

Lecture 10

Derandomization, quick sort, and stable matching I

Chinmay Nirkhe | CSE 421 Winter 2026

W

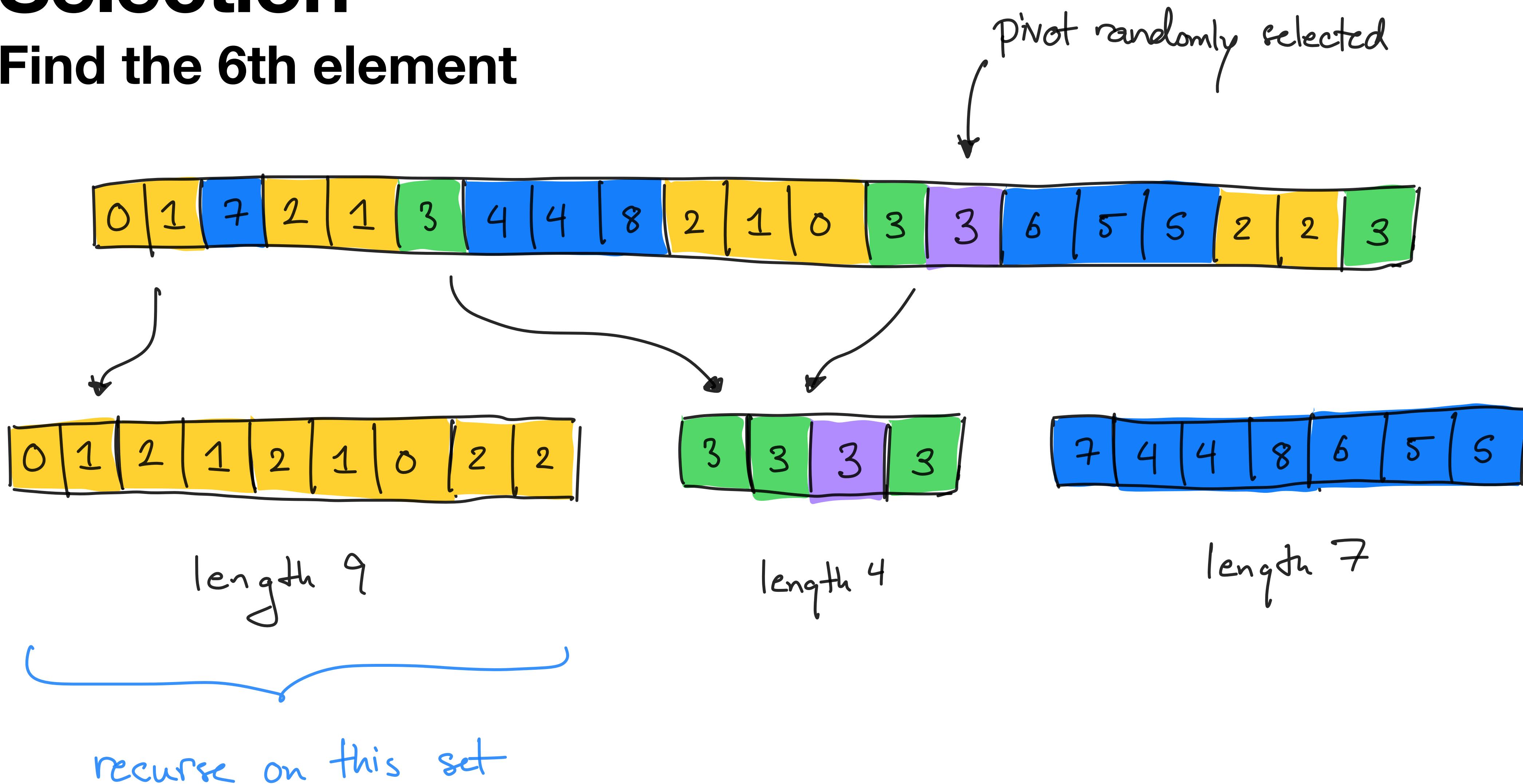
Writing quality

- I've been pretty pleased with the solutions for the first two sets of problems
- You may have gotten a -0.01 for writing quality
 - Take the time to chat with TAs about how to improve writing
 - Take a look at the solutions to see how we might write the solutions
 - Emulate solutions from section worksheets
- Starting with set 4, we will increase the deduction for writing quality if egregious

Previously in CSE 421...

Selection

Find the 6th element



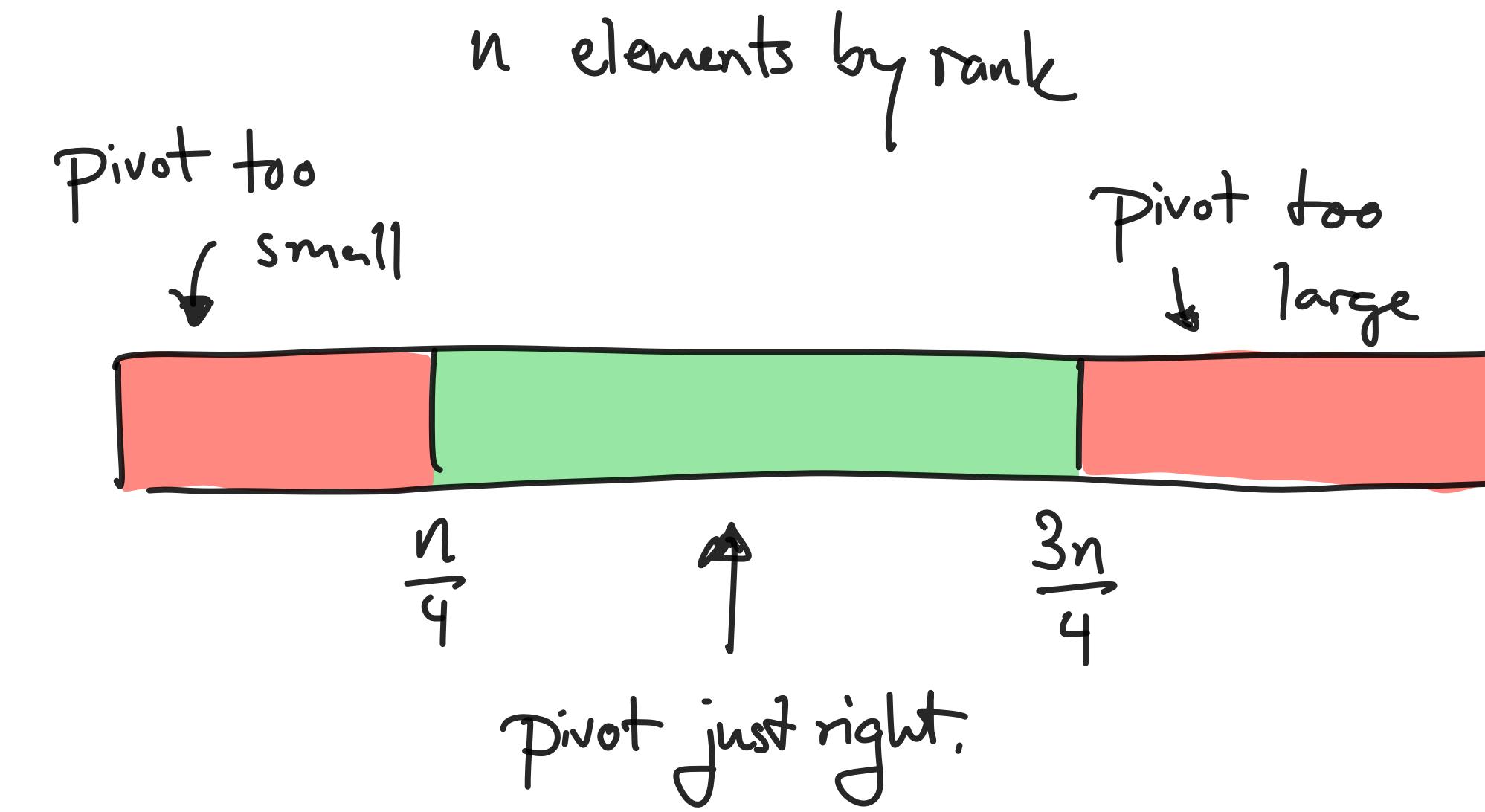
Selection

- **Recursive algorithm** $\text{Selection}(X, k)$:
 - Randomly sample j from $[n]$. Call x_j the “**pivot**”.
 - Filter X into X_L , X_E , and X_R based on if $x_i < x_j$, $x_i = x_j$, or $x_i > x_j$.
 - If $|X_L| \geq k$, recursively output $\text{Selection}(X_L, k)$.
 - Else if, $|X_L| + |X_E| \geq k$, output x_j .
 - Else, recursively output $\text{Selection}(X_R, k - |X_L| - |X_E|)$.

Runtime analysis

- In order to apply the master theorem, we would need to argue that each recursive call was reducing the input size from n to n/b for $b > 1$
 - $T(n) = T(n/b) + cn \implies T(n) = \frac{c}{1 - 1/b}n$
- However, each call may not reduce the size from n to n/b
- Depends on how close the randomly chosen x_j is to the middle
 - If pivot x_j was the largest element, then $|X_L| = n - 1$, $|X_E| = 1$, and $|X_R| = 0$.
 - Decreases instance size from n to $n - 1$.
 - Fortunately, the probability this occurs is $1/n$.

Runtime analysis



- **Amortized analysis:**
 - If pivot x_j is the ℓ -th element, then the next problem is of size $\leq \max\{\ell, n - \ell\}$.
 - With probability $\geq 1/2$, pivot x_j is the ℓ -th element for $\ell \in \{n/4, \dots, 3n/4\}$.
 - The expected compute in reducing from n -sized instance to a $3n/4$ -sized instance is $O(n)$.
 - Total **expected** runtime: $T(n) = T(3n/4) + O(n) \implies T(n) = O(n)$.

Runtime analysis

- **Amortized analysis:**

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- The expected compute in reducing from n -sized instance to a $3n/4$ -sized instance is $O(n)$.
 - $\geq 1/2$ probability, shrinks in 1 reduction.
 - $\geq 1/4$ probability, shrinks in 2 reductions.
 - ... $\geq 1/2^j$ probability, shrinks in j reductions ...
- Expected compute is $\leq O(n) \cdot \left(\frac{1}{2} + \frac{1}{4} \cdot 2 + \frac{1}{8} \cdot 3 + \dots \right) = O(n) \cdot 2$
- Total **expected** runtime: $T(n) = T(3n/4) + O(n) \implies T(n) = O(n)$.

