CSE 444: Database Internals

Lecture 25 Replication

References

- Ullman Book Chapter 20.6
- Database management systems.

Ramakrishnan and Gehrke.

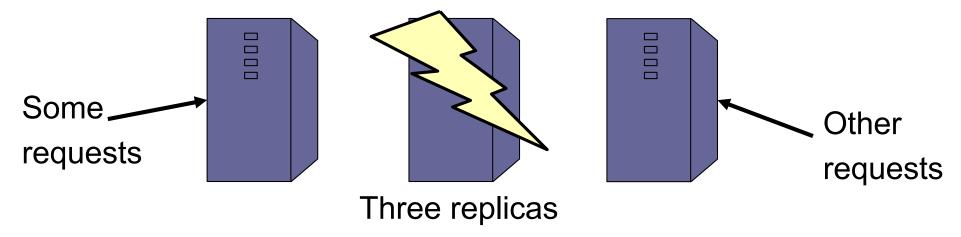
Third Ed. Chapter 22.11

Outline

- Goals of replication
- Three types of replication
 - Eager replication
 - Lazy replication
 - Two-tier replication

Goals of Replication

- Goal 1: availability
- Goal 2: performance



 But, it's easy to build a replicated system that reduces performance and availability

Synchronous Replication

- Also called eager replication
- All updates are applied to all replicas (or to a majority)
 as part of a single transaction (need two phase commit)
- Main goal: as if there was only one copy
 - Maintain consistency
 - Maintain one-copy serializability
 - I.e., execution of transactions has same effect as an execution on a non-replicated db
- Transactions must acquire global locks

Synchronous Master Replication

- One master for each object holds primary copy
 - The "Master" is also called "Primary"
 - To update object, transaction must acquire a lock at the master
 - Lock at the master is global lock
- Master propagates updates to replicas synchronously
 - Updates propagate as part of the same distributed transaction
 - Need to run 2PC at the end
 - For example, using triggers

Crash Failures

- What happens when a secondary crashes?
 - Nothing happens
 - When secondary recovers, it catches up
- What happens when the master/primary fails?
 - Blocking would hurt availability
 - Must chose a new primary: run election

Network Failures

- Network failures can cause trouble...
 - Secondaries think that primary failed
 - Secondaries elect a new primary
 - But primary can still be running
 - Now have two primaries!

Majority Consensus

- To avoid problem, only majority partition can continue processing at any time
- In general,
 - Whenever a replica fails or recovers...
 - a set of communicating replicas must determine...
 - whether they have a majority before they can continue

Synchronous Group Replication

With n copies

- Exclusive lock on x copies is global exclusive lock
- Shared lock on s copies is global shared lock
- Must have: 2x > n and s + x > n

Majority locking

- s = x = [(n+1)/2]
- No need to run any reconfiguration algorithms
- Read-locks-one, write-locks-all
 - s=1 and x = n, high read performance
 - Need to make sure algo runs on quorum of computers

Synchronous Replication Properties

- Favours consistency over availability
 - Only majority partition can process requests
 - There appears to be a single copy of the db
- High runtime overhead
 - Must lock and update at least majority of replicas
 - Two-phase commit
 - Runs at pace of slowest replica in quorum
 - So overall system is now slower
 - Higher deadlock rate (transactions take longer)

Asynchronous Replication

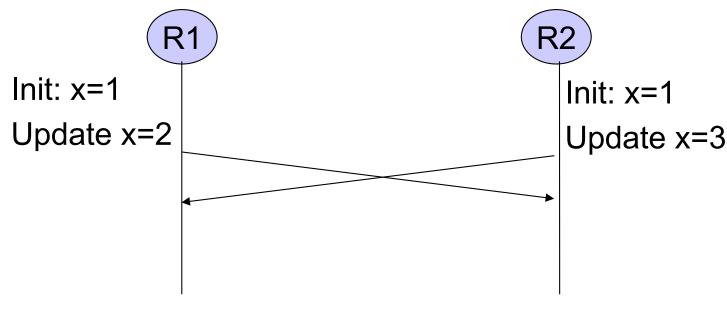
- Also called lazy replication
- Also called optimistic replication
- Main goals: availability and performance
- Approach
 - One replica updated by original transaction
 - Updates propagate asynchronously to other replicas

Asynchronous Master Replication

- One master holds primary copy
 - Transactions update primary copy
 - Master asynchronously propagates updates to replicas, which process them in same order (e.g. through log shipping)
 - Ensures single-copy serializability
- What happens when master/primary fails?
 - Can lose most recent transactions when primary fails!
 - After electing a new primary, secondaries must agree who is most up-to-date

