# Lecture 02: SQL

Wednesday, March 31st, 2010

### Accessing SQL Server

- Host: IISQLSRV.cs.washington.edu
- Authentication: SQL Server Authentication
- User: YOUR\_USER\_NAME@u.washington.edu
- Password: 'cse444login!' (without the quotes)
- Change your password!

#### Outline

- Data in SQL
- Simple Queries in SQL (6.1)
- Queries with more than one relation (6.2)
- Subqueries (6.3)

#### SQL

- Data Definition Language (DDL)
  - Create/alter/delete tables and their attributes
  - Following lectures...
- Data Manipulation Language (DML)
  - Query one or more tables discussed next!
  - Insert/delete/modify tuples in tables

Table name

Tables in SQL

Attribute names

**Product** 

Key

<u>PName</u>	Price	Category	Manufacturer
Gizmo	\$19.99	Gadgets	GizmoWorks
Powergizmo	\$29.99	Gadgets	GizmoWorks
SingleTouch	\$149.99	Photography	Canon
MultiTouch	\$203.99	Household	Hitachi

# Data Types in SQL

- Atomic types:
  - Characters: CHAR(20), VARCHAR(50)
  - Numbers: INT, BIGINT, SMALLINT, FLOAT
  - Others: MONEY, DATETIME, …
- Record (aka tuple)
  - Has atomic attributes
- Table (relation)
  - A set of tuples

# Simple SQL Query

#### **Product**

PName	Price	Category	Manufacturer
Gizmo	\$19.99	Gadgets	GizmoWorks
Powergizmo	\$29.99	Gadgets	GizmoWorks
SingleTouch	\$149.99	Photography	Canon
MultiTouch	\$203.99	Household	Hitachi

SELECT \*
FROM Product
WHERE category='Gadgets'





PName	Price	Category	Manufacturer
Gizmo	\$19.99	Gadgets	GizmoWorks
Powergizmo	\$29.99	Gadgets	GizmoWorks

# Simple SQL Query

#### **Product**

PName	Price	Category	Manufacturer
Gizmo	\$19.99	Gadgets	GizmoWorks
Powergizmo	\$29.99	Gadgets	GizmoWorks
SingleTouch	\$149.99	Photography	Canon
MultiTouch	\$203.99	Household	Hitachi

SELECT PName, Price, Manufacturer FROM Product WHERE Price > 100



"selection" and "projection"

PName	Price	Manufacturer
SingleTouch	\$149.99	Canon
MultiTouch	\$203.99	Hitachi

#### **Details**

Case insensitive:

SELECT = Select = select

Product = product

BUT: 'Seattle' ≠ 'seattle'

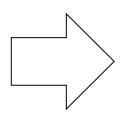
• Constants:

'abc' - yes

"abc" - no

# **Eliminating Duplicates**

SELECT DISTINCT category
FROM Product



Category

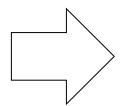
Gadgets

Photography

Household

#### Compare to:

SELECT category FROM Product



Category

Gadgets

Gadgets

Photography

Household

# Ordering the Results

```
SELECT pname, price, manufacturer

FROM Product
```

WHERE category='gizmo' AND price > 50

ORDER BY price, pname

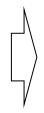
Ties are broken by the second attribute on the ORDER BY list.

Ordering is ascending, unless you specify the DESC keyword.

PName	Price	Category	Manufacturer
Gizmo	\$19.99	Gadgets	GizmoWorks
Powergizmo	\$29.99	Gadgets	GizmoWorks
SingleTouch	\$149.99	Photography	Canon
MultiTouch	\$203.99	Household	Hitachi

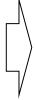
SELECT DISTINCT category
FROM Product

**ORDER BY** category





SELECT Category
FROM Product
ORDER BY PName





SELECT DISTINCT category
FROM Product
ORDER BY PName





# Keys and Foreign Keys

#### Company

	<u>CName</u>	StockPrice	Country
Key	GizmoWorks	25	USA
	Canon	65	Japan
	Hitachi	15	Japan

#### **Product**

<u>PName</u>	Price	Category	Manufacturer
Gizmo	\$19.99	Gadgets	GizmoWorks
Powergizmo	\$29.99	Gadgets	GizmoWorks
SingleTouch	\$149.99	Photography	Canon
MultiTouch	\$203.99	Household	Hitachi

Foreign key

#### **Joins**

Product (<u>pname</u>, price, category, manufacturer) Company (<u>cname</u>, stockPrice, country)

Find all products under \$200 manufactured in Japan; return their names and prices.

