

CSE 444: Database Internals

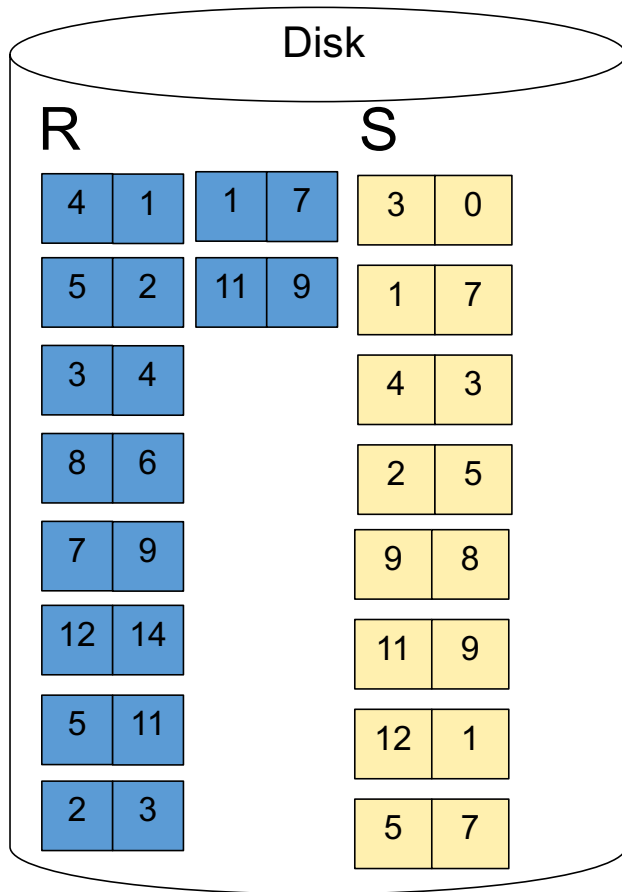
Lecture 8

Operator Algorithms (part 2)

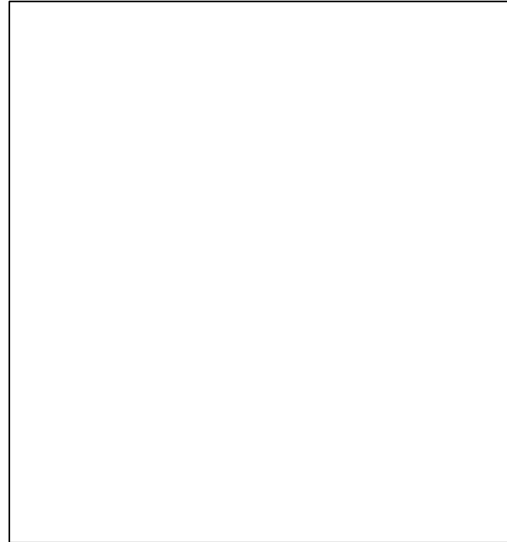
WHITE BOARD EXAMPLE

Relation R: 10 pages
Relation S: 8 pages
Memory: 5 pages

Each page holds 2 tuples
We only show value of join attribute for each tuple

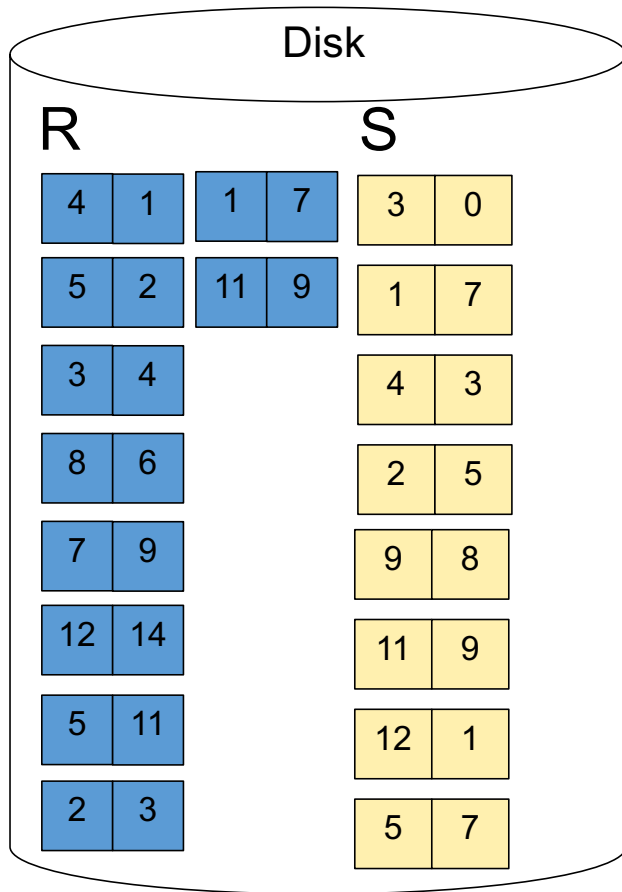


Memory M = 5 pages



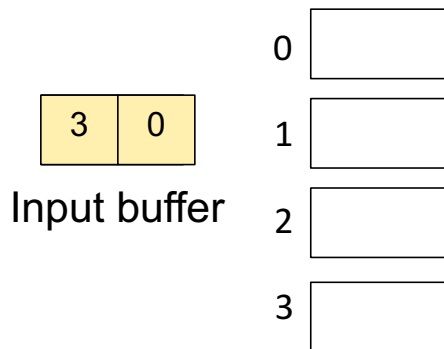
Grace Join
Also called
Partitioned Hash-Join Example

Step 1: Read relation S one page at a time and hash into M-1 (=4 buckets)

