

CSE 444: Database Internals

Lecture 11 Query Optimization (part 2)

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Reminders

- Homework 2 due tonight by 11pm in drop box
- Lab 2 is due on Monday by 11pm

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Query Optimizer Overview

- **Input:** A logical query plan
- **Output:** A good physical query plan
- **Basic query optimization algorithm**
 - Enumerate alternative plans (logical and physical)
 - Compute estimated cost of each plan
 - Compute number of I/Os
 - Optionally take into account other resources
 - Choose plan with lowest cost
 - This is called cost-based optimization

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The Three Parts of an Optimizer

- Cost estimation
 - Based on cardinality estimation
- Search space
 - Algebraic laws, restricted types of join trees
- Search algorithm
 - Will discuss next

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Search Algorithm

- Dynamic programming (in class)
 - Based on System R (aka Selinger) style optimizer[1979]
 - Limited to joins: *join reordering algorithm*
 - Bottom-up
- Rule-based algorithm (will not discuss)
 - Database of rules (=algebraic laws)
 - Usually: dynamic programming
 - Usually: top-down

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Dynamic Programming

Originally proposed in System R [1979]

- Only handles single block queries:

```
SELECT list
FROM R1, ..., Rn
WHERE cond1 AND cond2 AND ... AND condk
```

- Some heuristics for search space enumeration:
 - Selections down
 - Projections up
 - Avoid cartesian products

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SELECT list
FROM R1, ..., Rn
WHERE cond1 AND cond2 AND ... AND condn

Dynamic Programming

- For each subquery $Q \subseteq \{R_1, \dots, R_n\}$ compute the following:
 - $T(Q)$ = the estimated size of Q
 - $\text{Plan}(Q)$ = a best plan for Q
 - $\text{Cost}(Q)$ = the estimated cost of that plan

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SELECT list
FROM R1, ..., Rn
WHERE cond1 AND cond2 AND ... AND condn

Dynamic Programming

- Step 1:** For each $\{R_i\}$ do:
 - $T(\{R_i\}) = T(R_i)$
 - $\text{Plan}(\{R_i\})$ = access method for R_i
 - $\text{Cost}(\{R_i\})$ = cost of access method for R_i

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SELECT list
FROM R1, ..., Rn
WHERE cond1 AND cond2 AND ... AND condn

Dynamic Programming

- Step 2:** For each $Q \subseteq \{R_1, \dots, R_n\}$ of size k do:
 - $T(Q)$ = use estimator
 - Consider all partitions $Q = Q' \cup Q''$
compute $\text{cost}(\text{Plan}(Q') \bowtie \text{Plan}(Q''))$
 - $\text{Cost}(Q)$ = the smallest such cost
 - $\text{Plan}(Q)$ = the corresponding plan
- Note**
 - If we restrict to left-linear trees: Q'' = single relation
 - May want to avoid cartesian products

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SELECT list
FROM R1, ..., Rn
WHERE cond1 AND cond2 AND ... AND condn

Dynamic Programming

- Step 3:** Return $\text{Plan}(\{R_1, \dots, R_n\})$

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SELECT *
FROM R, S, T, U
WHERE cond1 AND cond2 AND ...

Example

- $R \bowtie S \bowtie T \bowtie U$
- Assumptions:

T(R) = 2000
T(S) = 5000
T(T) = 3000
T(U) = 1000
- Every join selectivity is 0.001

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T(R) = 2000
T(S) = 5000
T(T) = 3000
T(U) = 1000

Assume
B(..) = T(..)/10

Subquery	T	Plan	Cost
R	2000		
S	5000		
T	3000		
U	1000		
RS			
RT			
RU			
ST			
SU			
TU			
RST			
RSU			
RTU			
STU			
RSTU			

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T(R) = 2000
T(S) = 5000
T(T) = 3000
T(U) = 1000

Assume
 $B(.) = T(.)/10$

Subquery	T	Plan	Cost
R	2000		
S	5000		
T	3000		
U	1000		
RS	10000		
RT	6000		
RU	2000		
ST	15000		
SU	5000		
TU	3000		
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

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T(R) = 2000
T(S) = 5000
T(T) = 3000
T(U) = 1000

Assume
 $B(.) = T(.)/10$

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000		
T	3000		
U	1000		
RS	10000		
RT	6000		
RU	2000		
ST	15000		
SU	5000		
TU	3000		
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

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T(R) = 2000
T(S) = 5000
T(T) = 3000
T(U) = 1000

Assume
 $B(.) = T(.)/10$

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000		
U	1000		
RS	10000		
RT	6000		
RU	2000		
ST	15000		
SU	5000		
TU	3000		
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

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T(R) = 2000
T(S) = 5000
T(T) = 3000
T(U) = 1000

Assume
 $B(.) = T(.)/10$

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000	Table scan	300
U	1000	Unclustered index scan U.F	1000
RS	10000		
RT	6000		
RU	2000		
ST	15000		
SU	5000		
TU	3000		
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

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T(R) = 2000
T(S) = 5000
T(T) = 3000
T(U) = 1000

Assume
 $B(.) = T(.)/10$

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000	Table scan	300
U	1000	Unclustered index scan U.F	1000
RS	10000	R $\bowtie$ S nested loop join	...
RT	6000		
RU	2000		
ST	15000		
SU	5000		
TU	3000		
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

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T(R) = 2000
T(S) = 5000
T(T) = 3000
T(U) = 1000