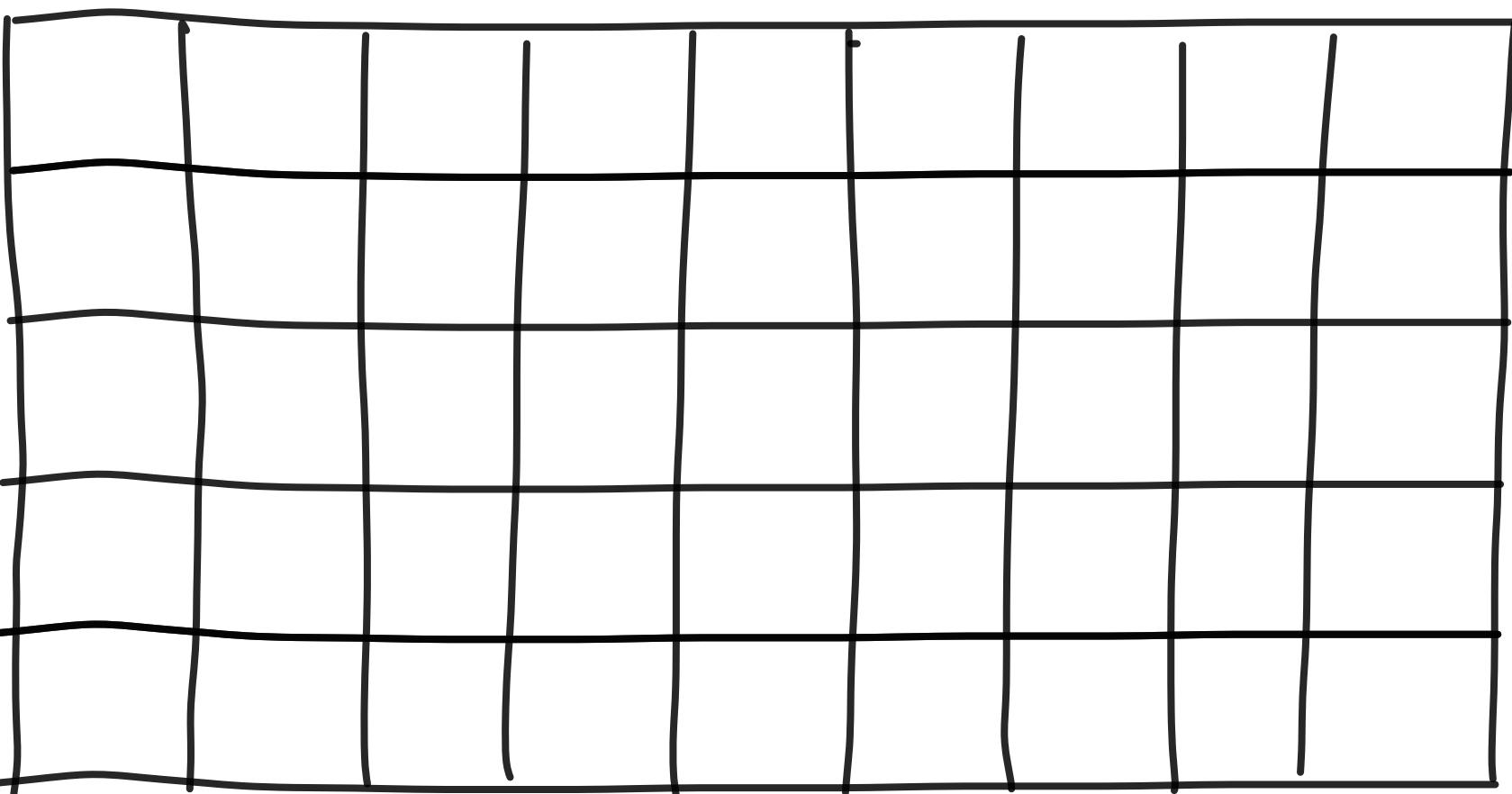
Derandomization

- The worst case runtime is $O(n^2)$.
 - Only happens with $2^{-\Omega(n \log n)}$ probability.
- But, is there an algorithm that didn't require randomness?
 - Why?
- Recall, if we could guarantee that the pivot x_j was in the middle half, then each recursion would decrease in size by $3/4$.
- **Blum-Pratt-Floyd-Rivest-Tarjan (1973)**: Calculate a pivot in the middle $4n/10$ in time $O(n)$.

Pivot selection algorithm

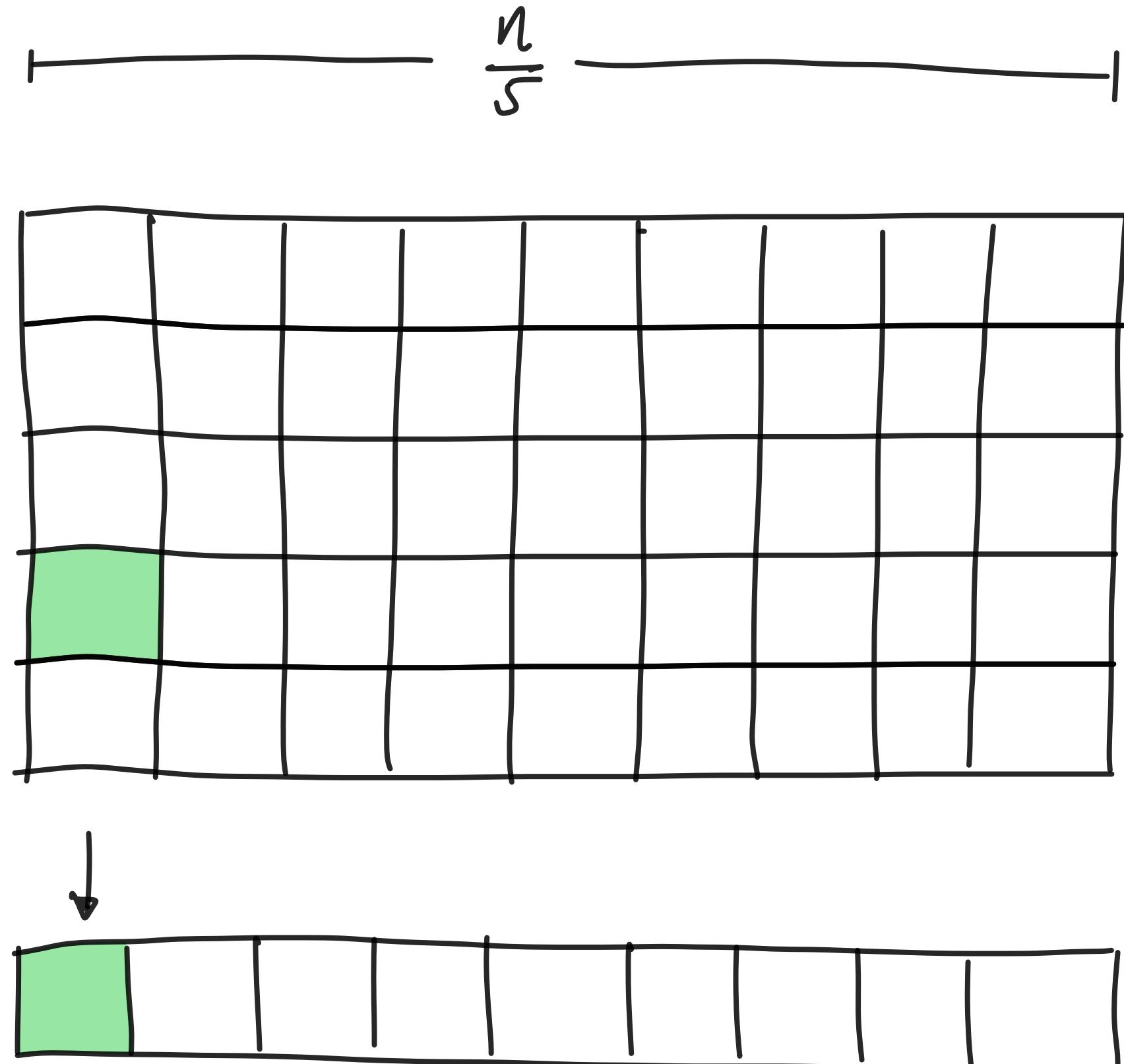
- Express the n elements as a $5 \times (n/5)$ matrix of elements

$$\frac{n}{5}$$



Pivot selection algorithm

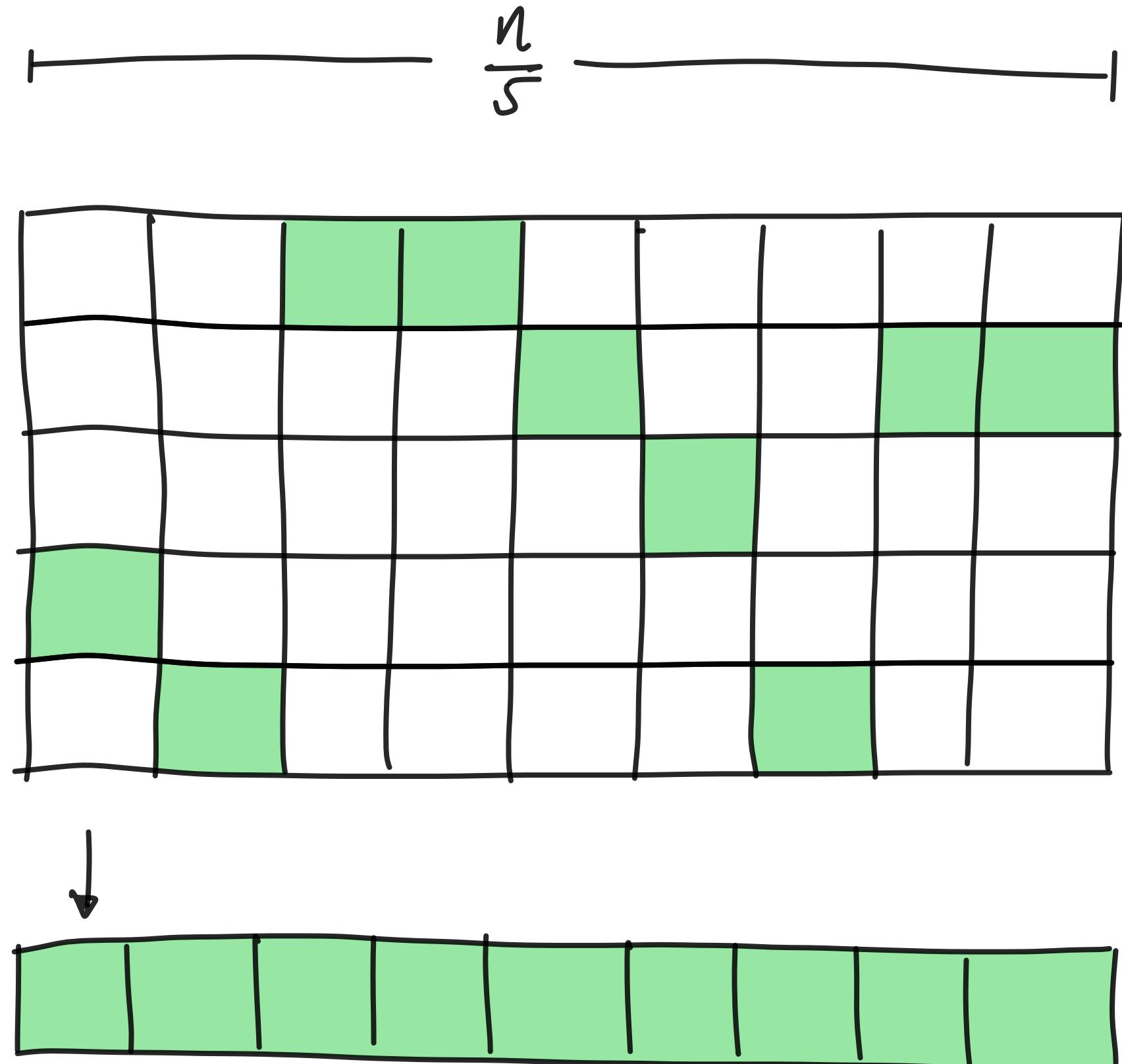
- Express the n elements as a $5 \times (n/5)$ matrix of elements
- Calculate the medians of each of the columns:
$$Y = (y_1, y_2, \dots, y_{n/5})$$