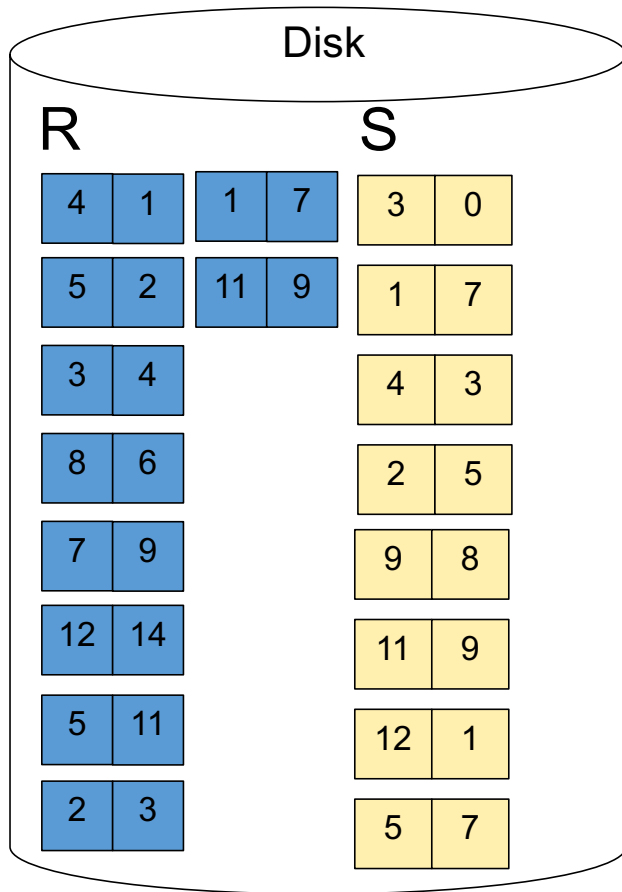


Memory M = 5 pages

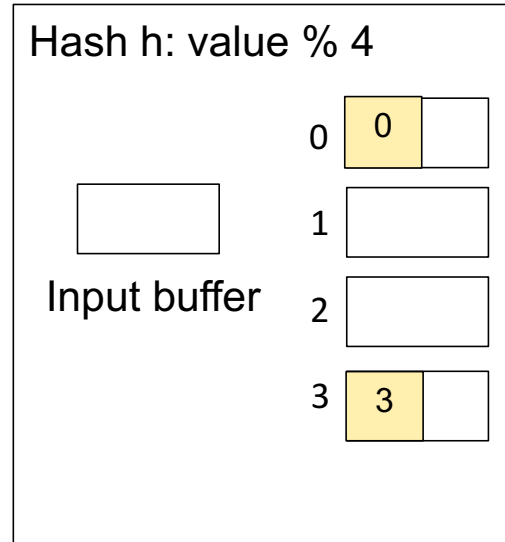
Hash h: value % 4



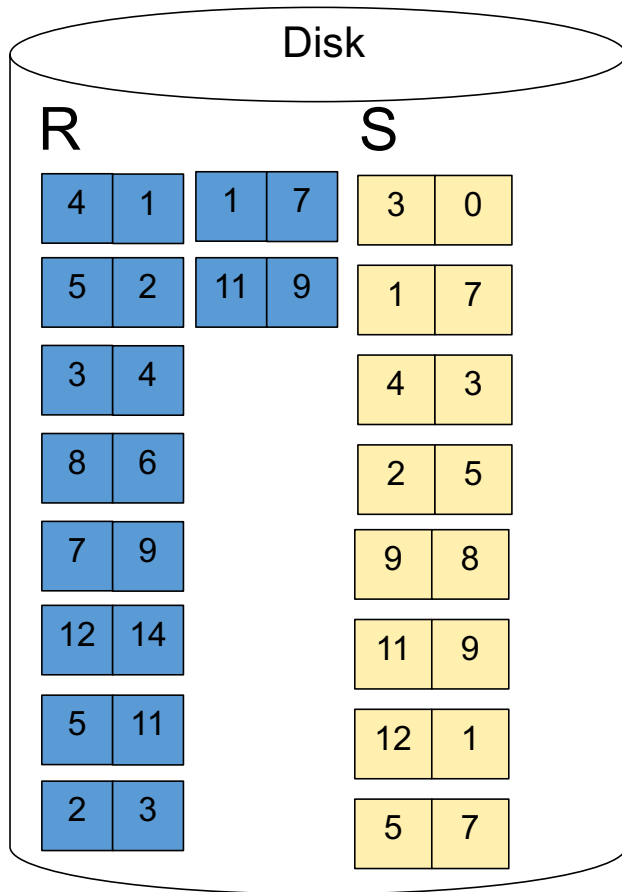
Step 1: Read relation S one page at a time and hash into the 4 buckets



Memory M = 5 pages

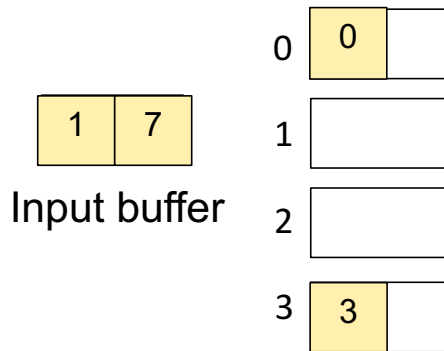


Step 1: Read relation S one page at a time and hash into the 4 buckets

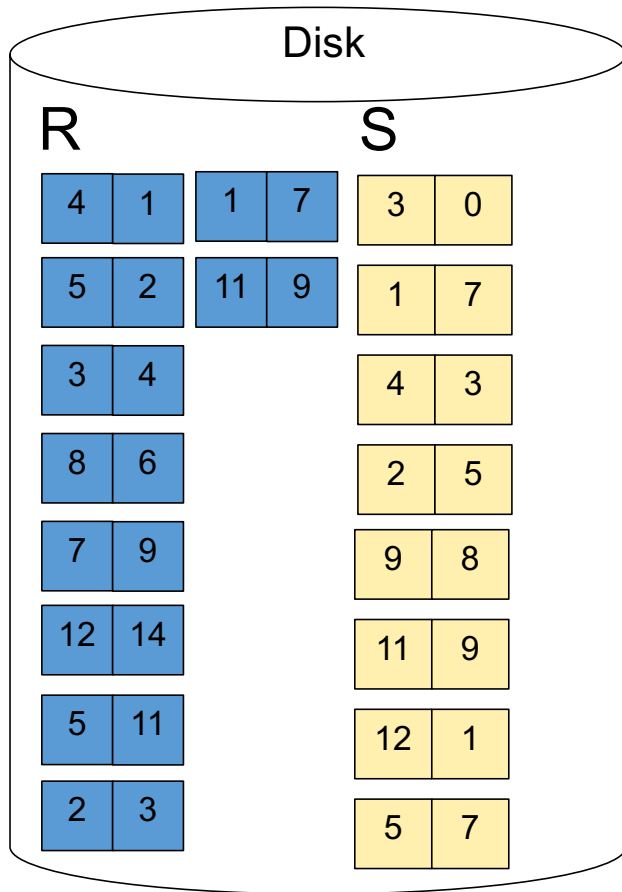


Memory M = 5 pages

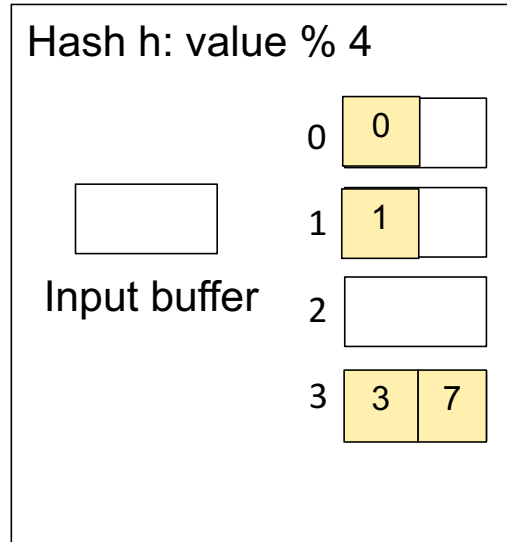
Hash h: value % 4



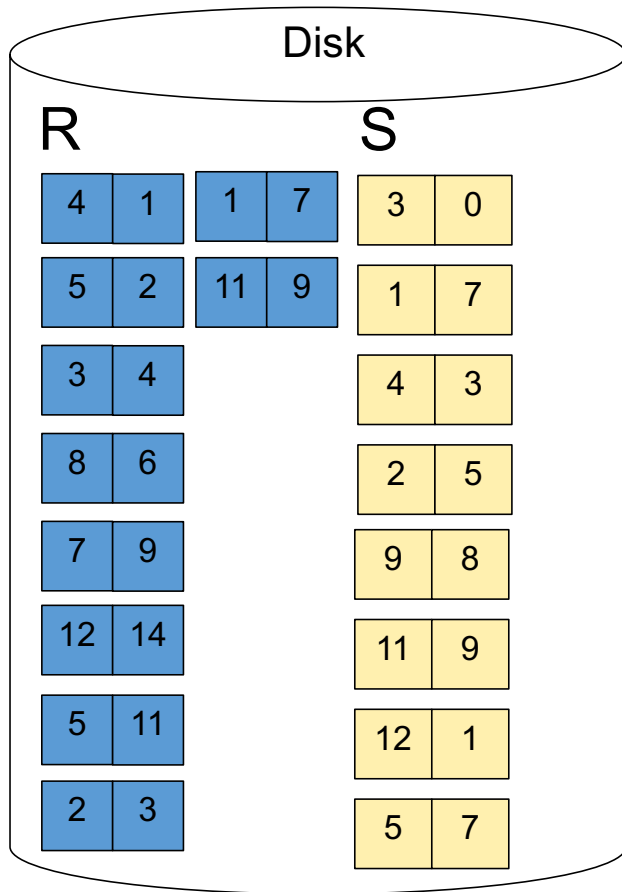
Step 1: Read relation S one page at a time and hash into the 4 buckets



