

# Lecture 21: Query Optimization (3)

Friday, May 21, 2010

# Announcement

- No Homework 4
- **BUT PLEASE:**
  - Study the remaining material
  - Do good job on Project 4

# Outline

- Search space
- Algorithms for enumerating query plans
- Estimating the cost of a query plan

# Computing the Cost of a Plan

- Collect statistical summaries of stored data
- Estimate size in a bottom-up fashion
- Estimate cost by using the estimated size

# Statistics on Base Data

- **Collected information for each relation**
  - Number of tuples (cardinality)
  - Indexes, number of keys in the index
  - Number of physical pages, clustering info
  - Statistical information on attributes
    - Min value, max value, number distinct values
    - Histograms
  - Correlations between columns (hard)
- **Collection approach: periodic, using sampling**

# Size Estimation Problem

```
S = SELECT list  
    FROM R1, ..., Rn  
    WHERE cond1 AND cond2 AND ... AND condk
```

Given  $T(R_1), T(R_2), \dots, T(R_n)$   
Estimate  $T(S)$

How can we do this ? Note: doesn't have to be exact.

# Size Estimation Problem

```
S = SELECT list  
    FROM R1, ..., Rn  
    WHERE cond1 AND cond2 AND ... AND condk
```

Remark:  $T(S) \leq T(R1) \times T(R2) \times \dots \times T(Rn)$

# Selectivity Factor

- Each condition *cond* reduces the size by some factor called *selectivity factor*
- Assuming independence, multiply the selectivity factors



# Example

R(A,B)  
S(B,C)  
T(C,D)

```
SELECT *  
FROM R, S, T  
WHERE R.B=S.B and S.C=T.C and R.A<40
```

$T(R) = 30k$ ,  $T(S) = 200k$ ,  $T(T) = 10k$

Selectivity of  $R.B = S.B$  is  $1/3$

Selectivity of  $S.C = T.C$  is  $1/10$

Selectivity of  $R.A < 40$  is  $1/2$

What is the estimated size of the query output ?

# Rule of Thumb

- If selectivities are unknown, then:  
selectivity factor = 1/10  
[System R, 1979]

# Selectivities from Statistics

- Condition is  $A = c$  /\* value selection on R \*/
  - Selectivity =  $1/V(R,A)$
- Condition is  $A < c$  /\* range selection on R \*/
  - Selectivity =  $(c - \text{Low}(R, A)) / (\text{High}(R,A) - \text{Low}(R,A))T(R)$
- Condition is  $A = B$  /\*  $R \bowtie_{A=B} S$  \*/
  - Selectivity =  $1 / \max(V(R,A), V(S,A))$
  - (will explain next)

# Selectivity of $R \bowtie_{A=B} S$

Assumptions:

- Containment of values: if  $V(R,A) \subseteq V(S,B)$ , then the set of A values of R is included in the set of B values of S
  - Note: this indeed holds when A is a foreign key in R, and B is a key in S
- Preservation of values: for any other attribute B,  $V(R \bowtie_A S, B) = V(R, B) \cap V(S, B)$

# Selectivity of $R \bowtie_{A=B} S$

Assume  $V(R,A) \leq V(S,B)$

- Each tuple  $t$  in  $R$  joins with  $T(S)/V(S,B)$  tuple(s) in  $S$
- Hence  $T(R \bowtie_{A=B} S) = T(R) T(S) / V(S,B)$

In general:  $T(R \bowtie_{A=B} S) = T(R) T(S) / \max(V(R,A), V(S,B))$

# Size Estimation for Join

Example:

- $T(R) = 10000$ ,  $T(S) = 20000$
- $V(R,A) = 100$ ,  $V(S,B) = 200$
- How large is  $R \bowtie_{A=B} S$  ?