the median is one of
the 5 elements in the
column

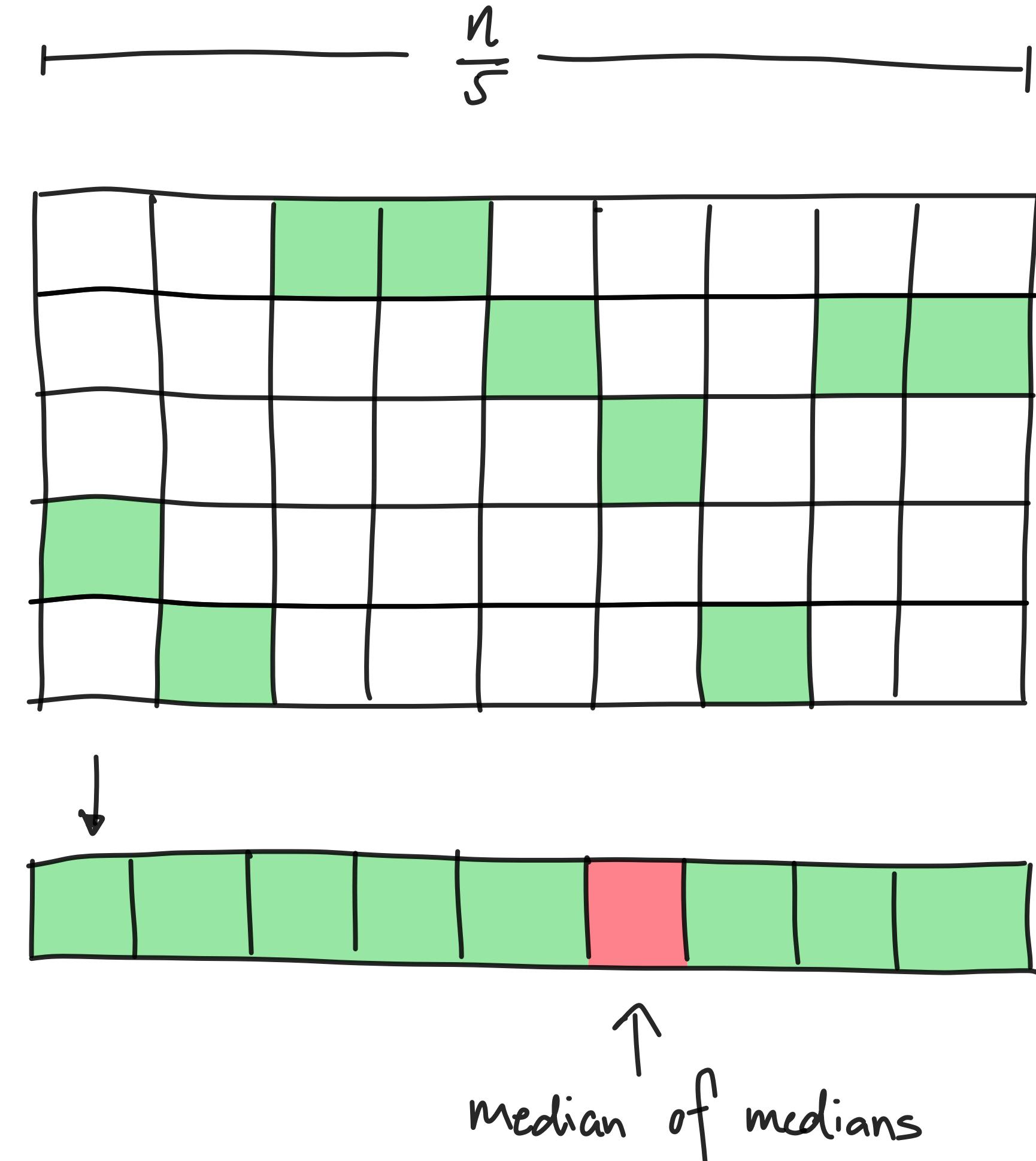
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- Calculate the medians of each of the columns:
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- Choose the pivot as the median of the medians:
$$p \leftarrow \text{median}(Y)$$



Pivot selection algorithm

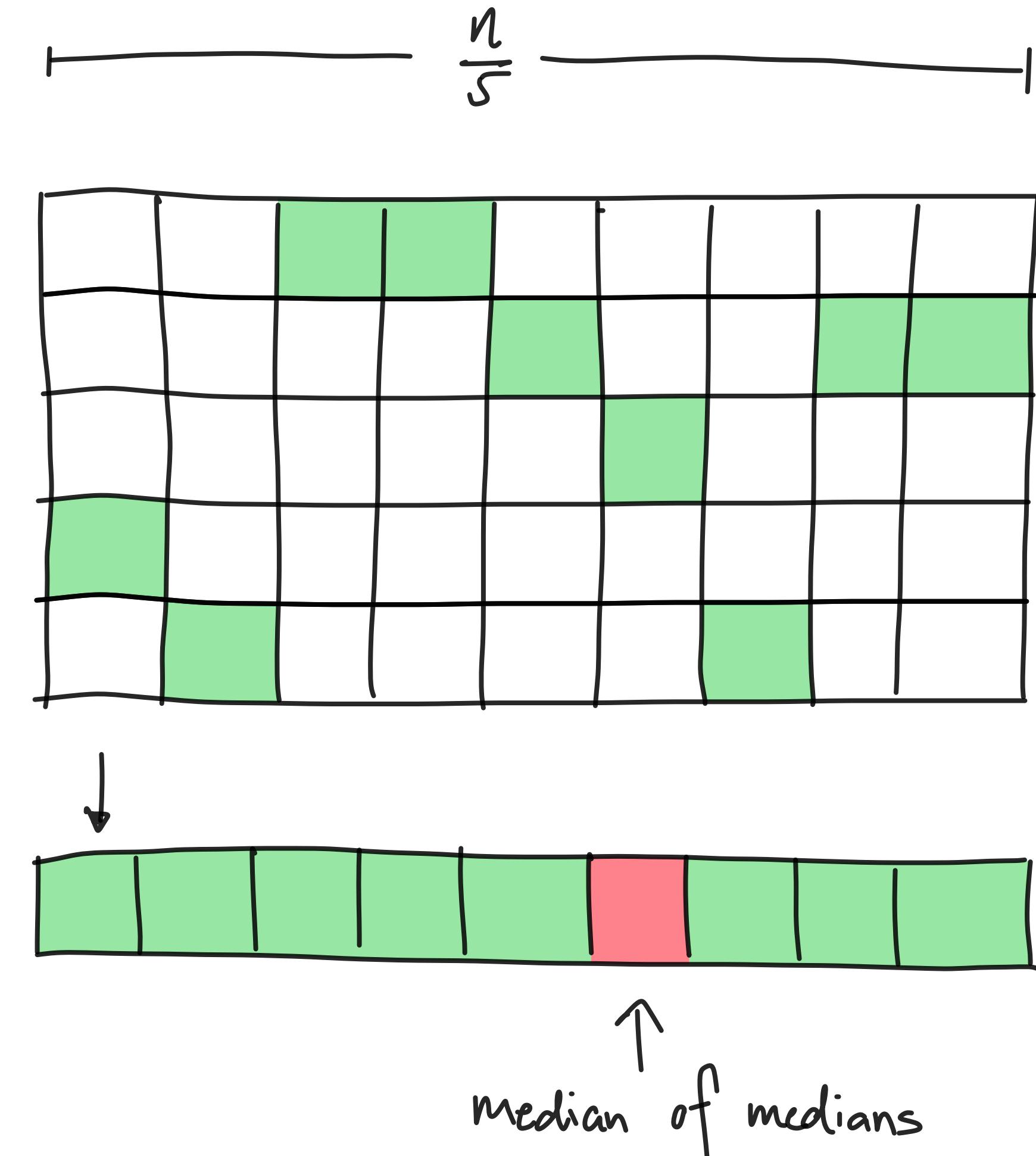
Runtime analysis

- Express the n elements as a $5 \times (n/5)$ matrix of elements
- Calculate the medians of each of the columns:
$$Y = (y_1, y_2, \dots, y_{n/5})$$

$$\leftarrow O(1) \text{ per col.}$$

$$\text{Total } O(n).$$
- Choose the pivot as the median of the medians:
$$p \leftarrow \text{median}(Y)$$
 $T(n/5)$ recursively

Only semantic.