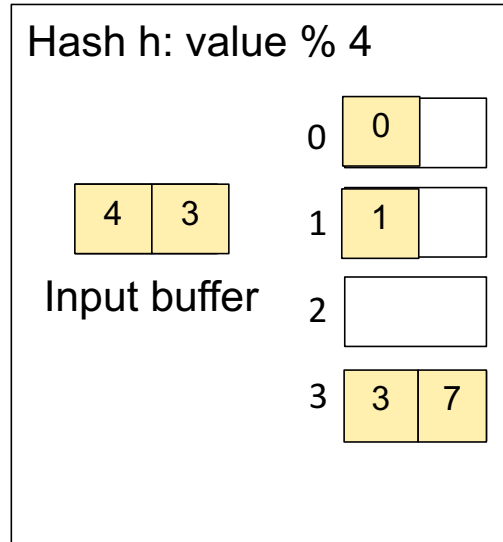
Memory M = 5 pages



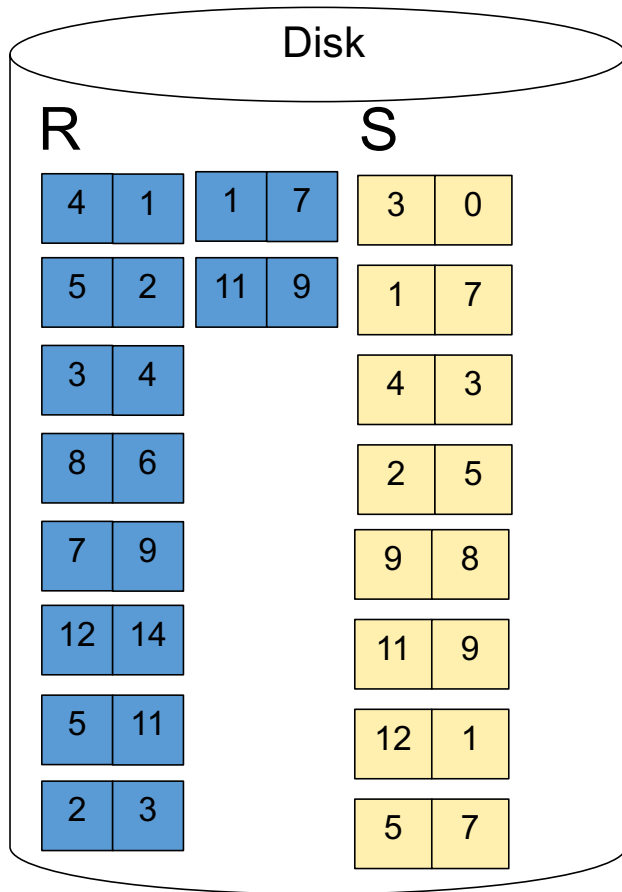
Step 1: Read relation S one page at a time and hash into the 4 buckets



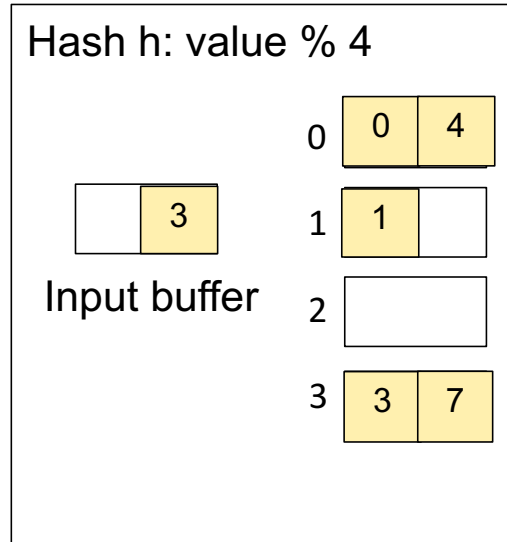
Memory M = 5 pages



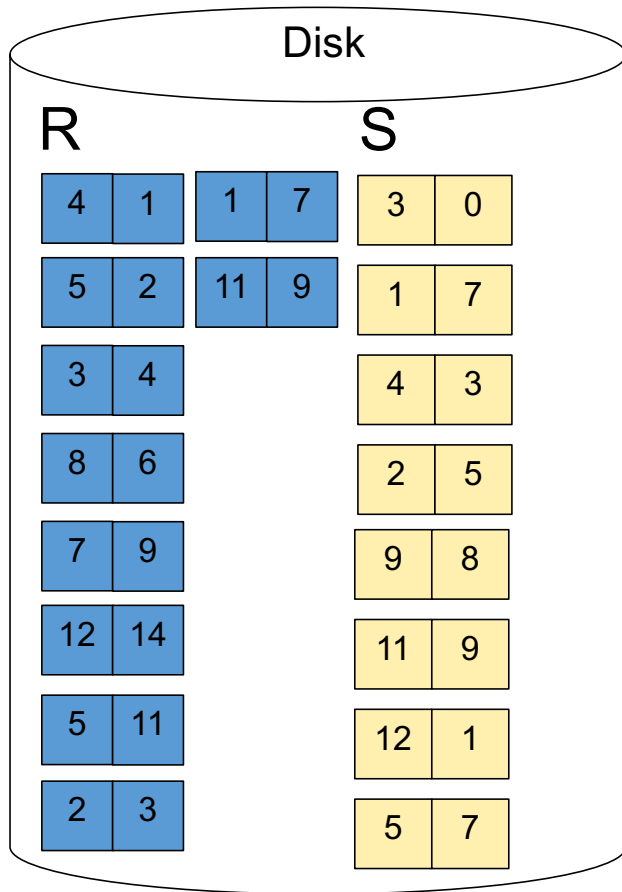
Step 1: Read relation S one page at a time and hash into the 4 buckets
 When a bucket fills up, flush it to disk



Memory M = 5 pages

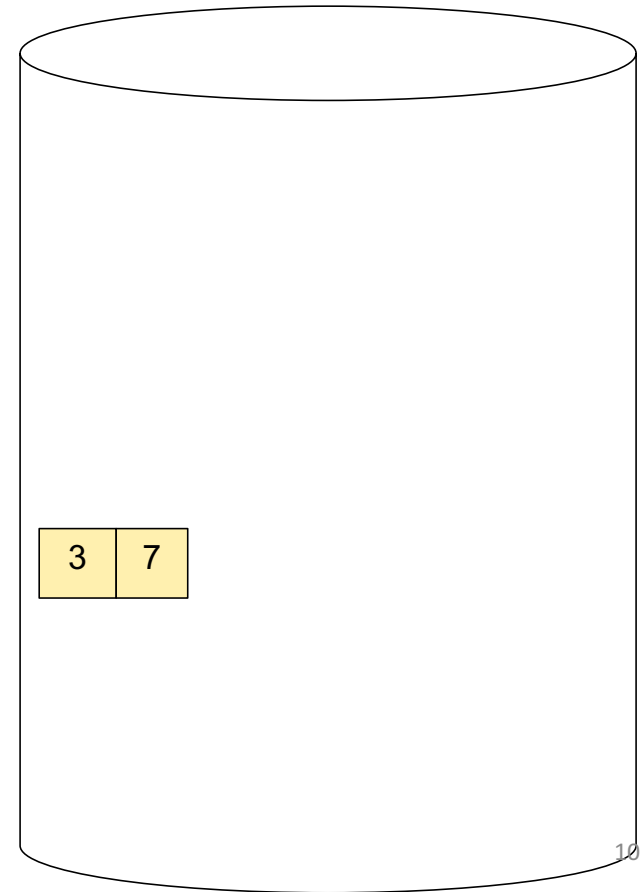
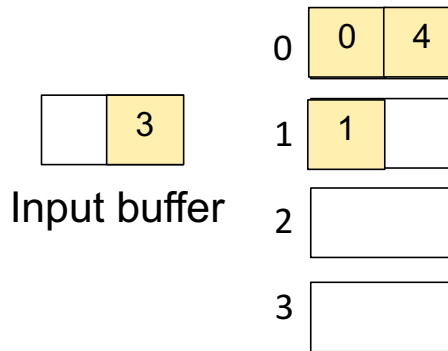


Step 1: Read relation S one page at a time and hash into the 4 buckets
 When a bucket fills up, flush it to disk