SELECT PName, Price and Company

Product, Company

WHERE Manufacturer=CName AND Country='Japan'

AND Price <= 200

#### **Joins**

#### **Product**

#### **PName Price** Category Manufacturer Gadgets GizmoWorks Gizmo 19.99 GizmoWorks Powergizmo \$29.99 Gadgets 110 0 SingleTouch Photography Canon MultiTouch \$203.99 Household Hitaum

#### Company

Cname	StockPrice	Country
Gizillovvoiks	25	USA
Canon	65	Japan
Hitachi	15	Japan

**SELECT** PName, Price

FROM Product, Company

WHERE Manufacturer=CName AND Country='Japan'

AND Price <= 200



PName	Price	
SingleTouch	\$149.99	

# Tuple Variables

Person(<u>pname</u>, address, worksfor)

Company(cname, address)

SELECT DISTINCT pname, address

FROM Person, Company

WHERE worksfor = cname

SELECT DISTINCT Person.pname, Company.address

FROM Person, Company

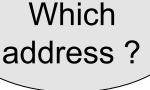
WHERE Person.worksfor = Company.cname

SELECT DISTINCT x.pname, y.address

FROM Person AS x, Company AS y

WHERE x.worksfor = y.cname

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#### In Class

Product (pname, price, category, manufacturer)
Company (cname, stockPrice, country)

Find all Chinese companies that manufacture products both in the 'toy' category

**SELECT** cname

**FROM** 

**WHERE** 

#### In Class

Product (pname, price, category, manufacturer)
Company (cname, stockPrice, country)

Find all Chinese companies that manufacture products both in the 'electronic' and 'toy' categories

**SELECT** cname

**FROM** 

WHERE

#### Meaning (Semantics) of SQL Queries

```
SELECT a_1, a_2, ..., a_k
FROM R_1 AS x_1, R_2 AS x_2, ..., R_n AS x_n
WHERE Conditions
```

```
\label{eq:formula} \begin{array}{l} \text{Answer} = \{\} \\ \text{for } x_1 \text{ in } R_1 \text{ do} \\ \text{for } x_2 \text{ in } R_2 \text{ do} \\ & \dots \\ \text{for } x_n \text{ in } R_n \text{ do} \\ & \text{if Conditions} \\ & \text{then } \text{Answer} = \text{Answer} \cup \{(a_1, \dots, a_k)\} \\ \text{return } \text{Answer} \end{array}
```

# Using the Formal Semantics

What do these queries compute?

SELECT DISTINCT R.A FROM R, S WHERE R.A=S.A

Returns  $R \cap S$ 

```
SELECT DISTINCT R.A
FROM R, S, T
WHERE R.A=S.A OR R.A=T.A
```

If  $S \neq \emptyset$  and  $T \neq \emptyset$ then returns  $R \cap (S \cup T)$ else returns  $\emptyset$ 

# Joins Introduce Duplicates

Product (pname, price, category, manufacturer)
Company (cname, stockPrice, country)

Find all countries that manufacture some product in the 'Gadgets' category.