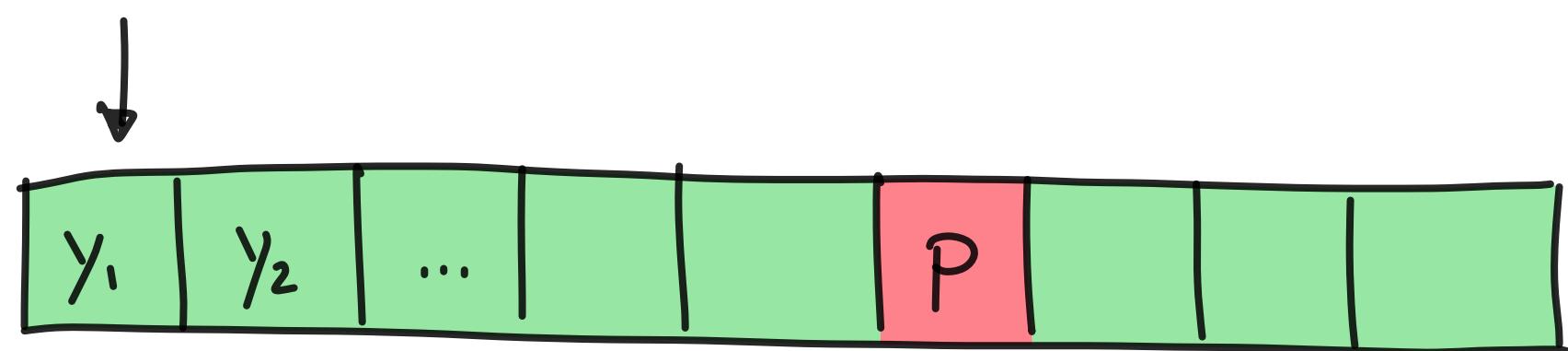
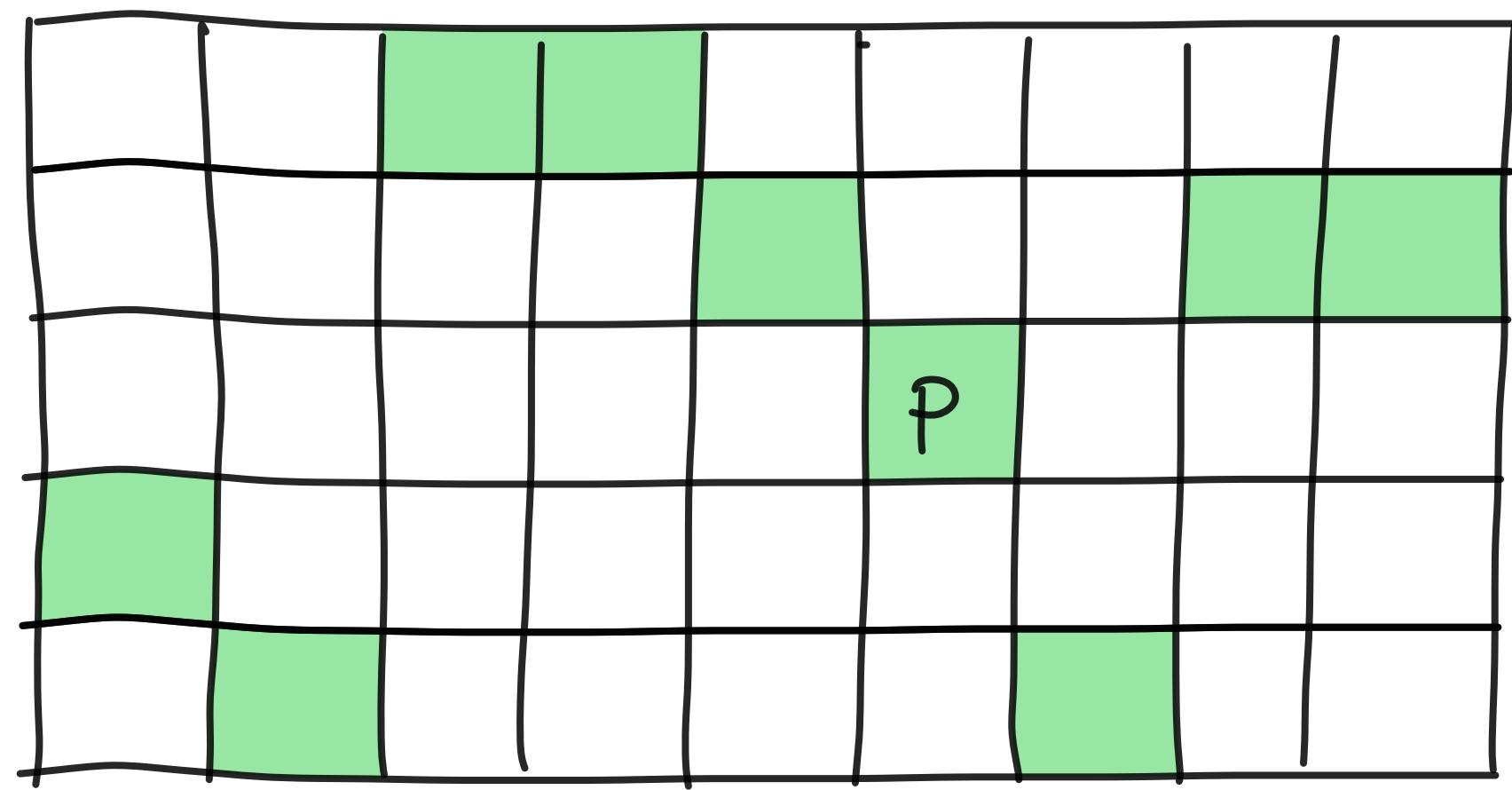


$$\text{Total time: } T(n) = T\left(\frac{n}{5}\right) + O(n) \Rightarrow T(n) = O(n).$$

Pivot selection algorithm

Proof of correctness

$$\frac{n}{5}$$

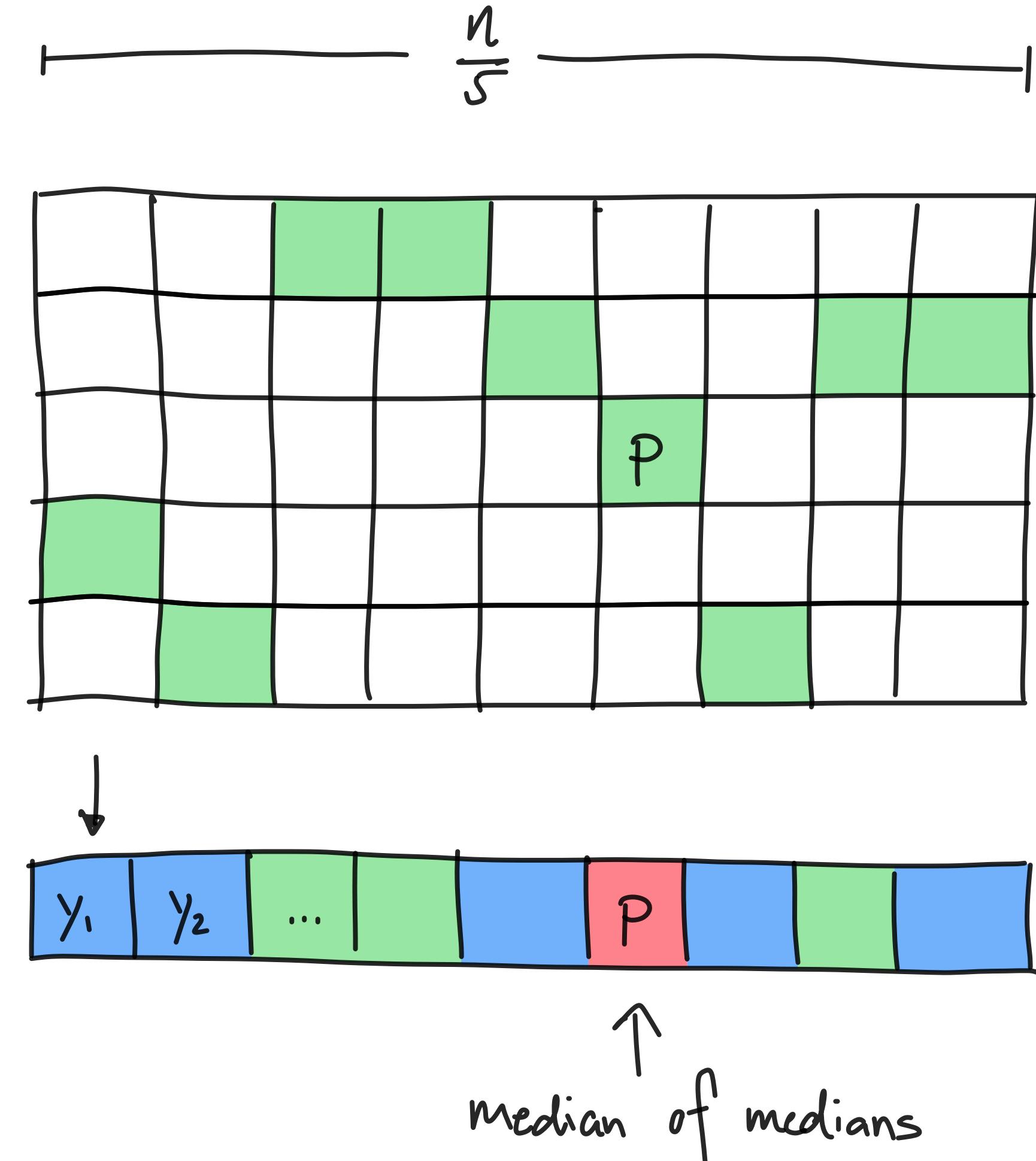


↑
median of medians

Pivot selection algorithm

Proof of correctness

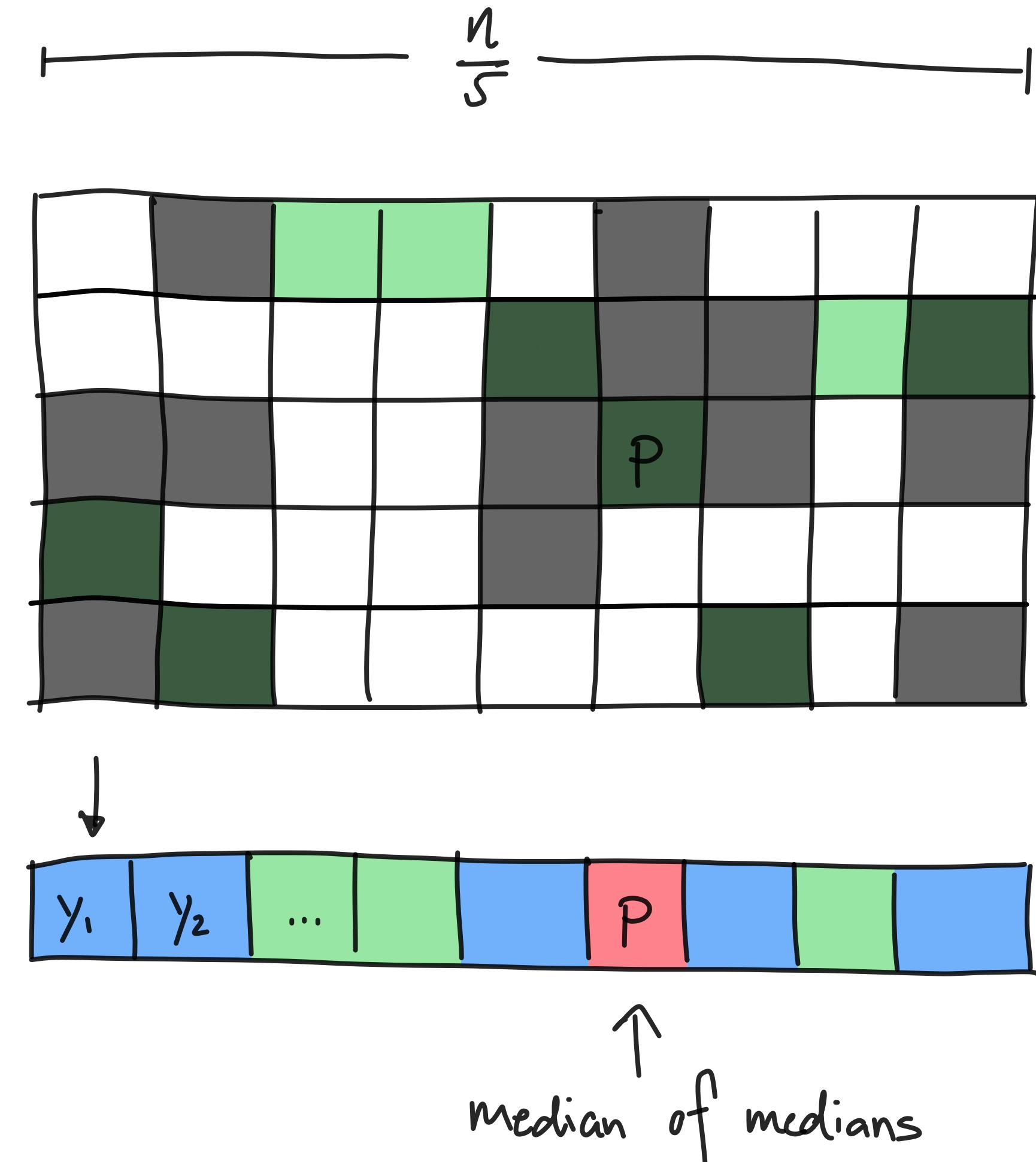
- There are $\geq n/10$ columns such that $y_j \geq p$.