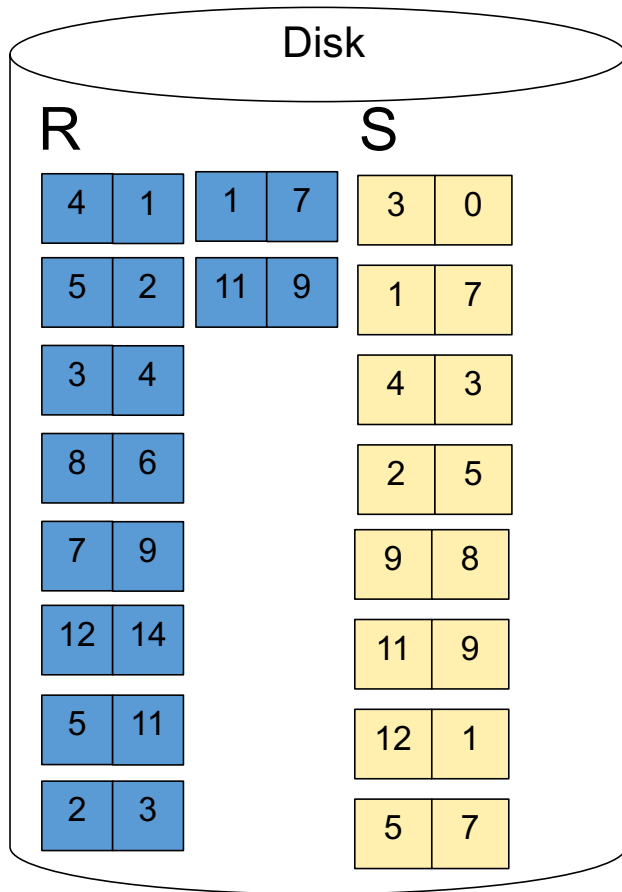


Memory M = 5 pages

Hash h: value % 4

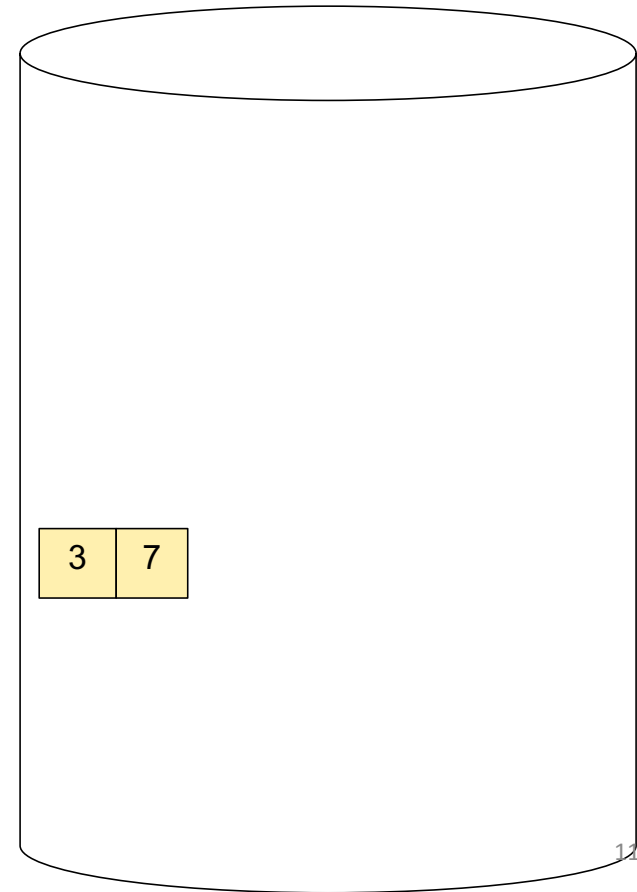
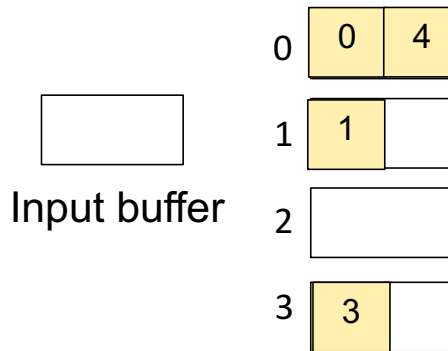


Step 1: Read relation S one page at a time and hash into the 4 buckets
 When a bucket fills up, flush it to disk

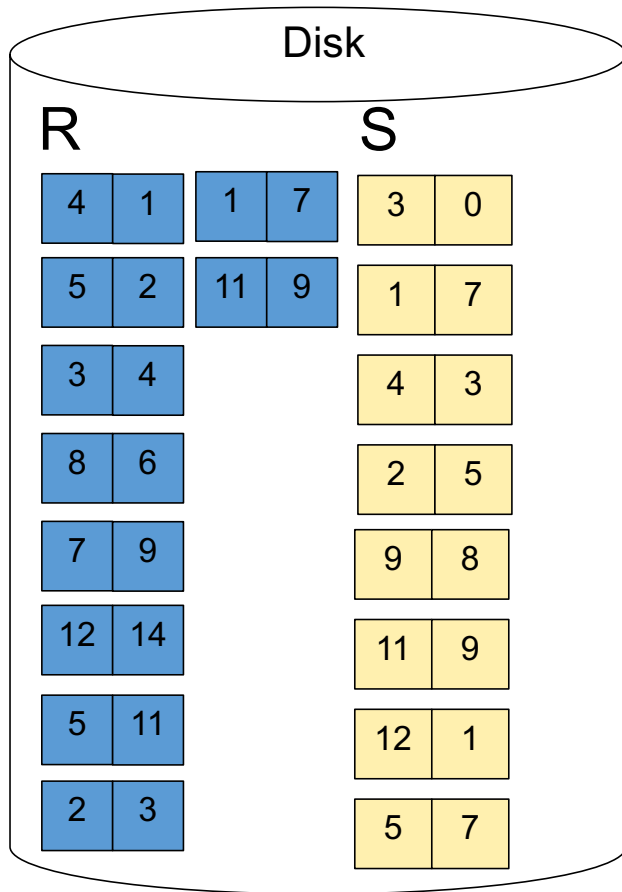


Memory M = 5 pages

Hash h: value % 4

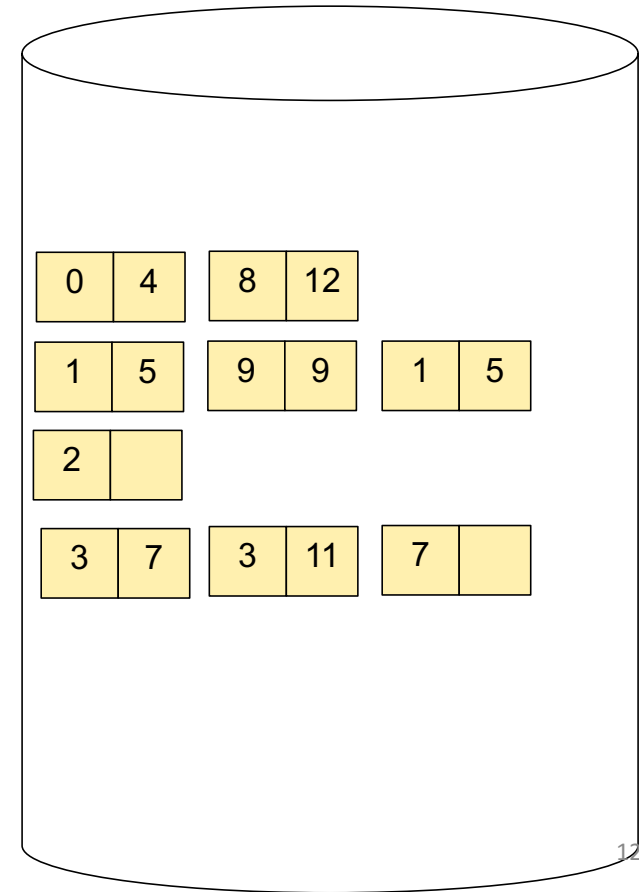
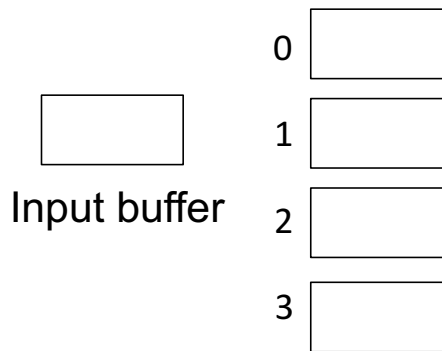


Step 1: Read relation S one page at a time and hash into the 4 buckets
 At the end, we get relation S back on disk split into 4 buckets

