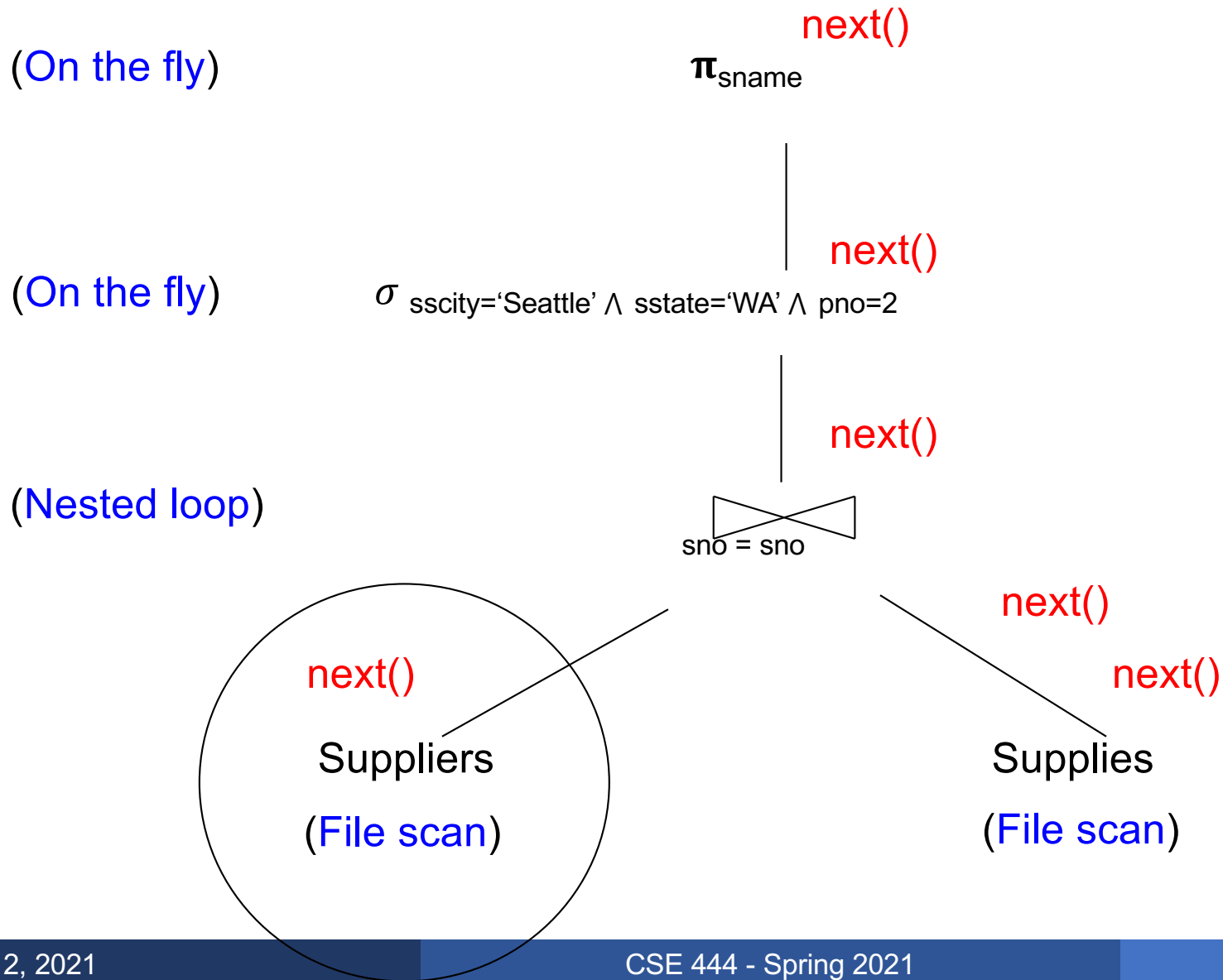
- A DBMS stores data on disk by breaking it into *pages*
  - A page is the size of a disk block.
  - A page is the unit of disk IO
- Buffer manager caches these pages in memory
- Access methods do the following:
  - They organize pages into collections called DB *files*
  - They organize data inside pages
  - They provide an API for operators to access data in these files
- Discussion:
  - OS vs DBMS files
  - OS vs DBMS buffer manager

# Query Execution





# Query Execution



# Query Execution In SimpleDB

open()

next()

**SeqScan**

Operator at  
bottom of plan

open()

next()