Pivot selection algorithm

Proof of correctness

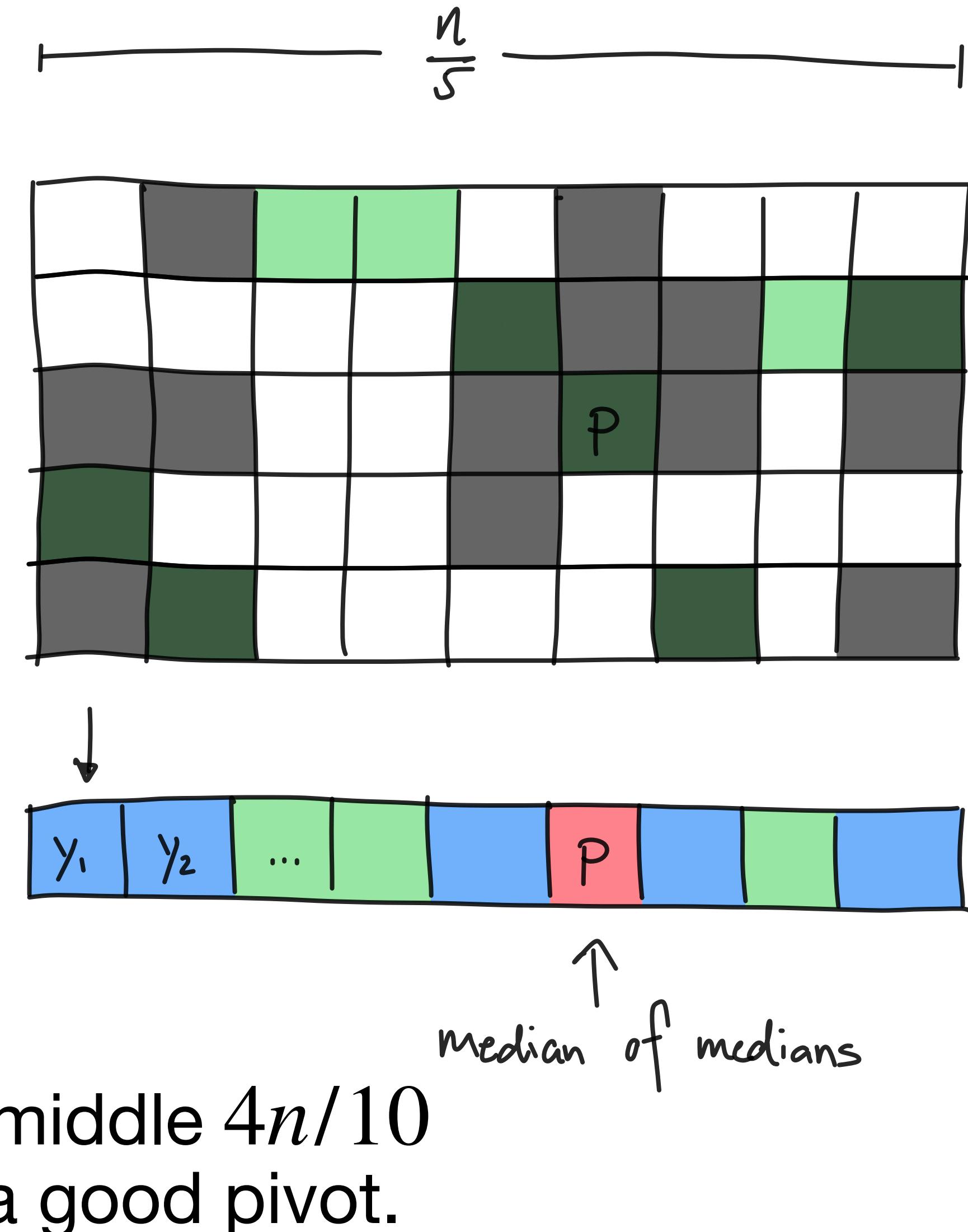
- There are $\geq n/10$ columns such that $y_j \geq p$.
- In each such column, there are 3 elements $\geq y_j$.



Pivot selection algorithm

Proof of correctness

- There are $\geq n/10$ columns such that $y_j \geq p$.
- In each such column, there are 3 elements $\geq y_j$.
- Therefore, there are $\geq 3n/10$ elements $\geq p$.
- Similarly, there are $\geq 3n/10$ elements $\leq p$.
 - So, p is in the middle $4n/10$ elements and a good pivot.



Median/Selection algorithm

$$\begin{aligned}\text{Total: } T(n) &= T\left(\frac{7}{10}n\right) + T\left(\frac{n}{5}\right) + O(n) \\ &\Rightarrow T(n) = O(n)\end{aligned}$$

- **Input:** $(X, k) \in \mathbb{R}^n \times [n]$
- **Output:** the k -th item in the list X
- **Algorithm:**

- Calculate $p \leftarrow \text{median-of-medians}(X)$ in a $5 \times (n/5)$ division.
- Filter X into X_L, X_E , and X_R based on p
- If $|X_L| \geq k$, recurse Selection(X_L, k)
- Else if $|X_L| + |X_E| \geq k$, return p
- Else, return Selection($X_R, k - |X_L| - |X_E|$).

↑ Pset problem on how to analyze this
generalization of Master theorem

recursive $T\left(\frac{n}{5}\right) + O(n)$

recursive $T\left(\frac{7}{10}n\right)$

Quicksort algorithm

- The algorithm we just analyzed, “Quickselect”, can be generalized to sorting
- **Sorting algorithm** Quicksort(X):
 - Pick a pivot p (either randomized or with median-of-medians)
 - Filter X into X_L, X_E, X_R by comparing elements with p
 - Concatenate Sort(X_L), X_E , Sort(X_R).

↑ ↑
Computing expected runtime is challenging due to
variable size

Quicksort algorithm

Runtime analysis

- Runtime depends on pivot selection
- **Median-of-medians:**
 - $T(n) \leq T(\alpha n) + T(n - \alpha n) + O(n)$ for $\alpha \in [0.3, 0.7]$
 - $T(n) = O(n \log n)$ by analysis in problem 54
- **Choose random element:**
 - Worst case: $O(n^2)$ time
 - Amortized: $O(n \log n)$ (next!)

Quicksort algorithm

Runtime analysis

[for random choice of pivot]

- **Observations:**

- The runtime of Quicksort is proportional to the number of comparisons
- The algorithm only compares two elements if one is the pivot
- Let $Y = (y_1, \dots, y_n)$ be the sorted version of the input.
- Let $p_{ij} = \Pr [y_i \text{ and } y_j \text{ are compared}]$
- **Claim:** $p_{ij} \leq \frac{2}{j - i + 1}$ when $i < j$.

Expected number of comparisons:

$$\begin{aligned} \sum_{i < j} p_{ij} &\leq 2 \sum_{i=1}^n \sum_{j=i+1}^n \frac{1}{j - i + 1} \\ &= 2 \sum_{i=1}^n \sum_{k=1}^{n-i+1} \frac{1}{k+1} \\ &= 2 \sum_{i=1}^n \left(\underbrace{\frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n-i+1}}_{\leq \log(n-i+1) + 1} \right) \\ &\leq \log(n) + 1 \\ &\leq 2n \log n + 2n. \end{aligned}$$

Runtime of quicksort = $O(n \log n)$.

Proof of claim

- **Claim:** $p_{ij} \leq \frac{2}{j - i + 1}$ when $i < j$.
- **Proof:**
 - $y_i \leq y_j$ and y_i and y_j are compared at most once
 - Comparisons only occur when one of them is the pivot
 - Case 1: $y_i, y_j \in X_E$ and we never recurse on X_E
 - Case 2: $y_i \in X_E, y_j \in X_R$ and we never compare between X_L, X_E , and X_R
 - Case 3: $y_i \in X_L, y_j \in X_E$ and we never compare between X_L, X_E , and X_R
 - If and when y_i and y_j are compared during $\text{sort}(X')$ then $y_i, y_{i+1}, y_{i+2}, \dots, y_j \in X'$
 - Can be formally proven via induction
 - So $|X'| \geq j - i + 1$.
 - Probability that either y_i or y_j is chosen as pivot is $\leq \frac{2}{j - i + 1}$.

Sorting in the real world

- **Quicksort**

- Fast almost always, especially for in-memory sorting.
- Works well with caches due to good locality of reference.
- In practice,
 - Don't filter X_L , X_E , and X_R . Use in-place swaps.
 - When n is small, insertion sorting is a better base case.
 - Pick pivot randomly for small n , median of 3 random values for medium n , and median-of-medians on 9 elements for large n
 - Never actually run the median-of-medians pivot finding routine

Sorting in the real world

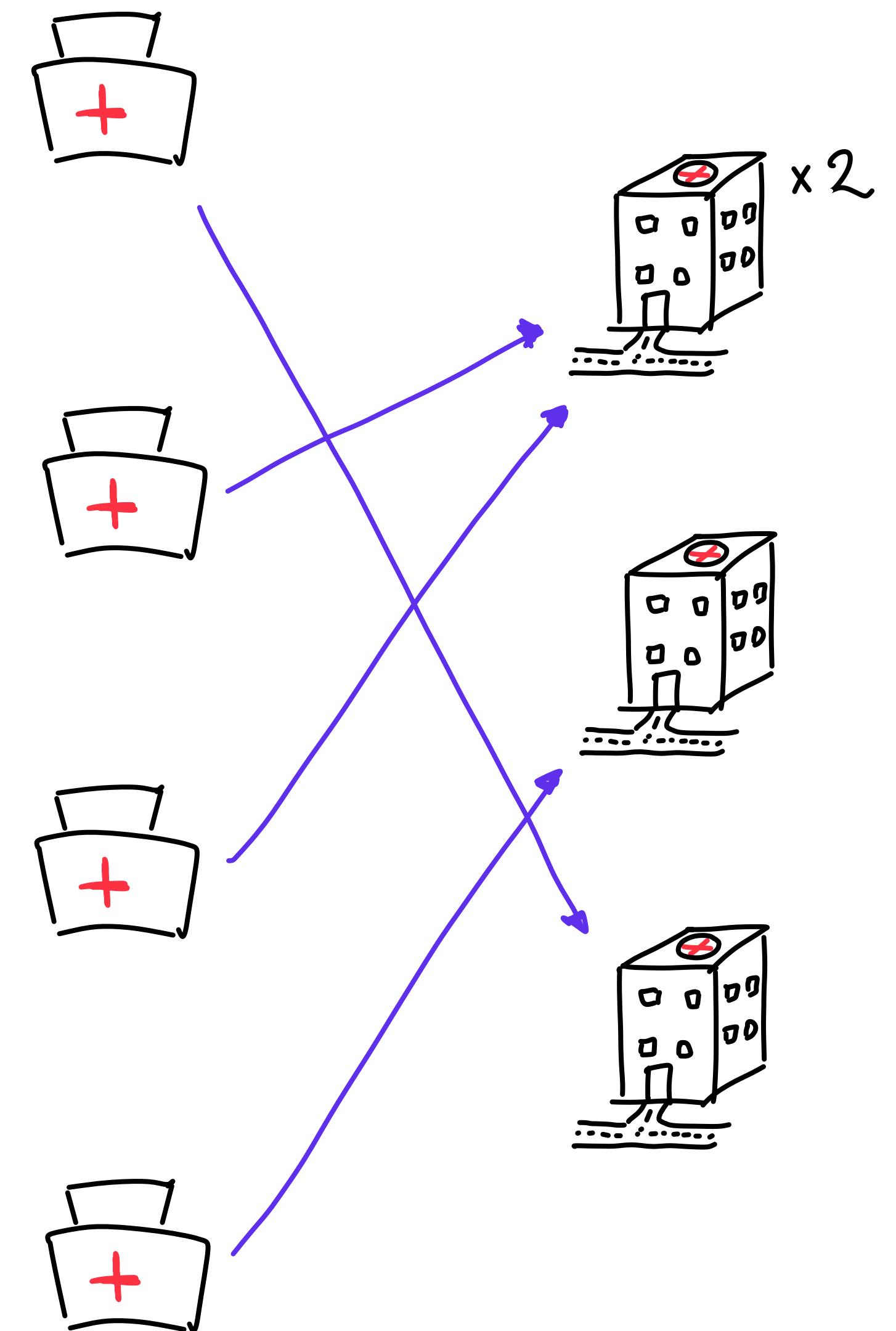
- **Mergesort**
 - Used when data is expressed as a linked list and RAM access to entries in the middle of the list is non-existent
 - Sorting over a dataset that cannot be stored in memory
 - Uses $O(n)$ extra space when sorting arrays over Quicksort