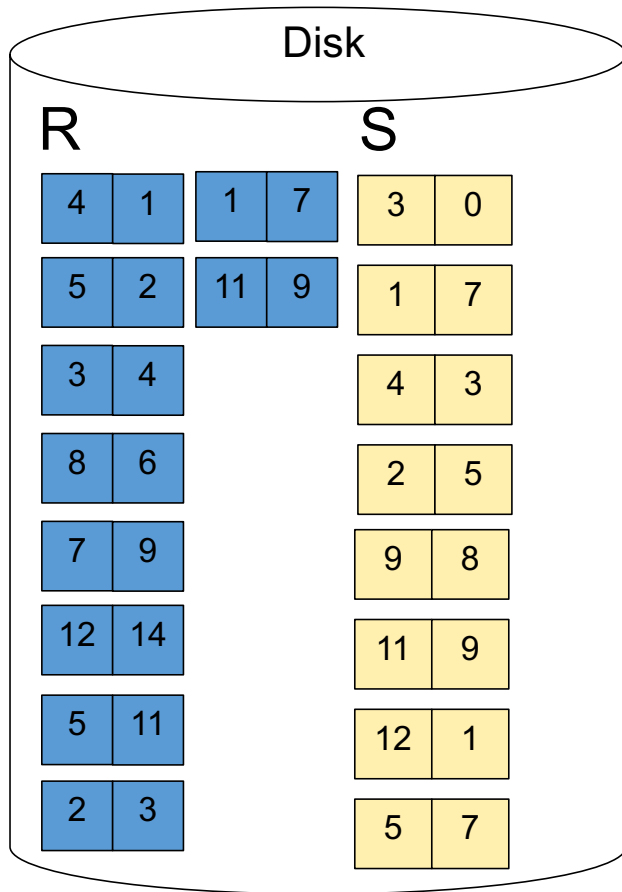


Memory M = 5 pages

Hash h: value % 4

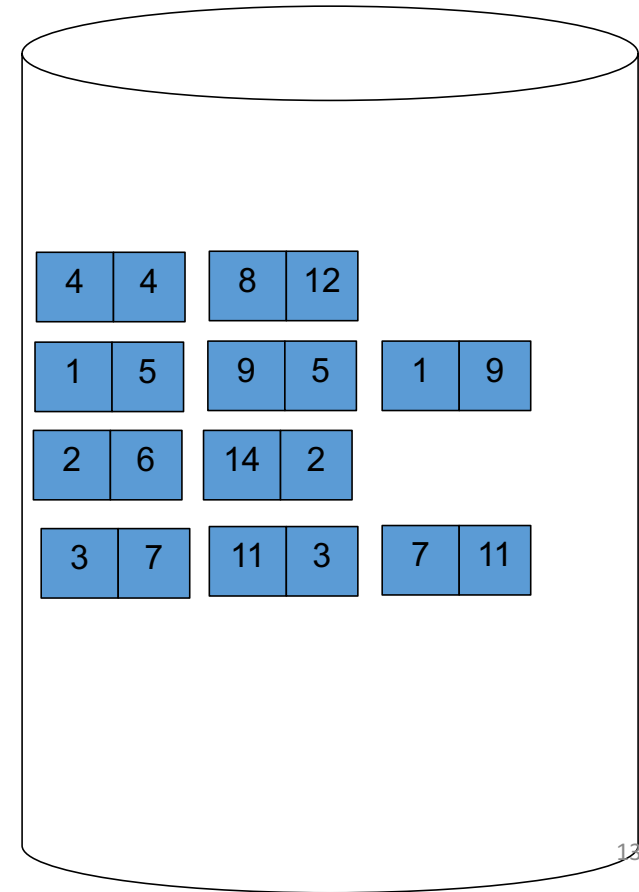
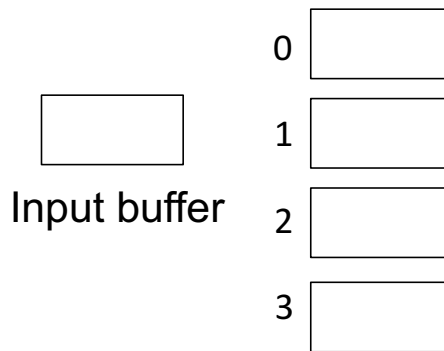


Step 2: Read relation R one page at a time and hash into same 4 buckets

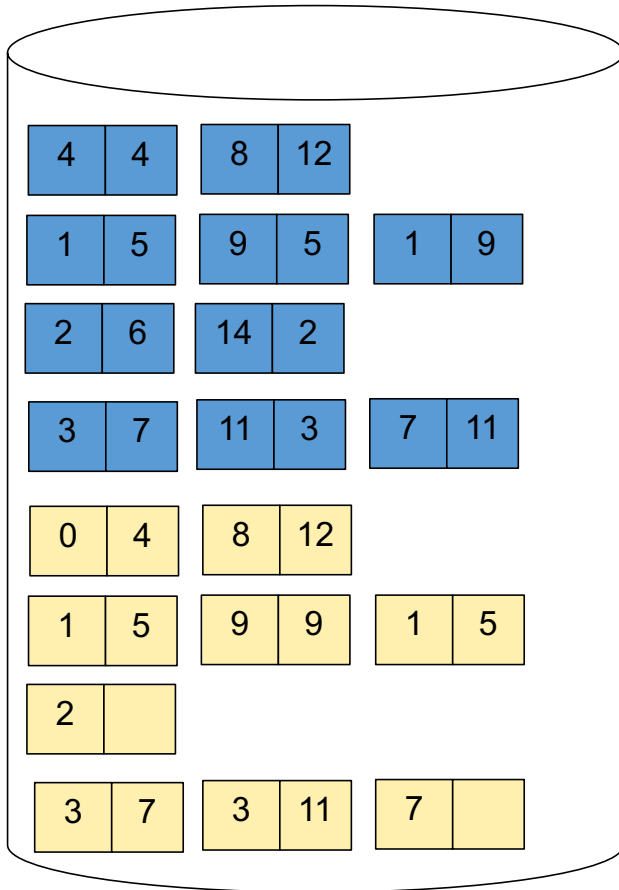


Memory M = 5 pages

Hash h: value % 4

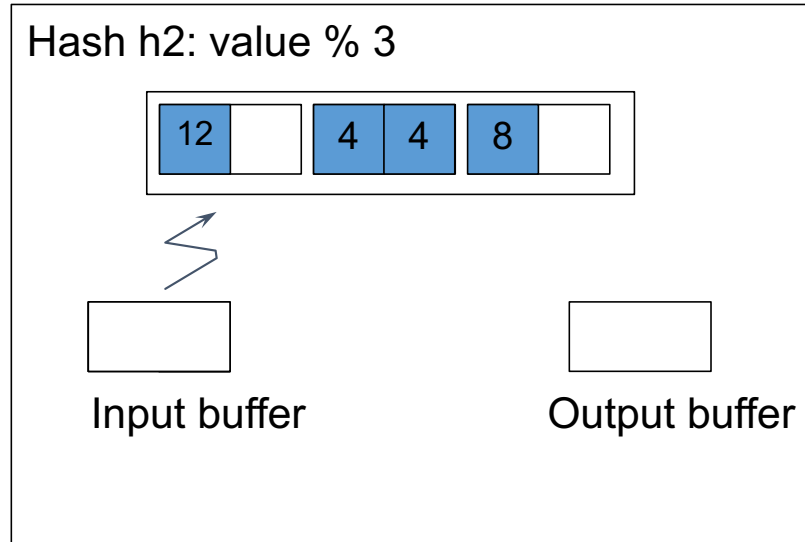


Step 3: Read one partition of R and create hash table in memory using a different hash function



