

CSE 431 Spring 2015

Assignment #4

Due: Monday, May 4, 2015

Reading assignment: Read Sections 7.1 and 7.2 of Sipser's text.

Problems:

1. Let $J = \{w \mid w = 0x \text{ for some } x \in A_{TM} \text{ or } w = 1y \text{ for some } y \in \overline{A_{TM}}\}$. Show that neither J nor \overline{J} is Turing-recognizable.
2. Show that there is an undecidable language contained in 1^* .
3. Which of the following problems are decidable? Justify each answer:
 - (a) Given a Turing machine M , does M accept 0101?
 - (b) Given Turing machines M and N , is $L(N)$ the complement of $L(M)$?
 - (c) Given a Turing machine M , integers a and b and an input x , does M run for more than $a|x|^2 + b$ steps on input x ?
4. Prove that if K and L are decidable by Turing machines running in polynomial time then so are $K \cup L$, KL , and \overline{L} .
5. Let $TRI = \{\langle G \rangle \mid G \text{ is an undirected graph that contains a triangle}\}$. Prove that there is a polynomial-time Turing machine that decides TRI .
6. (Bonus) Show that the following problem is undecidable: Given a Turing machine M and integers a and b , does there exist an input x on which M runs for more than $a|x|^2 + b$ steps on input x ?
7. (Bonus) We showed previously that neither EQ_{TM} nor its complement is Turing-recognizable. Your problem is to show that, despite this, if you had a magic black box that decided A_{TM} that you could call repeatedly on different inputs, then you could decide $\overline{EQ_{TM}}$.