Sorting in the real world

- **Insertion sort**
 - Best when data is almost sorted already
 - $O(n^2)$ when far from sorted
- **Heap sort** - memory efficient choice
- **Bucket sort** - distribution aware sorting
- Etc...

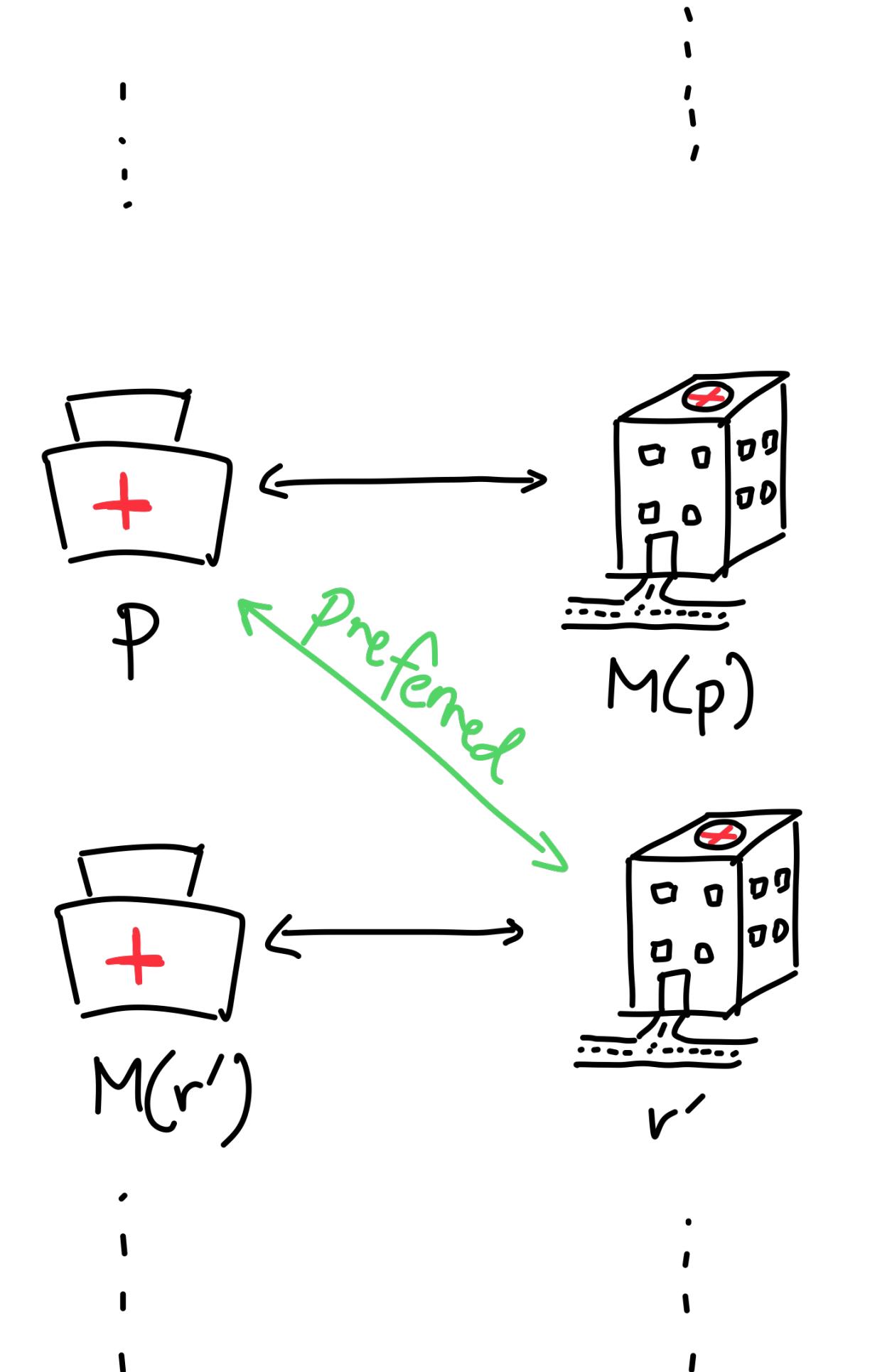
The matching problem

- **Goal:** Given a set of preferences amongst hospital and residents, design an admissions process to allocate residents to hospitals.
- What might we want to optimize for?
- When do we know we have achieved the optimal solution?
- What properties does our optimal solution have?



A notion of *stability*

- Lets assume there are n residents and n hospitals for now.
- A matching M is n disjoint pairs (p, r) assigning hospital r to resident p .
- A resident-hospital pair (resident p , hospital r') is **unstable** for M if both
 - resident p prefers hospital r' to their assigned hospital $M(p)$.
 - hospital r' prefers resident p to their assigned resident $M(r')$.
- A matching is **stable** if the matching has no **unstable** pairs.
 - Natural and desirable condition. *Self-interest* will prevent side-deals from being made.



Cartoon of unstable (p, r')

Can we design an algorithm to find a stable matching?

And does a stable matching necessarily exist?

- **Input to the problem:**

- Two groups of n people: one group P and the other group R .
- For each $p \in P$, a ranking from 1 to n of the group R .
- For each $r \in R$, a ranking from 1 to n of the group P .

- **Output of the problem:**

- A list of n disjoint pairs M . The matching should be stable with respect to the input rankings.

	1 st	2 nd	3 rd
X	A	B	C
Y	B	A	C
Z	A	B	C

	1 st	2 nd	3 rd
A	Y	X	Z
B	X	Y	Z
C	X	Y	Z

Example 1: Is the following matching stable?

favorite ↴ least fav. ↴

$X \leftrightarrow C$
 $Y \leftrightarrow B$
 $Z \leftrightarrow A$

	1 st	2 nd	3 rd
X	A	B	C
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Z	A	B	C

	1 st	2 nd	3 rd
A	Y	X	Z
B	X	Y	Z
C	X	Y	Z

Example 1: Is the following matching stable?

No.

	1 st	2 nd	3 rd
X	A	B	C
Y	B	A	C
Z	A	B	C

favorite

least fav.

$X \leftrightarrow C$
 $Y \leftrightarrow B$
 $Z \leftrightarrow A$

	1 st	2 nd	3 rd
A	Y	X	Z
B	X	Y	Z
C	X	Y	Z

mutually preferred
change

Example 2: Is the following matching stable?

favorite ↴ least fav. ↴

	1 st	2 nd	3 rd
X	A	B	C
Y	B	A	C
Z	A	B	C

$X \leftrightarrow A$
 $Y \leftrightarrow B$
 $Z \leftrightarrow C$

	1 st	2 nd	3 rd
A	Y	X	Z
B	X	Y	Z
C	X	Y	Z

Example 2: Is the following matching stable?

YES.

favorite ↴ least fav. ↴

	1 st	2 nd	3 rd
X	A	B	C
Y	B	A	C
Z	A	B	C

	1 st	2 nd	3 rd
A	Y	X	Z
B	X	Y	Z
C	X	Y	Z

X \longleftrightarrow A
Y \longleftrightarrow B
Z \longleftrightarrow C

The propose and reject algorithm

Gale & Shapley 1962

The group P proposes and the group R receives



```
Initialize each person to be free.

while (some p in P is free) {
    Choose some free p in P
    r = 1st person on p's preference list to whom p has not yet proposed
    if (r is free)
        tentatively match (p,r)    //p and r both engaged, no longer free
    else if (r prefers p to current tentative match p')
        replace (p',r) by (p,r)    //p now engaged, p' now free
    else
        r rejects p
}
```

Gale-Shapley walkthrough

$$n = 4.$$

We will walkthrough alg, staying blind to the remainder of the input until we have queried it.