Memory M = 5 pages

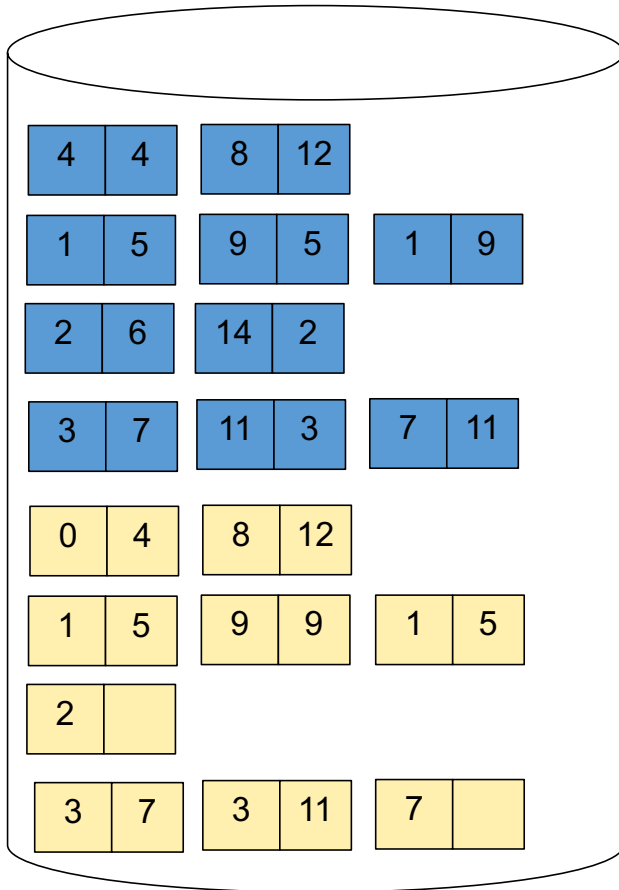
Hash h_2 : value % 3



Step 4: Scan matching partition of S and probe the hash table

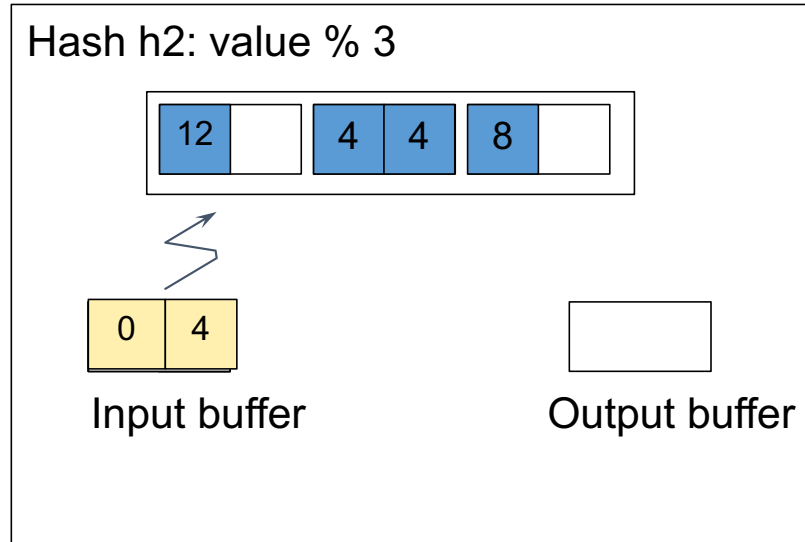
Step 5: Repeat for all the buckets

Total cost: $3B(R) + 3B(S)$



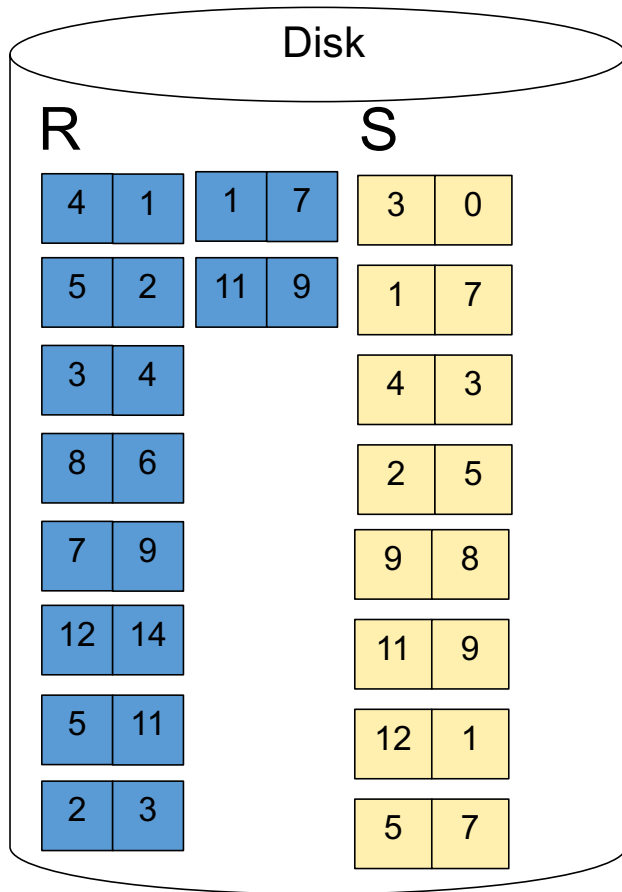
Memory $M = 5$ pages

Hash h_2 : $\text{value} \% 3$

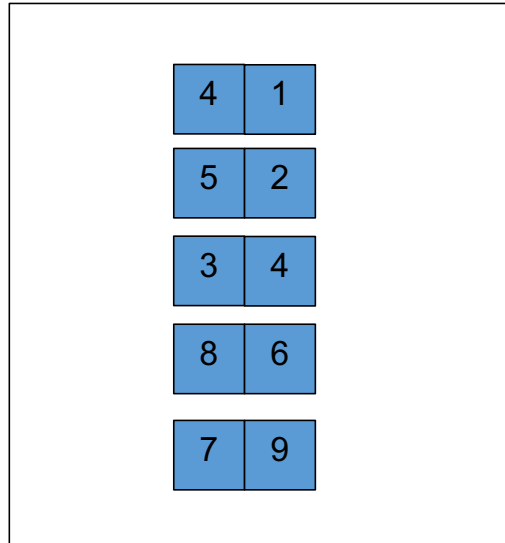


External Merge Join

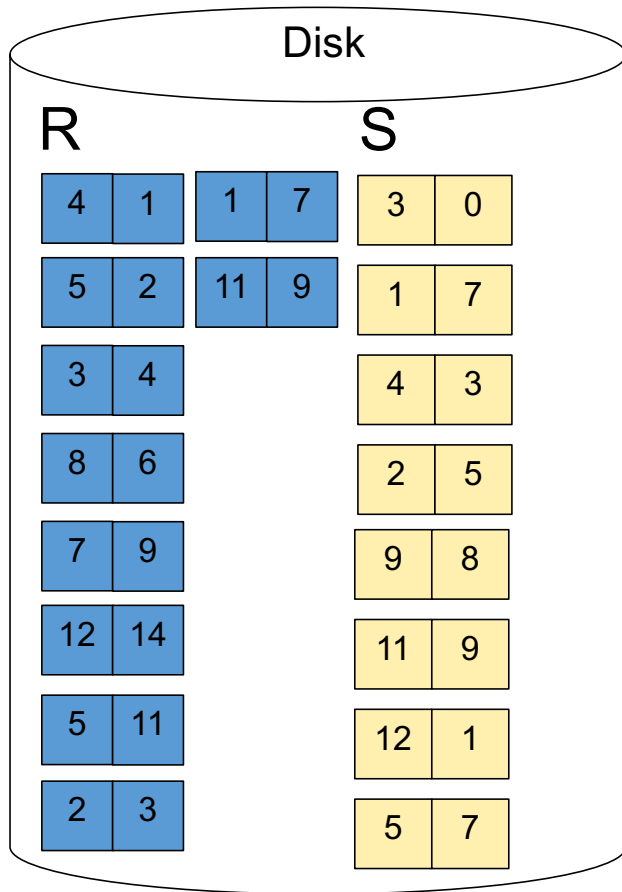
Step 1: Read M pages of R and sort in memory



