



- What happens when a secondary crashes? – Nothing happens
  - When secondary recovers, it catches up
- What happens when the master/primary fails?

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- Blocking would hurt availability
- Must chose a new primary: run election



- Secondaries think that primary failed
- Secondaries elect a new primary
- But primary can still be running
- Now have two primaries!

# Majority Consensus

• To avoid problem, only majority partition can continue processing at any time

- In general,
  - Whenever a replica fails or recovers...
  - a set of communicating replicas must determine...
  - whether they have a majority before they can continue

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#### • With n copies

- Exclusive lock on x copies is global exclusive lock
- Shared lock on s copies is global shared lock
- Must have: 2x > n and s + x > n
- Majority locking
  - -s = x = [(n+1)/2]
  - No need to run any reconfiguration algorithms
- Read-locks-one, write-locks-all
  - s=1 and x = n, high read performance
  - Need to make sure algo runs on quorum of computers

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## Synchronous Replication Properties

- Favours consistency over availability
- Only majority partition can process requests
- There appears to be a single copy of the db
- · High runtime overhead
  - Must lock and update at least majority of replicas
  - Two-phase commit
  - Runs at pace of slowest replica in quorum
  - So overall system is now slower
  - Higher deadlock rate (transactions take longer)

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### Asynchronous Replication

- Also called lazy replication
- Also called optimistic replication
- · Main goals: availability and performance
- · Approach
  - One replica updated by original transaction
  - Updates propagate asynchronously to other replicas

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- One master holds primary copy
  - Transactions update primary copy
  - Master asynchronously propagates updates to replicas, which process them in same order (e.g. through log shipping)
  - Ensures single-copy serializability
- What happens when master/primary fails?
  - Can lose most recent transactions when primary fails!
    After electing a new primary, secondaries must agree who is most up-to-date

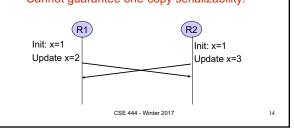
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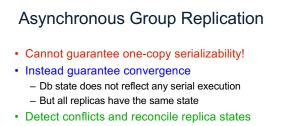
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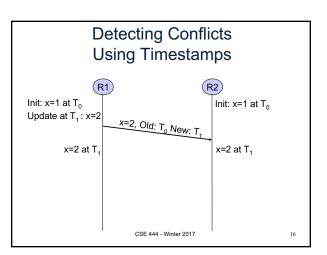


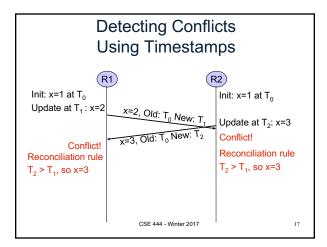
Cannot guarantee one-copy serializability!





- · Different reconciliation techniques are possible
  - Manual
  - Most recent timestamp wins
  - Site A wins over site B
  - User-defined rules, etc. CSE 444 - Winter 2017







#### **Two-Tier Replication**

- Benefits of lazy master and lazy group
- · Each object has a master with primary copy
- When disconnected from master
  - Secondary can only run tentative transactions
- When reconnects to master
  - Master reprocesses all tentative transactions
  - Checks an acceptance criterion
  - If passes, we now have final commit order
  - Secondary undoes tentative and redoes committed

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#### Conclusion

- Replication is a very important problem
  - Fault-tolerance (various forms of replication)
  - Caching (lazy master)
  - Warehousing (lazy master)
  - Mobility (two-tier techniques)
- Replication is complex, but basic techniques and trade-offs are very well known
  - Synchronous or asynchronous replication
  - Master or quorum

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