FAV ↓ LEAST ↓

ALPHA			
BRAVO			
CHARLIE			
DELTA			

FAV ↓ LEAST ↓

PAPA			
QUEBEC			
ROMEO			
SIERRA			

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Gale-Shapley walkthrough

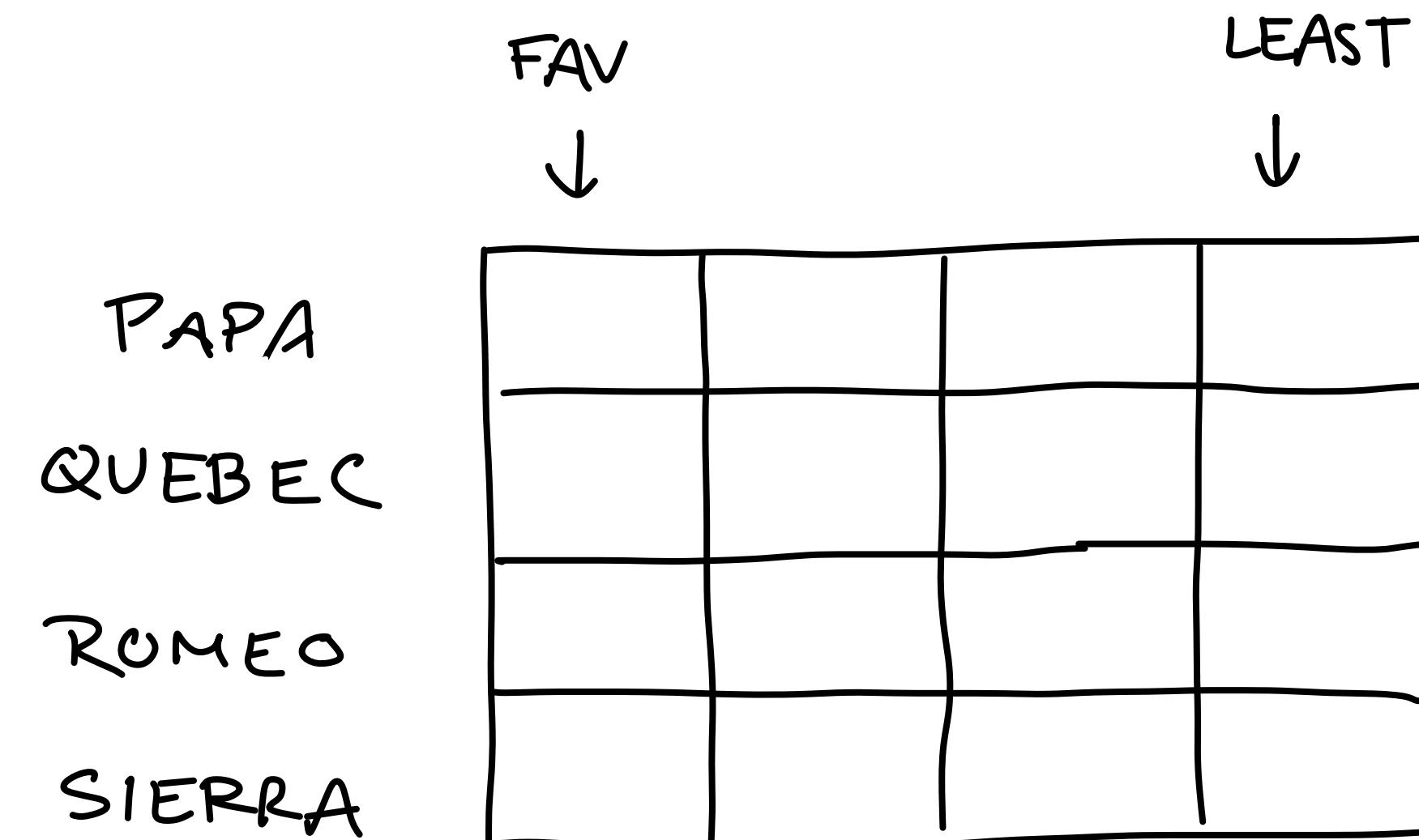
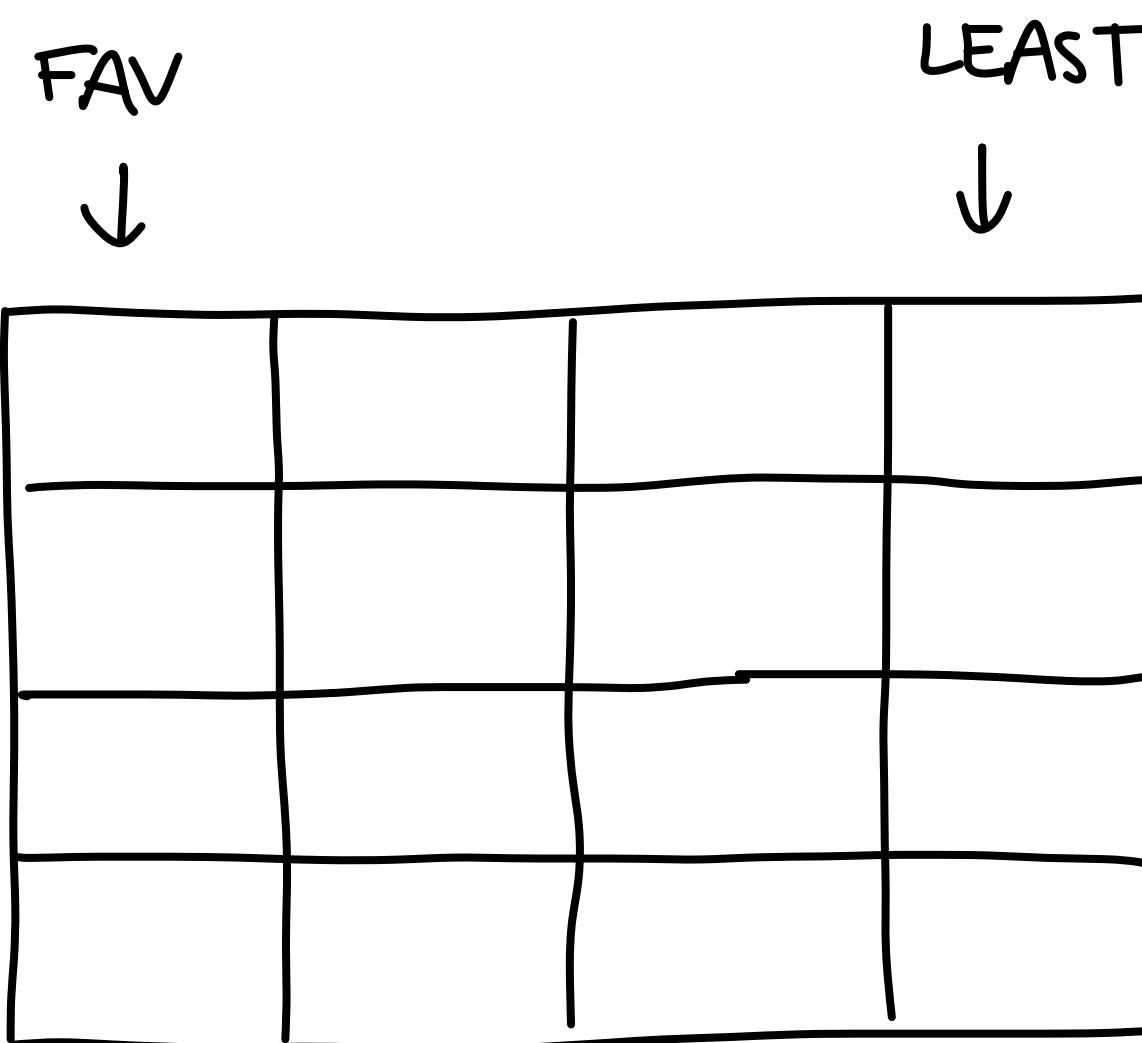
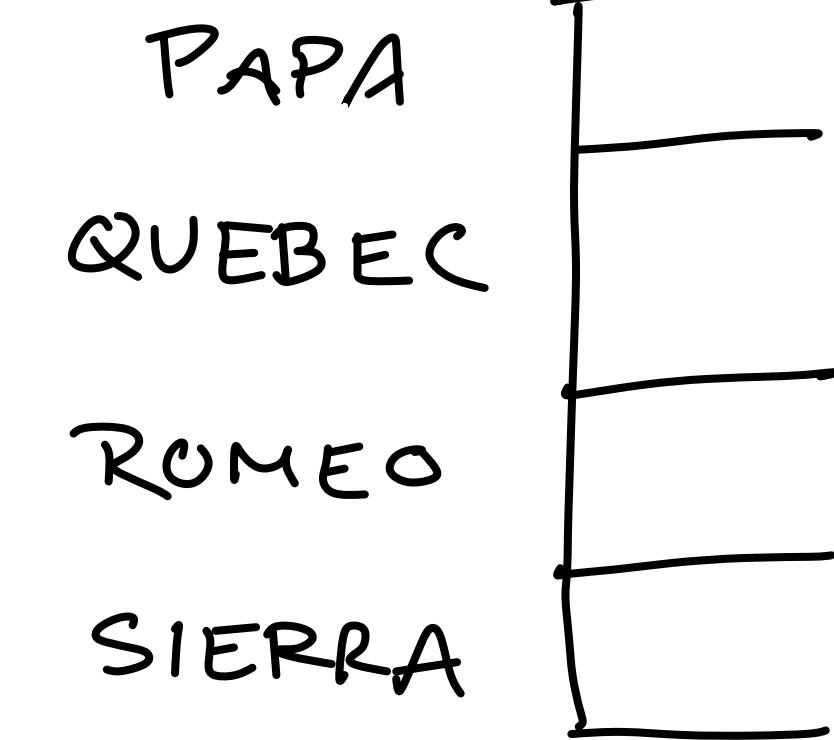
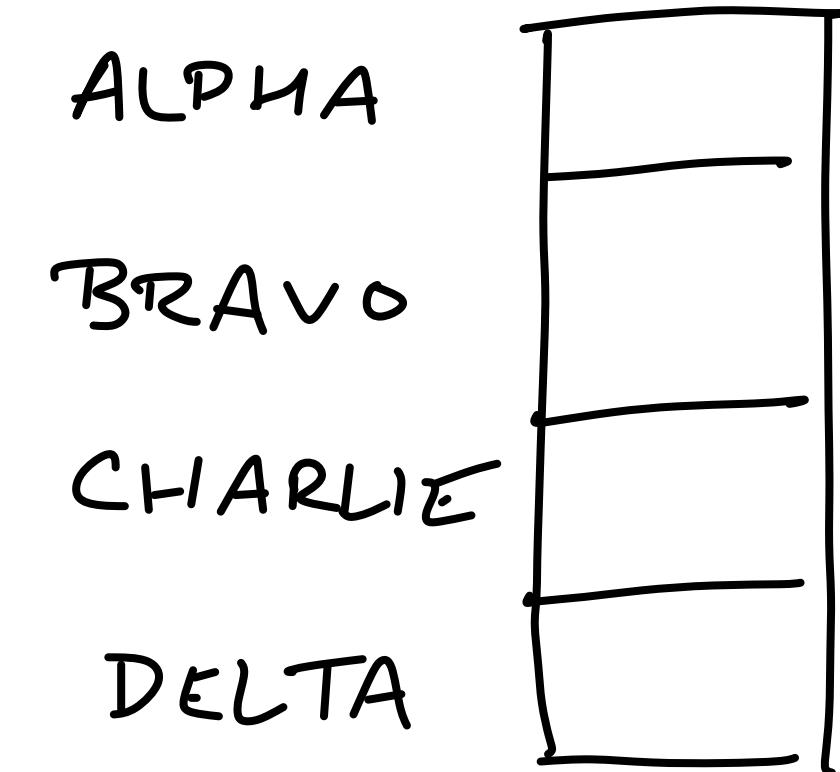
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Current partner:



