#### CSE 444: Database Internals

Lecture 24 Replication

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## References

- · Book Chapter 20.6
- · Database management systems.

Ramakrishnan and Gehrke.

Third Ed. Chapter 22.11 (more info than our main book)

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requests

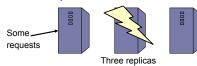
## Outline

- · Goals of replication
- · Three types of replication
  - Eager replication
  - Lazy replication
  - Two-tier replication

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# Goals of Replication

- · Goal 1: availability
- Goal 2: performance



• But, it's easy to build a replicated system that reduces performance and availability

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# **Eager Replication**

- · Also called synchronous replication
- All updates are applied to all replicas (or to a majority) as part of a single transaction (need two phase commit)
- · Main goal: as if there was only one copy
  - Maintain consistency
  - Maintain one-copy serializability
  - I.e., execution of transactions has same effect as an execution on a non-replicated db
- Transactions must acquire global locks

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## Eager Master

- One master for each object holds primary copy
  - The "Master" is also called "Primary"
  - To update object, transaction must acquire a lock at the master
  - Lock at the master is global lock

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## Crash Failures

- What happens when a secondary crashes?
  - Nothing happens
  - When secondary recovers, it catches up
- · What happens when the master/primary fails?
  - Blocking would hurt availability
  - Must chose a new primary: run election

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## **Network Failures**

- · Network failures can cause trouble...
  - Secondaries think that primary failed
  - Secondaries elect a new primary
  - But primary can still be running
  - Now have two primaries!

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## **Majority Consensus**

- · To avoid problem, only majority partition can continue processing at any time
- · In general,
  - Whenever a replica fails or recovers...
  - a set of communicating replicas must determine...
  - whether they have a majority before they can continue

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## **Eager Group**

- · With n copies
  - Exclusive lock on x copies is global exclusive lock
  - Shared lock on s copies is global shared lock
  - Must have: 2x > n and s + x > n
- · Majority locking
  - s = x = [(n+1)/2]
  - No need to run any reconfiguration algorithms
- · Read-locks-one, write-locks-all
  - s=1 and x = n, high read performance
  - Need to make sure algo runs on quorum of computers

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## **Eager Replication Properties**

- · Favours consistency over availability
  - Only majority partition can process requests
  - There appears to be a single copy of the db
- · High runtime overhead
  - Must lock and update at least majority of replicas
  - Two-phase commit
  - Runs at pace of slowest replica in quorum
  - So overall system is now slower
  - Higher deadlock rate (transactions take longer)

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## Lazy Replication

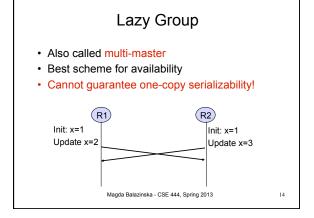
- · Also called asynchronous replication
- · Also called optimistic replication
- · Main goals: availability and performance
- · Approach
  - One replica updated by original transaction
  - Updates propagate asynchronously to other replicas

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# Lazy Master

- · One master holds primary copy
  - Transactions update primary copy
  - Master asynchronously propagates updates to replicas, which process them in same order
  - Ensures single-copy serializability
- · What happens when master/primary fails?
  - Can lose most recent transactions when primary fails!
  - After electing a new primary, secondaries must agree who is most up-to-date

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# Lazy Group

- · Cannot guarantee one-copy serializability!
- · Instead guarantee convergence
  - Db state does not reflect any serial execution
  - But all replicas have the same state
- · Detect conflicts and reconcile replica states
- · Different reconciliation techniques are possible
  - Manual
  - Most recent timestamp wins
  - Site A wins over site B
  - User-defined rules, etc.

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Detecting Conflicts
Using Timestamps

R1

Init: x=1 at T<sub>0</sub>
Update at T<sub>1</sub>: x=2

x=2, Old: T<sub>0</sub> New: T<sub>1</sub>

x=2 at T<sub>1</sub>

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#### **Detecting Conflicts Using Timestamps** (R1) Init: x=1 at T<sub>0</sub> Init: x=1 at T<sub>o</sub> Update at T<sub>1</sub>: x=2 x=2, Old: T<sub>0</sub> New: T Update at T<sub>2</sub>: x=3 Conflict! Conflict! Reconciliation rule Reconciliation rule $T_2 > T_1$ , so x=3 $T_2 > T_1$ , so x=3 Magda Balazinska - CSE 444, Spring 2013 17

# Lazy Group Replication Properties

- · Favours availability over consistency
  - Can read and update any replica
  - High runtime performance
- · Weak consistency
  - Conflicts and reconciliation

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# **Two-Tier Replication**

- · Benefits of lazy master and lazy group
- Each object has a master with primary copy
- When disconnected from master
  - Secondary can only run tentative transactions
- · When reconnects to master
  - Master reprocesses all tentative transactions
  - Checks an acceptance criterion
  - If passes, we now have final commit order
  - Secondary undoes tentative and redoes committed

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## Conclusion

- · Replication is a very important problem
  - Fault-tolerance (various forms of replication)
  - Caching (lazy master)
  - Warehousing (lazy master)
  - Mobility (two-tier techniques)
- Replication is complex, but basic techniques and trade-offs are **very well known** 
  - Eager or lazy replication
  - Master or no master
  - For eager replication: use quorum

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