**SELECT** Country

FROM Product, Company

WHERE Manufacturer=CName AND Category='Gadgets'

# Joins Introduce Duplicates

#### **Product**

<u>Name</u>	Price	Category	Manufacturer
Gizmo	\$19.99	Gadgets	GizmoWorks
Powergizmo	\$29.99	Gadgets	GizmoWorks
SingleTouch	\$149.99	Photography	Canon
MultiTouch	\$203.99	Household	Hitachi

#### Company

<u>Cname</u>	StockPrice	Country
GizmoWorks	25	USA
Canon	65	Japan
Hitachi	15	Japan

**SELECT** Country

FROM Product, Company

WHERE Manufacturer=CName AND Category='Gadgets'



Duplicates!
Remember to add DISTINCT

Country
USA
USA

# Subqueries

- A subquery is another SQL query nested inside a larger query
- Such inner-outer queries are called nested queries
- A subquery may occur in:
  - A SELECT clause
  - 2. A FROM clause
  - 3. A WHERE clause

Rule of thumb: avoid writing nested queries when possible; keep in mind that sometimes it's impossible

# 1. Subqueries in SELECT

```
Product (<u>pname</u>, price, company)
Company(<u>cname</u>, city)
```

For each product return the city where it is manufactured

```
SELECT X.pname, (SELECT Y.city
FROM Company Y
WHERE Y.cname=X.company)
FROM Product X
```

What happens if the subquery returns more than one city?

# 1. Subqueries in SELECT

```
Product (<u>pname</u>, price, company)
Company(<u>cname</u>, city)
```

Whenever possible, don't use a nested queries:

SELECT pname, (SELECT city FROM Company WHERE cname=company)
FROM Product

SELECT pname, city
FROM Product, Company
WHERE cname=company

We have "unnested" the query

# 1. Subqueries in SELECT

```
Product (<u>pname</u>, price, company)
Company(<u>cname</u>, city)
```

Compute the number of products made in each city

```
SELECT DISTINCT city, (SELECT count(*)
FROM Product
WHERE cname=company)
FROM Company
```

Better: we can unnest by using a GROUP BY (next lecture)

### 2. Subqueries in FROM

```
Product (<u>pname</u>, price, company)
Company(<u>cname</u>, city)
```

Find all products whose prices is > 20 and < 30

```
SELECT X.city
FROM (SELECT * FROM Product AS Y WHERE Y.price > 20) AS X
WHERE X.price < 30
```

Unnest this query!

```
Product (<u>pname</u>, price, company) Existential quantifiers
Company(<u>cname</u>, city)
```

Find all cities that make <u>some</u> products with price < 100

#### Using EXISTS:

```
SELECT DISTINCT Company.city
FROM Company
WHERE EXISTS (SELECT *
FROM Product
WHERE company = cname and Produc.price < 100)
```

Product (<u>pname</u>, price, company) Existential quantifiers
Company(<u>cname</u>, city)

Find all cities that make <u>some</u> products with price < 100

Predicate Calculus (a.k.a. First Order Logic)

 $\{y \mid \exists x. Company(x,y) \land (\exists z. \exists p. Product(z,p,x) \land p < 100)\}$ 

```
Product (<u>pname</u>, price, company) Existential quantifiers
Company(<u>cname</u>, city)
```

Find all cities that make <u>some</u> products with price < 100

#### Using IN

```
SELECT DISTINCT Company.city
FROM Company
WHERE Company.cname IN (SELECT Product.company
FROM Product
WHERE Produc.price < 100)
```

```
Product (<u>pname</u>, price, company) Existential quantifiers Company(<u>cname</u>, city)
```

Find all cities that make <u>some</u> products with price < 100

#### Using ANY:

```
SELECT DISTINCT Company.city
FROM Company
WHERE 100 > ANY (SELECT price
FROM Product
WHERE company = cname)
```

Product (pname, price, company) Existential quantifiers Company(<u>cname</u>, city)

Find all cities that make some products with price < 100

Now let's unnest it:

SELECT DISTINCT Company.cname

Company, Product FROM

WHERE Company.cname = Product.company and Product.price < 100

Product (<u>pname</u>, price, company) Company(<u>cname</u>, city) Universal quantifiers

Find all cities with companies that make <u>only</u> products with price < 100

Universal quantifiers are hard!