**Heap File Access Method**

In SimpleDB, SeqScan can  
find HeapFile in Catalog

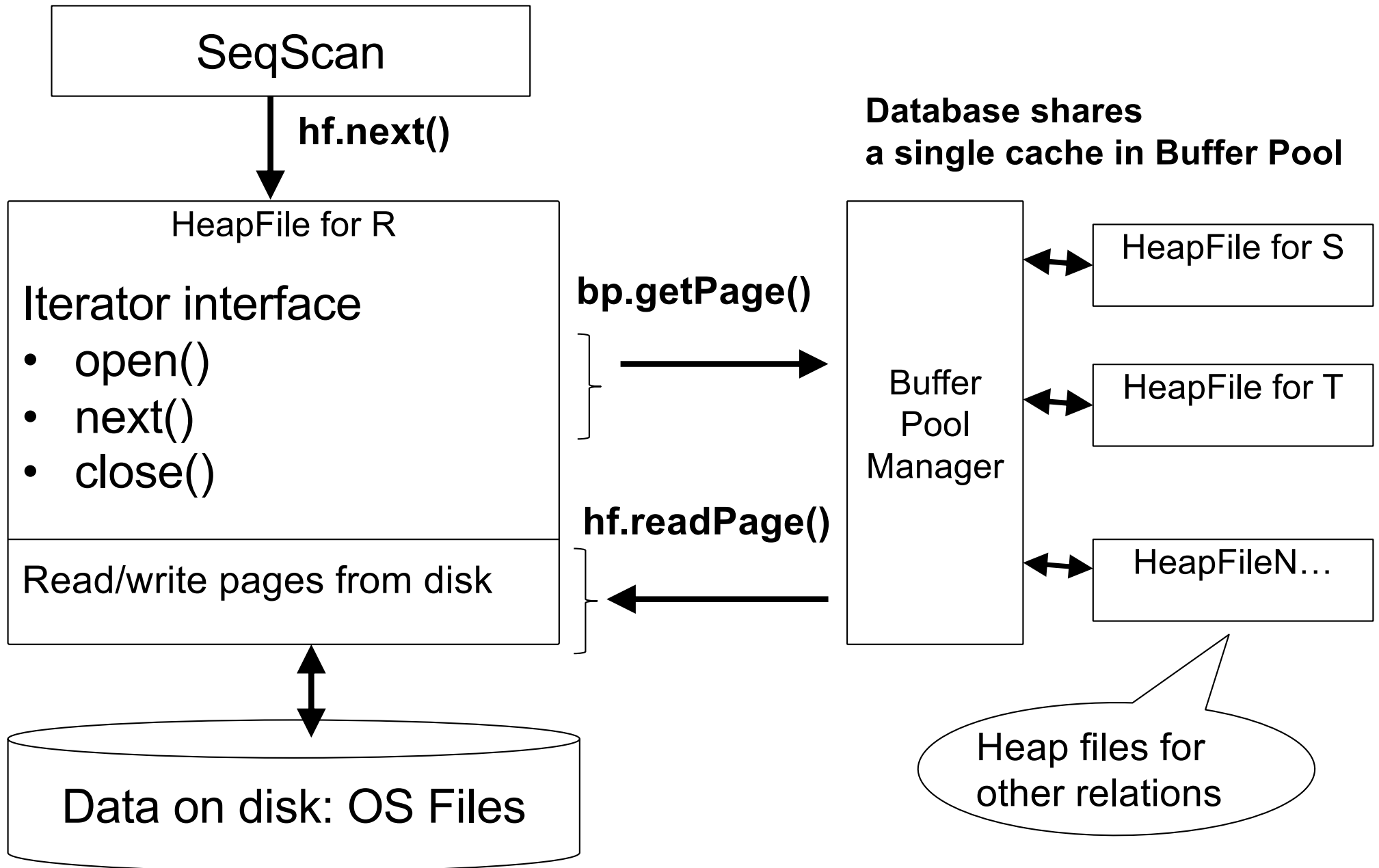
Offers iterator interface

- open()
- next()
- close()

Knows how to read/write pages from disk

But if Heap File reads data  
directly from disk, it will not  
stay cached in Buffer Pool!

# Query Execution In SimpleDB



# HeapFile In SimpleDB

- Data is stored on disk in an OS file. HeapFile class knows how to “decode” its content
- Control flow:

SeqScan calls methods such as "iterate" on the HeapFile Access Method

During the iteration, the HeapFile object needs to call the BufferManager.getPage() method to ensure that necessary pages get loaded into memory.

The BufferManager will then call HeapFile .readPage()/writePage() page to actually read/write the page.