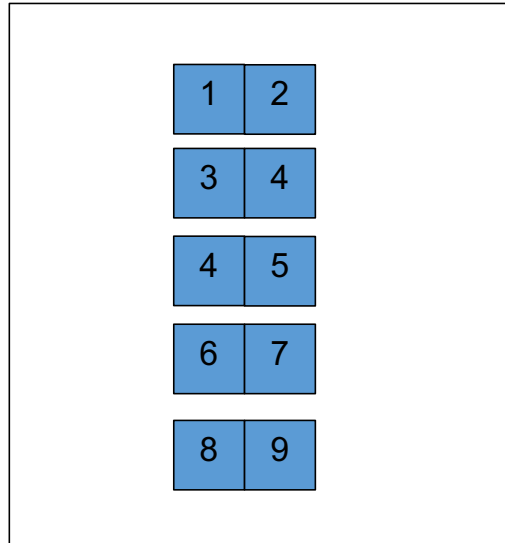
Memory M = 5 pages



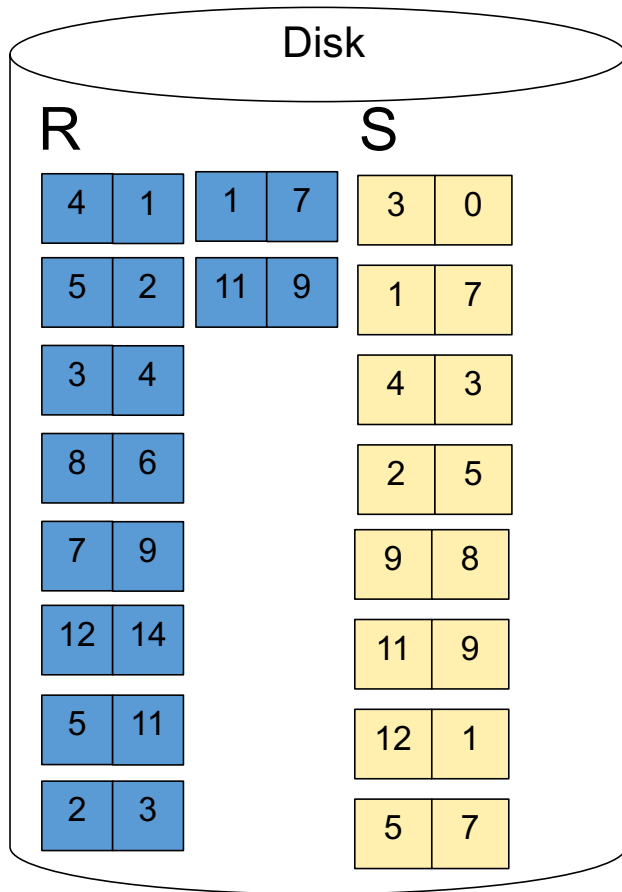
Step 1: Read M pages of R and sort in memory



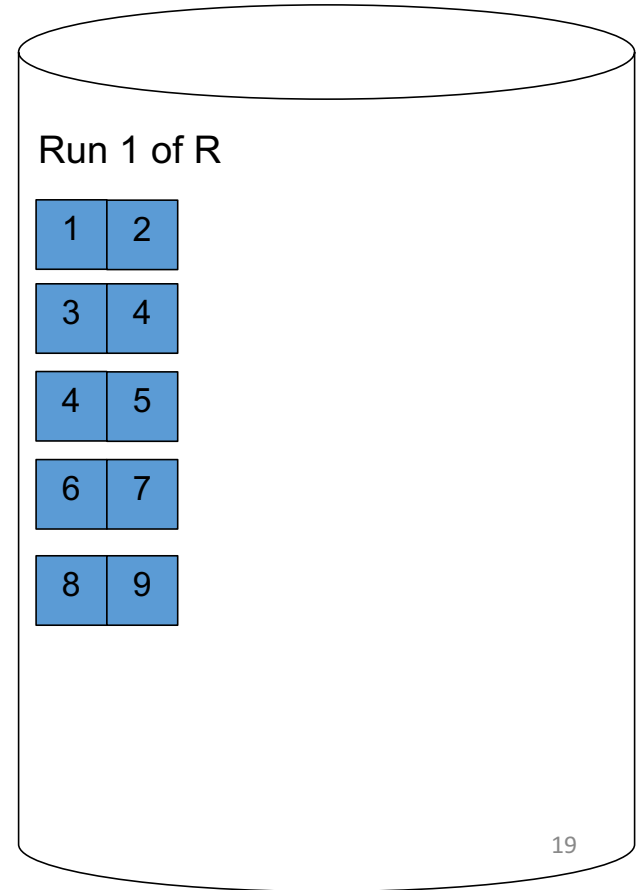
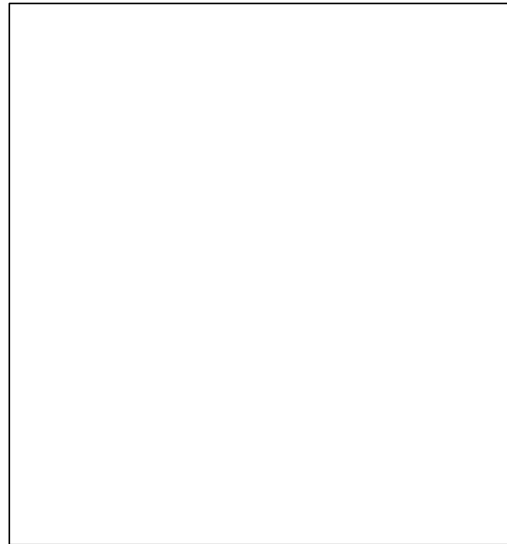
Memory M = 5 pages



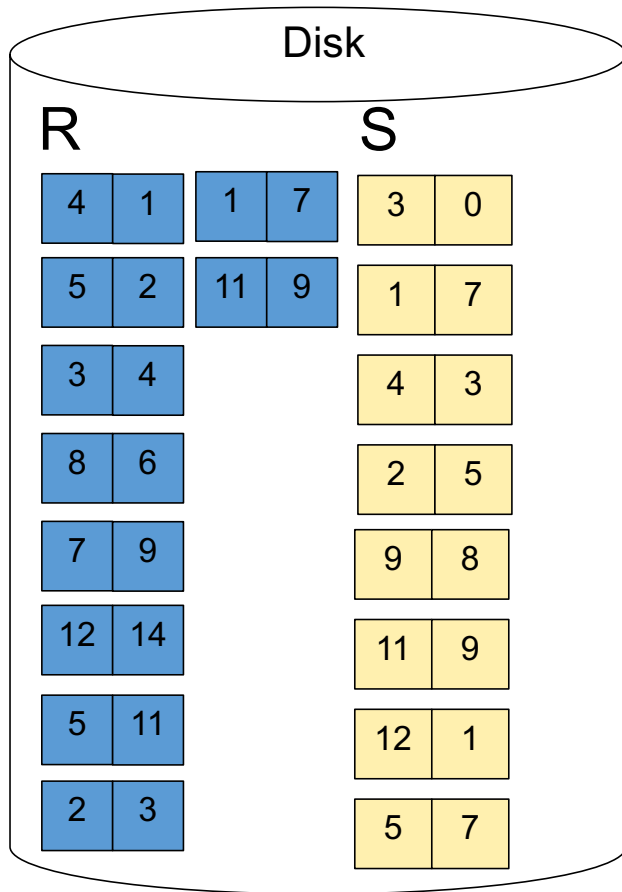
Step 1: Read M pages of R and sort in memory, then write to disk



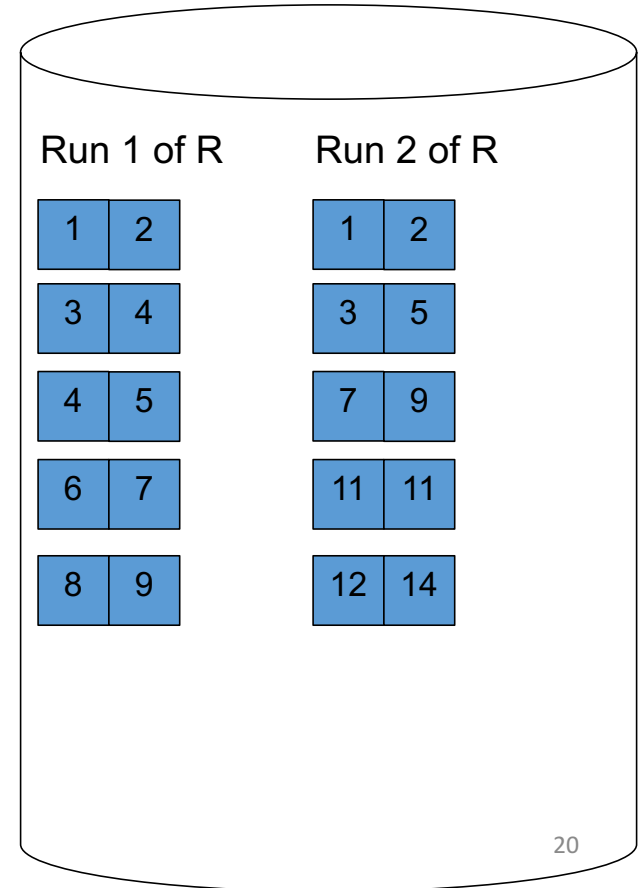
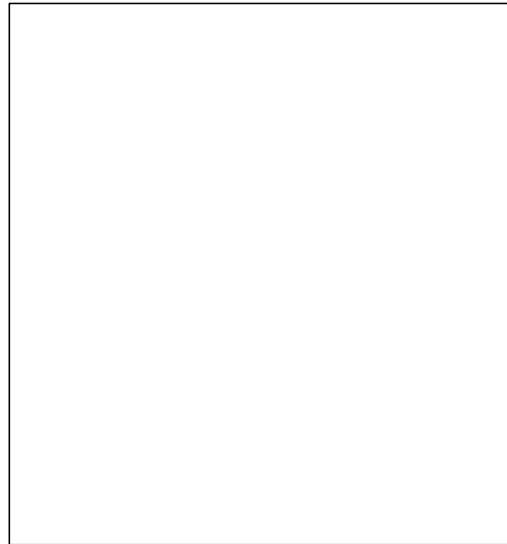
Memory M = 5 pages



