

CSE 444: Database Internals

Lecture 25 Replication

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References

- Ullman Book Chapter 20.6
- **Database management systems.**
Ramakrishnan and Gehrke.
Third Ed. **Chapter 22.11**

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Outline

- Goals of replication
- Three types of replication
 - Eager replication
 - Lazy replication
 - Two-tier replication

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Goals of Replication

- Goal 1: availability
- Goal 2: performance



- But, it's easy to build a replicated system that reduces performance and availability

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Eager Replication

- Also called **synchronous replication**
- All updates are applied to all replicas (or to a majority) as part of a single transaction (need two phase commit)
- Main goal: as if there was only one copy
 - Maintain **consistency**
 - Maintain **one-copy serializability**
 - I.e., execution of transactions has same effect as an execution on a non-replicated db
- Transactions must acquire **global locks**

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Eager Master

- **One master for each object holds primary copy**
 - The "Master" is also called "Primary"
 - To update object, transaction must acquire a lock at the master
 - Lock at the master is global lock
- Master propagates updates to replicas synchronously
 - Updates propagate as part of the same distributed transaction
 - Need to run 2PC at the end
 - For example, using triggers

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Crash Failures

- What happens when a secondary crashes?
 - Nothing happens
 - When secondary recovers, it catches up
- What happens when the master/primary fails?
 - Blocking would hurt availability
 - Must choose a new primary: run election

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Network Failures

- Network failures can cause trouble...
 - Secondaries think that primary failed
 - Secondaries elect a new primary
 - But primary can still be running
 - Now have two primaries!

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Majority Consensus

- To avoid problem, only majority partition can continue processing at any time
- In general,
 - Whenever a replica fails or recovers...
 - a set of communicating replicas must determine...
 - whether they have a majority before they can continue

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Eager Group

- With n copies
 - Exclusive lock on x copies is global exclusive lock
 - Shared lock on s copies is global shared lock
 - Must have: $2x > n$ and $s + x > n$
- Majority locking
 - $s = x = \lceil (n+1)/2 \rceil$
 - No need to run any reconfiguration algorithms
- Read-locks-one, write-locks-all
 - $s=1$ and $x = n$, high read performance
 - Need to make sure algo runs on quorum of computers

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Eager Replication Properties

- Favours consistency over availability
 - Only majority partition can process requests
 - There appears to be a single copy of the db
- High runtime overhead
 - Must lock and update at least majority of replicas
 - Two-phase commit
 - Runs at pace of slowest replica in quorum
 - So overall system is now slower
 - Higher deadlock rate (transactions take longer)

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Lazy Replication

- Also called asynchronous replication
- Also called optimistic replication
- Main goals: availability and performance
- Approach
 - One replica updated by original transaction
 - Updates propagate asynchronously to other replicas

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Lazy Master

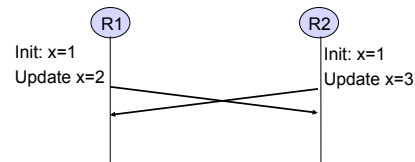
- One master holds primary copy
 - Transactions update primary copy
 - Master asynchronously propagates updates to replicas, which process them in same order (e.g. through log shipping)
 - Ensures single-copy serializability
- What happens when master/primary fails?
 - Can lose most recent transactions when primary fails!
 - After electing a new primary, secondaries must agree who is most up-to-date

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Lazy Group

- Also called **multi-master**
- Best scheme for availability
- **Cannot guarantee one-copy serializability!**



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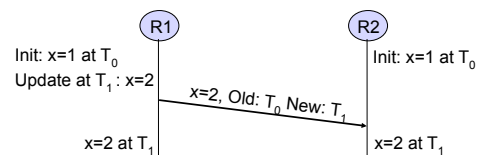
Lazy Group

- **Cannot guarantee one-copy serializability!**
- **Instead guarantee convergence**
 - Db state does not reflect any serial execution
 - But all replicas have the same state
- **Detect conflicts and reconcile replica states**
- Different reconciliation techniques are possible
 - Manual
 - Most recent timestamp wins
 - Site A wins over site B
 - User-defined rules, etc.

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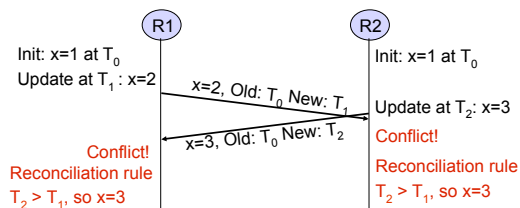
Detecting Conflicts Using Timestamps



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Detecting Conflicts Using Timestamps



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Lazy Group Replication Properties

- Favours **availability** over consistency
 - Can read and update any replica
 - High runtime performance
- **Weak consistency**
 - Conflicts and reconciliation

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Two-Tier Replication

- Benefits of lazy master and lazy group
- Each object has a master with primary copy
- When disconnected from master
 - Secondary can only run tentative transactions
- When reconnects to master
 - Master reprocesses all tentative transactions
 - Checks an acceptance criterion
 - If passes, we now have final commit order
 - Secondary undoes tentative and redoes committed

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Conclusion

- Replication is a very important problem
 - Fault-tolerance (various forms of replication)
 - Caching (lazy master)
 - Warehousing (lazy master)
 - Mobility (two-tier techniques)
- Replication is complex, but basic techniques and trade-offs are very well known
 - Eager or lazy replication
 - Master or quorum

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