Product (pname, price, company) Universal quantifiers Company(cname, city)

Find all cities with companies that make only products with price < 100

Predicate Calculus (a.k.a. First Order Logic)

 $\{y \mid \exists x. Company(x,y) \land (\forall z. \forall p. Product(z,p,x) \rightarrow p < 100)\}$ 

De Morgan's Laws:

$$\neg(A \land B) = \neg A \lor \neg B$$
  
 $\neg(A \lor B) = \neg A \land \neg B$   
 $\neg \forall x. P(x) = \exists x. \neg P(x)$   
 $\neg \exists x. P(x) = \forall x. \neg P(x)$ 

$$\neg(A \rightarrow B) = A \wedge \neg B$$

```
\{y \mid \exists x. Company(x,y) \land (\forall z. \forall p. Product(z,p,x) \rightarrow p < 100)\}
```

```
\{y \mid \exists x. Company(x,y) \land \neg (\exists z \exists p. Product(z,p,x) \land p \ge 100)\}
```

```
\{y \mid \exists x. Company(x,y)\} - \{y \mid \exists x. Company(x,y) \land (\exists z \exists p. Product(z,p,x) \land p \ge 100)\}
```

1. Find *the other* companies: i.e. s.t. <u>some</u> product ≥ 100

```
SELECT DISTINCT Company.city
FROM Company
WHERE Company.cname IN (SELECT Product.company
FROM Product
WHERE Produc.price >= 100
```

2. Find all companies s.t. <u>all</u> their products have price < 100

```
SELECT DISTINCT Company.city
FROM Company
WHERE Company.cname NOT IN (SELECT Product.company
FROM Product
WHERE Produc.price >= 100
```

```
Product (<u>pname</u>, price, company) Universal quantifiers Company(<u>cname</u>, city)
```

Find all cities with companies that make <u>only</u> products with price < 100

#### Using EXISTS:

```
SELECT DISTINCT Company.city
FROM Company
WHERE NOT EXISTS (SELECT *
FROM Product
WHERE company = cname and Produc.price >= 100)
```

```
Product (pname, price, company) Universal quantifiers
Company(cname, city)
```

Find all cities that make <u>some</u> products with price < 100

#### Using ALL:

```
SELECT DISTINCT Company.city
FROM Company
WHERE 100 > ALL (SELECT price
FROM Product
WHERE company = cname)
```

# Question for Database Fans and their Friends

Can we unnest the universal quantifier query?

#### **Monotone Queries**

- A query Q is monotone if:
  - Whenever we add tuples to one or more of the tables...
  - ... the answer to the query cannot contain fewer tuples
- **Fact**: all unnested queries are monotone
  - Proof: using the "nested for loops" semantics
- Fact: A query a universal quantifier is not monotone
- <u>Consequence</u>: we cannot unnest a query with a universal quantifier Dan Suciu -- 444 Spring 2010

#### Queries that must be nested

- Queries with universal quantifiers or with negation
- The drinkers-bars-beers example next
- This is a famous example from textbook on databases by Ullman

### The drinkers-bars-beers example

Likes(drinker, beer)
Frequents(drinker, bar)
Serves(bar, beer)

Challenge: write these in SQL

Find drinkers that frequent <u>some</u> bar that serves <u>some</u> beer they like.

```
x: \exists y. \exists z. Frequents(x, y) \land Serves(y,z) \land Likes(x,z)
```

Find drinkers that frequent only bars that serves some beer they like.

```
x: \forall y. Frequents(x, y)\Rightarrow (\exists z. Serves(y,z)\wedgeLikes(x,z))
```

Find drinkers that frequent <u>some</u> bar that serves <u>only</u> beers they like.

```
x: \exists y. Frequents(x, y) \land \forall z.(Serves(y,z) \Rightarrow Likes(x,z))
```

Find drinkers that frequent only bars that serves only beer they like.

```
x: \forall y. Frequents(x, y)\Rightarrow \forall z.(Serves(y,z)\Rightarrow Likes(x,z))
```