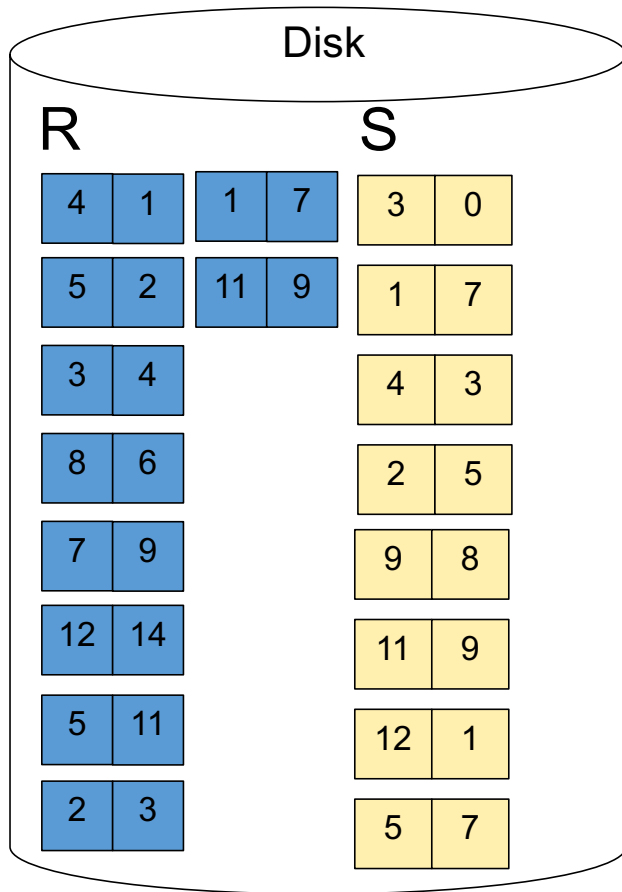
Step 2: Repeat for next M pages until all R is processed



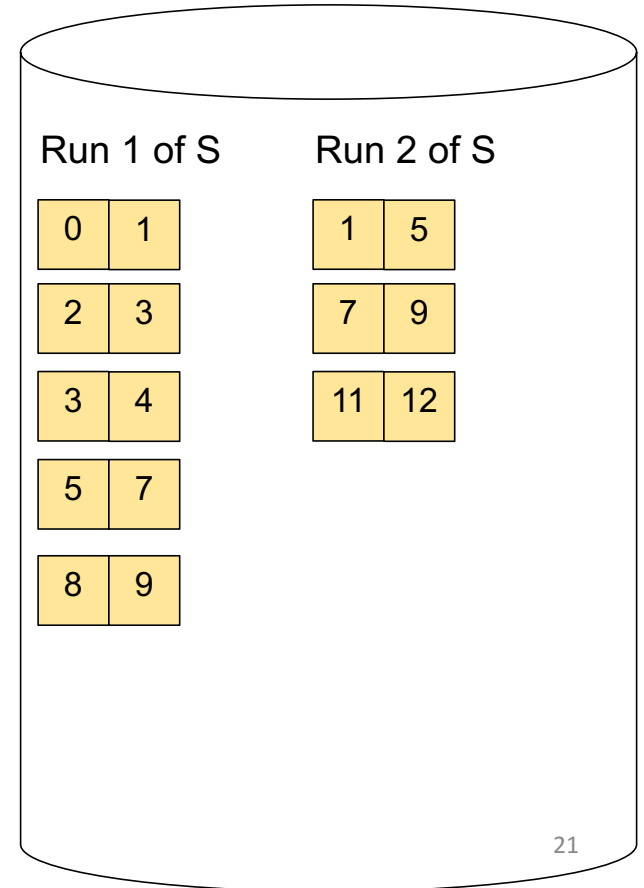
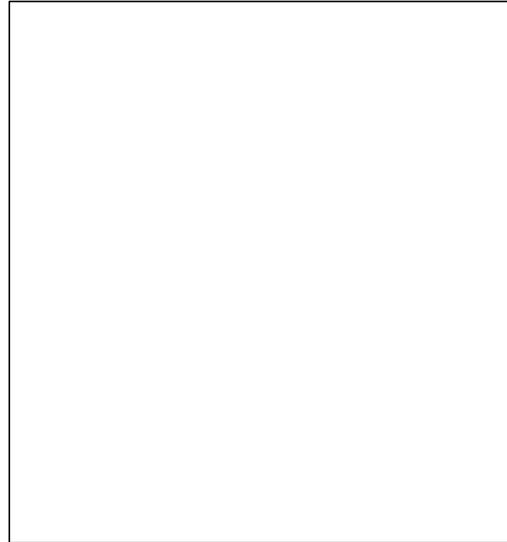
Memory M = 5 pages



Step 3: Do the same with S

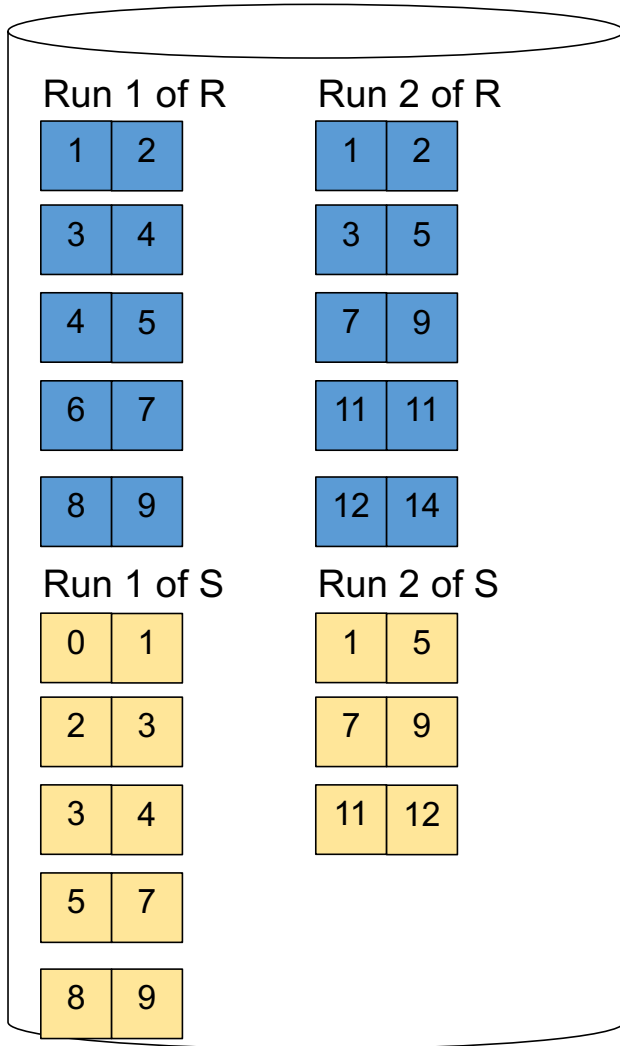


Memory M = 5 pages



Step 4: Join while merging sorted runs

Total cost: $3B(R) + 3B(S)$



Memory $M = 5$ pages

