

CSE 444 Practice Problems

Parallel DBMSs and MapReduce

1. Parallel Data Processing Algorithms

- (a) Describe how to compute the equi-join of two relations R and S in parallel. Describe how to compute a group by aggregation operation in parallel.

Solution:

See lecture notes.

2. System Comparison

- (a) In a parallel DBMS, why is it difficult to achieve linear speedup and linear scaleup?

Solution:

There are three key reasons why linear speedup and linear scaleup are difficult to achieve:

- i. Startup cost: The latency involved in starting an operation on many nodes may dominate the computation time.
- ii. Interference: Each new process competes for shared resources with the other processes (e.g., network bandwidth). This resource contention can limit the performance gains of adding more processes.
- iii. Skew: The time to complete a job is the time that the slowest partition takes to complete its job. When the variance dominates the mean, increased parallelism improves elapsed time only slightly.

- (b) List two features common to a traditional DBMS and MapReduce.

Clarification: Here, “traditional DBMS” means a traditional *parallel DBMS*.

Solution:

Many answers are possible including:

- i. Horizontal data and operator partitioning.
- ii. Distribution independence: applications need not know that the data is distributed.

- (c) List two features that are different between the two types of systems: i.e., features that are present in one system but not in the other. For example, you can give one feature present in MapReduce but absent in a parallel DBMS and one feature present in a parallel DBMS but missing from MapReduce (or any other combination).

Solution:

Again, different answers were possible including:

- i. A DBMS offers updates, transactions, indexing, and pipelined parallelism.
- ii. MapReduce has intra-query fault-tolerance and handles stragglers better.