Asynchronous Group Replication

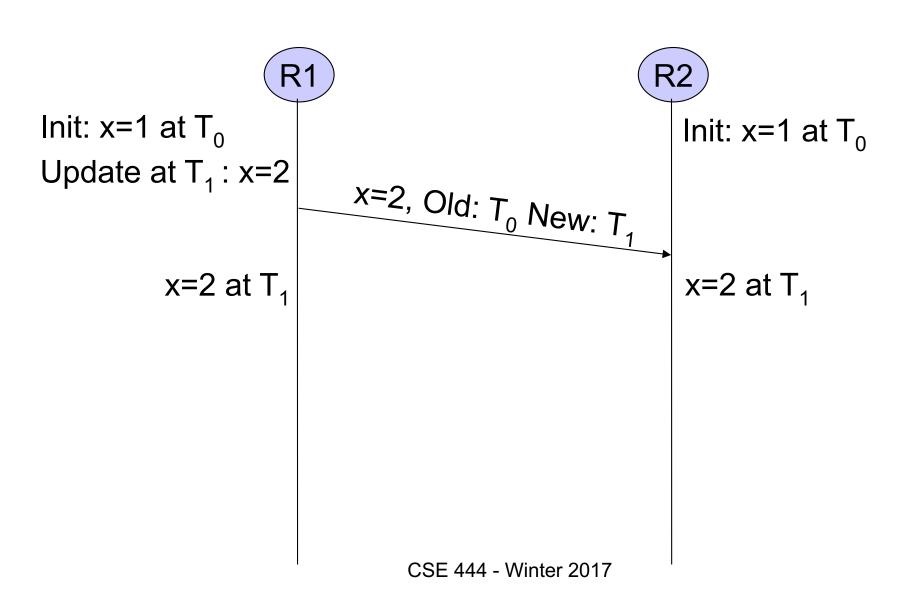
- Also called multi-master
- Best scheme for availability
- Cannot guarantee one-copy serializability!



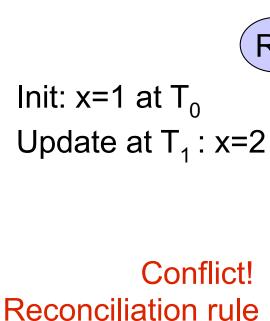
Asynchronous Group Replication

- Cannot guarantee one-copy serializability!
- Instead guarantee convergence
 - Db state does not reflect any serial execution
 - But all replicas have the same state
- Detect conflicts and reconcile replica states
- Different reconciliation techniques are possible
 - Manual
 - Most recent timestamp wins
 - Site A wins over site B
 - User-defined rules, etc.

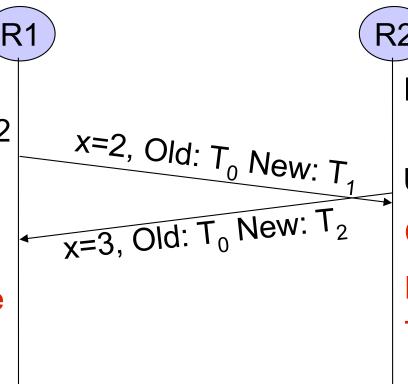
Detecting Conflicts Using Timestamps



Detecting Conflicts Using Timestamps



 $T_2 > T_1$, so x=3



Init: x=1 at T_0

Update at T_2 : x=3

Conflict!

Reconciliation rule

$$T_2 > T_1$$
, so x=3

Asynchronous Group Replication Properties

- Favours availability over consistency
 - Can read and update any replica
 - High runtime performance
- Weak consistency
 - Conflicts and reconciliation

Two-Tier Replication

- Benefits of lazy master and lazy group
- Each object has a master with primary copy
- When disconnected from master
 - Secondary can only run tentative transactions
- When reconnects to master
 - Master reprocesses all tentative transactions
 - Checks an acceptance criterion
 - If passes, we now have final commit order
 - Secondary undoes tentative and redoes committed

Conclusion

- Replication is a very important problem
 - Fault-tolerance (various forms of replication)
 - Caching (lazy master)
 - Warehousing (lazy master)
 - Mobility (two-tier techniques)
- Replication is complex, but basic techniques and trade-offs are very well known
 - Synchronous or asynchronous replication
 - Master or quorum