Assume
 $B(.) = T(.)/10$

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000	Table scan	300
U	1000	Unclustered index scan U.F	1000
RS	10000	R $\bowtie$ S nested loop join	...
RT	6000	R $\bowtie$ T index join	
RU	2000		
ST	15000		
SU	5000		
TU	3000		
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

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T(R) = 2000
T(S) = 5000
T(T) = 3000
T(U) = 1000

Assume
 $B(.) = T(.)/10$

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000	Table scan	300
U	1000	Unclustered index scan U.F	1000
RS	10000	R $\bowtie$ S nested loop join	...
RT	6000	R $\bowtie$ T index join	...
RU	2000	R $\bowtie$ U index join	...
ST	15000	S $\bowtie$ T hash join	...
SU	5000	...	...
TU	3000	...	...
RST	30000		
RSU	10000		
RTU	6000		
STU	15000		
RSTU	30000		

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T(R) = 2000
T(S) = 5000
T(T) = 3000
T(U) = 1000

Assume
 $B(.) = T(.)/10$

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000	Table scan	300
U	1000	Unclustered index scan U.F	1000
RS	10000	R $\bowtie$ S nested loop join	...
RT	6000	R $\bowtie$ T index join	...
RU	2000	R $\bowtie$ U index join	...
ST	15000	S $\bowtie$ T hash join	...
SU	5000	...	...
TU	3000	...	...
RST	30000	(RT) $\bowtie$ S hash join	...
RSU	10000	(SU) $\bowtie$ R merge join	...
RTU	6000	...	...
STU	15000	...	...
RSTU	30000	...	...

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T(R) = 2000
T(S) = 5000
T(T) = 3000
T(U) = 1000

Assume
 $B(.) = T(.)/10$

Subquery	T	Plan	Cost
R	2000	Clustered index scan R.A	200
S	5000	Table scan	500
T	3000	Table scan	300
U	1000	Unclustered index scan U.F	1000
RS	10000	R $\bowtie$ S nested loop join	...
RT	6000	R $\bowtie$ T index join	...
RU	2000	R $\bowtie$ U index join	...
ST	15000	S $\bowtie$ T hash join	...
SU	5000	...	...
TU	3000	...	...
RST	30000	(RT) $\bowtie$ S hash join	...
RSU	10000	(SU) $\bowtie$ R merge join	...
RTU	6000	...	...
STU	15000	...	...
RSTU	30000	(RT) $\bowtie$ (SU) hash join	...

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Discussion

- Need to consider both $R \bowtie S$ and $S \bowtie R$
 - Because the cost may be different!
- When computing the cheapest plan for $(Q) \bowtie R$, we may consider new access methods for R, e.g. an index look-up that makes sense only in the context of the join

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Discussion

Given a query with n relations R1, ..., Rn

- How many entries do we have in the dynamic programming table?
- For each entry, how many alternative plans do we need to inspect?

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SELECT list
FROM R1, ..., Rn
WHERE cond1 AND cond2 AND ... AND condk

Discussion

Given a query with n relations R1, ..., Rn

- How many entries do we have in the dynamic programming table?
 - A: $2^n - 1$
- For each entry, how many alternative plans do we need to inspect?
 - A: for each entry with k tables, examine $2^k - 2$ plans

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SELECT list
FROM R1, ..., Rn
WHERE cond1 AND cond2 AND ... AND condk

Reducing the Search Space

- Left-linear trees
- No cartesian products

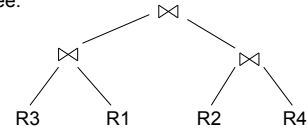
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Join Trees

- $R1 \bowtie R2 \bowtie \dots \bowtie Rn$

- Join tree:



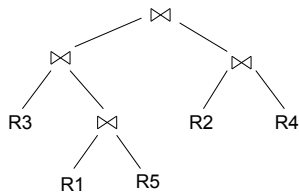
- A plan = a join tree
- A partial plan = a subtree of a join tree

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Types of Join Trees

- Bushy:

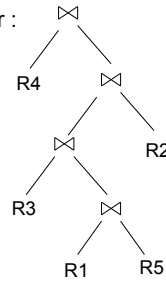


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Types of Join Trees

- Linear :

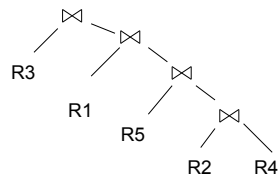


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Types of Join Trees

- Right deep:



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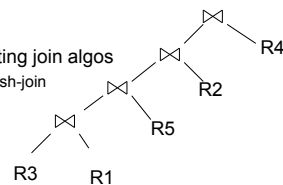
Types of Join Trees

- Left deep:

- Work well with existing join algos
 - Nested-loop and hash-join

- Facilitate pipelining

- Dynamic programming can be used with all trees



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Selinger Algorithm

Selinger enumeration algorithm considers

- Different logical and physical plans *at the same time*
- Cost of a plan is IO + CPU
- Concept of *interesting order* during plan enumeration
 - Same order as that requested by ORDER BY or GROUP BY
 - Attributes that appear in eq-join predicates

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More about the Selinger Algorithm

- Step 1: Enumerate all access paths for a single relation
 - File scan or index scan
 - Keep the cheapest for each *interesting order*
- Step 2: Consider all ways to join two relations
 - Use result from step 1 as the outer relation
 - Consider every other possible relation as inner relation
 - Estimate cost when using sort-merge or nested-loop join
 - Keep the cheapest for each *interesting order*
- Steps 3 and later: Repeat for three relations, etc.

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Example on the Board

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