# Histograms

- Statistics on data maintained by the RDBMS
- Makes size estimation much more accurate (hence, cost estimations are more accurate)

# Histograms

Employee(ssn, name, age)

$T(\text{Employee}) = 25000$ ,  $V(\text{Employee}, \text{age}) = 50$   
 $\min(\text{age}) = 19$ ,  $\max(\text{age}) = 68$

$\sigma_{\text{age}=48}(\text{Employee}) = ?$        $\sigma_{\text{age}>28 \text{ and } \text{age}<35}(\text{Employee}) = ?$

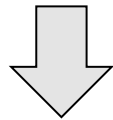


# Histograms

Employee(ssn, name, age)

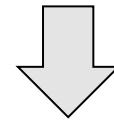
$T(\text{Employee}) = 25000$ ,  $V(\text{Employee}, \text{age}) = 50$   
 $\min(\text{age}) = 19$ ,  $\max(\text{age}) = 68$

$$\sigma_{\text{age}=48}(\text{Employee}) = ?$$



$$\text{Estimate} = 25000 / 50 = 500$$

$$\sigma_{\text{age}>28 \text{ and } \text{age}<35}(\text{Employee}) = ?$$



$$\text{Estimate} = 25000 * 6 / 60 = 2500$$

# Histograms

Employee(ssn, name, age)

$T(\text{Employee}) = 25000$ ,  $V(\text{Employee}, \text{age}) = 50$   
 $\min(\text{age}) = 19$ ,  $\max(\text{age}) = 68$

$\sigma_{\text{age}=48}(\text{Employee}) = ?$

$\sigma_{\text{age}>28 \text{ and } \text{age}<35}(\text{Employee}) = ?$

Age:	0..20	20..29	30-39	40-49	50-59	> 60
Tuples	200	800	5000	12000	6500	500

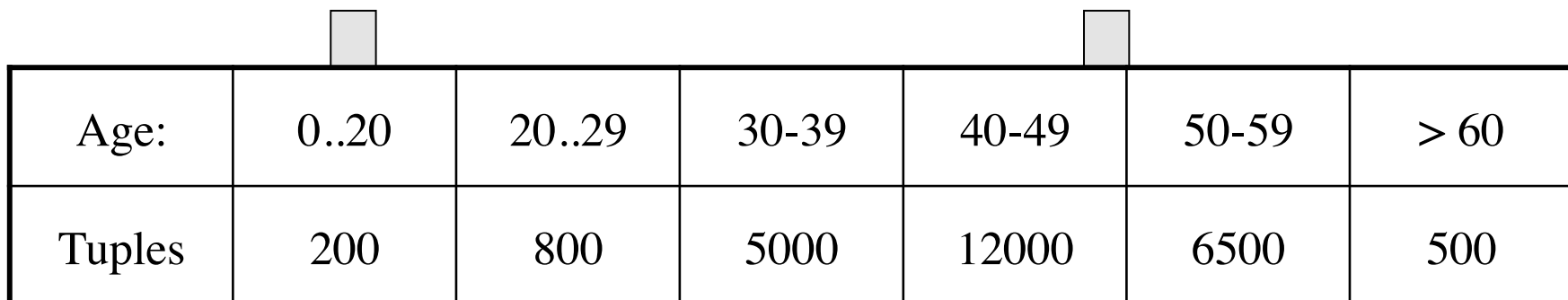
# Histograms

Employee(ssn, name, age)

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A histogram table with two rows: 'Age:' and 'Tuples'. The columns represent age ranges: '0..20', '20..29', '30-39', '40-49', '50-59', and '> 60'. Below the table, two arrows point from the '40-49' and '50-59' columns to the estimates below.

Age:	0..20	20..29	30-39	40-49	50-59	> 60
Tuples	200	800	5000	12000	6500	500

Estimate = 1200

Estimate =  $2 * 80 + 5 * 500 = 2660$  19

# Types of Histograms

- How should we determine the bucket boundaries in a histogram ?

# Types of Histograms

- How should we determine the bucket boundaries in a histogram ?
- Eq-Width
- Eq-Depth
- Compressed

Employee(ssn, name, age)

# Histograms

**Eq-width:**

Age:	0..20	20..29	30-39	40-49	50-59	> 60
Tuples	200	800	5000	12000	6500	500

**Eq-depth:**

Age:	0..20	20..29	30-39	40-49	50-59	> 60
Tuples	1800	2000	2100	2200	1900	1800

**Compressed:** store separately some highly frequent values: (48,1900)

# Difficult Questions on Histograms

- Small number of buckets
  - Hundreds, or thousands, but not more
  - WHY ?
- *Not* updated during database update, but recomputed periodically
  - WHY ?
- Multidimensional histograms rarely used
  - WHY ?

# Summary of Query Optimization

- Three parts:
  - search space, algorithms, size/cost estimation
- Ideal goal: find optimal plan. But
  - Impossible to estimate accurately
  - Impossible to search the entire space
- Goal of today's optimizers:
  - Avoid very bad plans