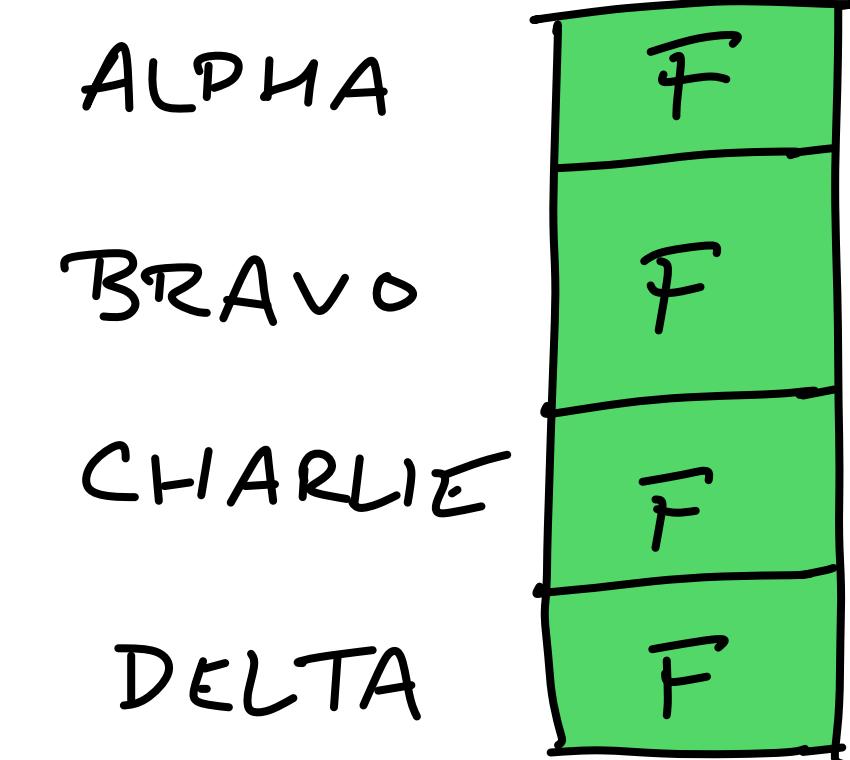
Gale-Shapley walkthrough

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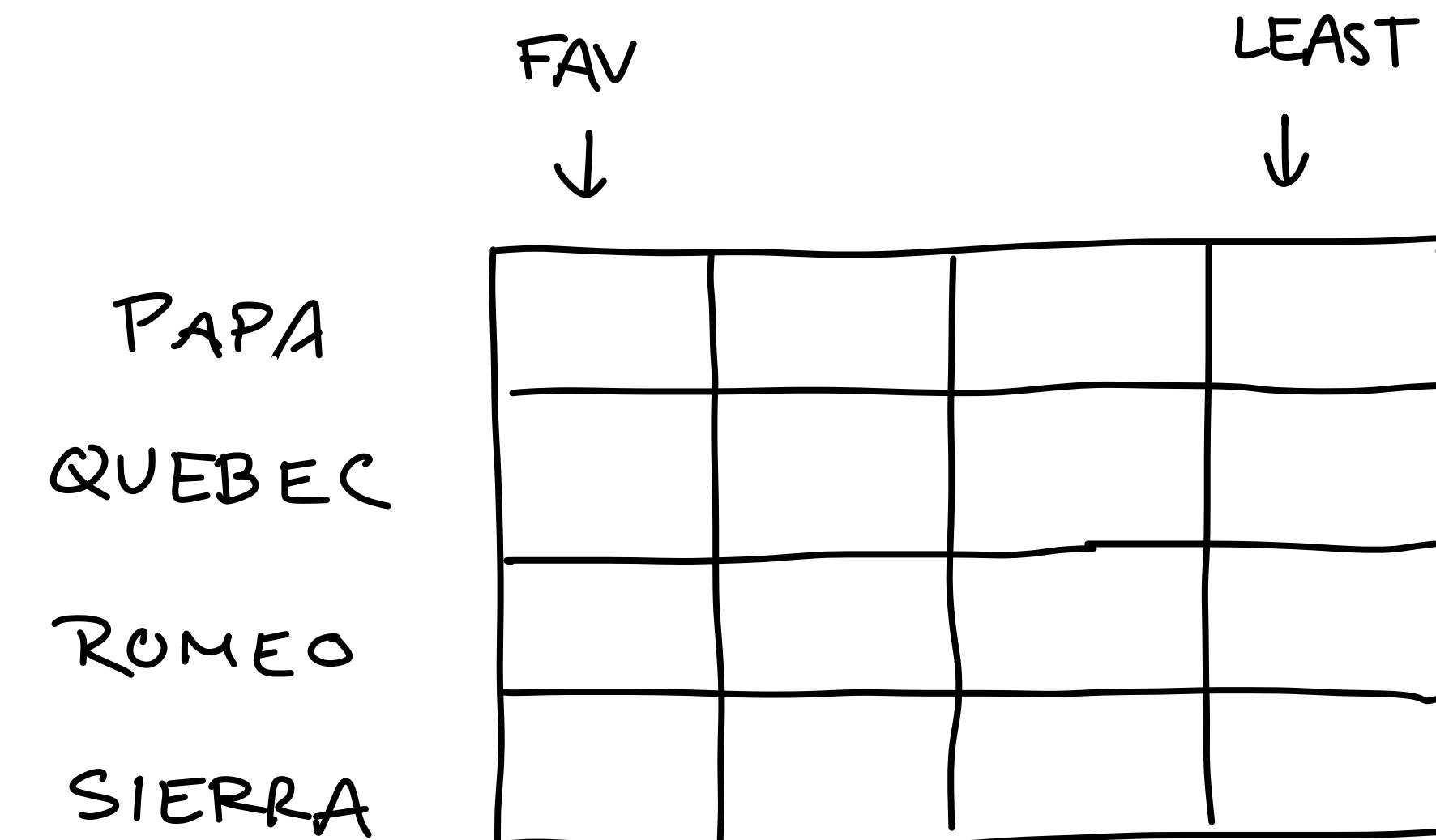
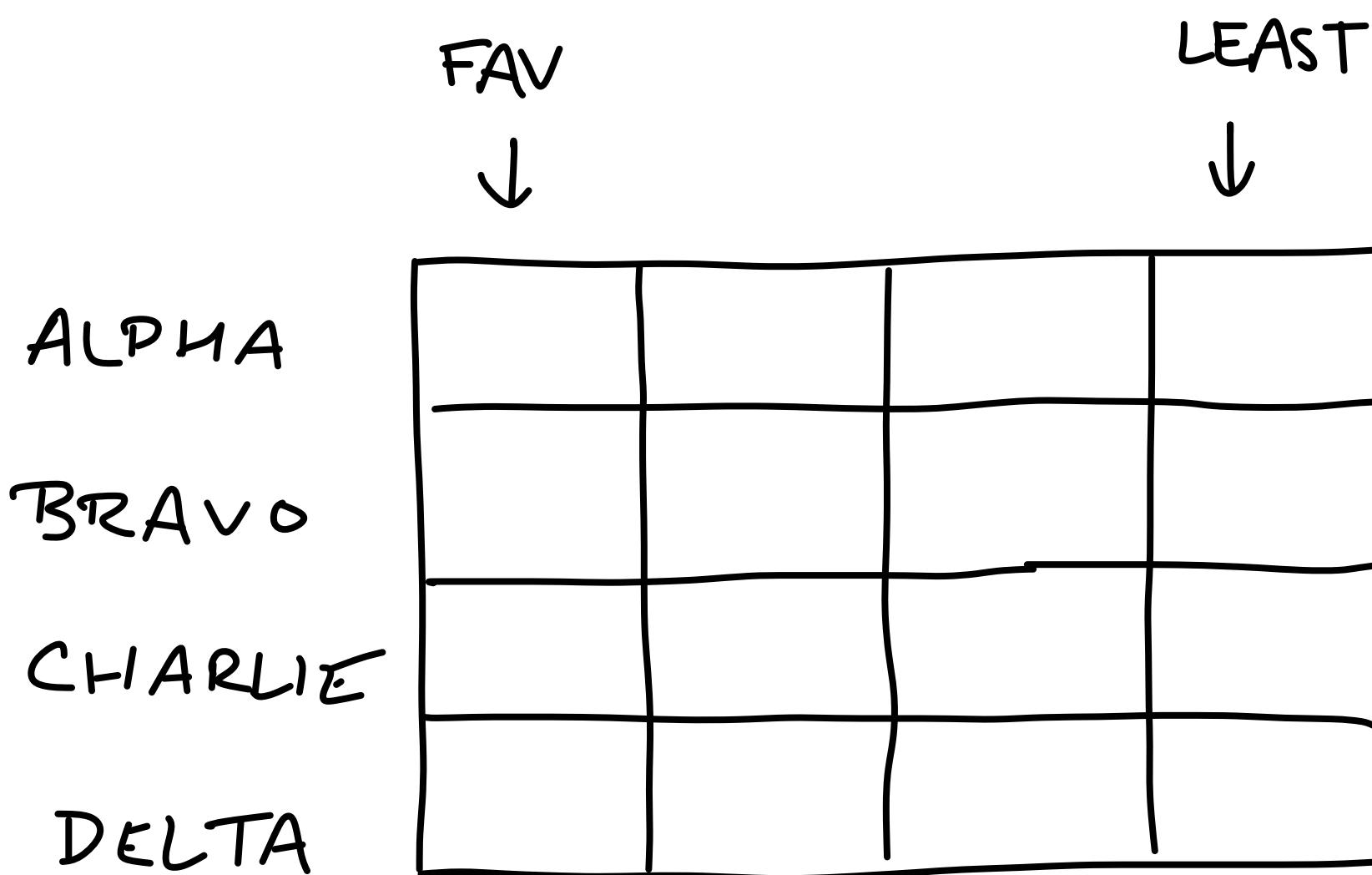
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Current partner:



PAPA QUEBEC ROMEO SIERRA



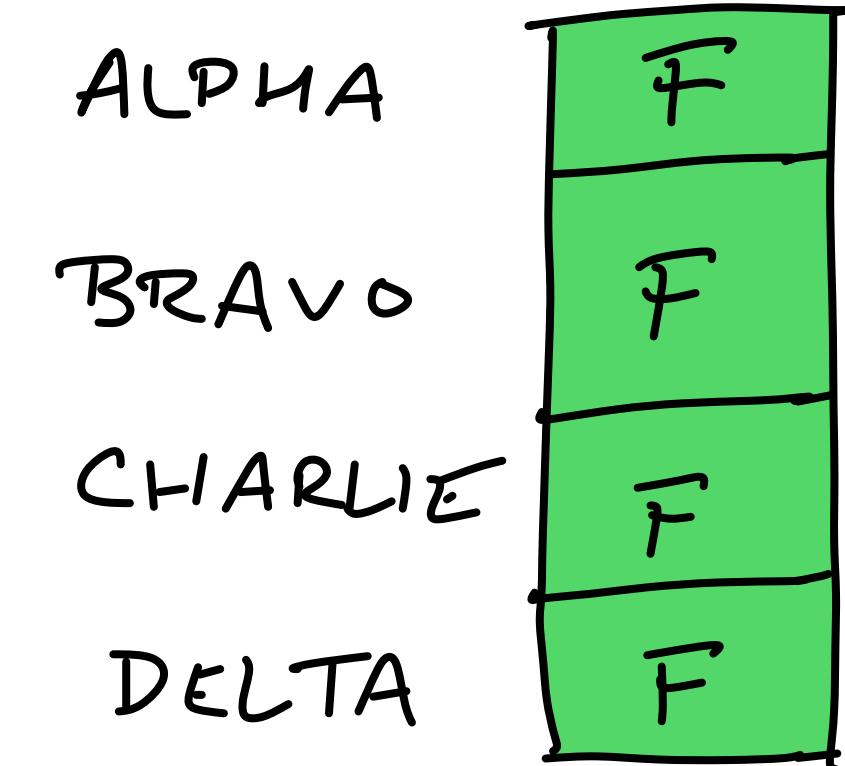
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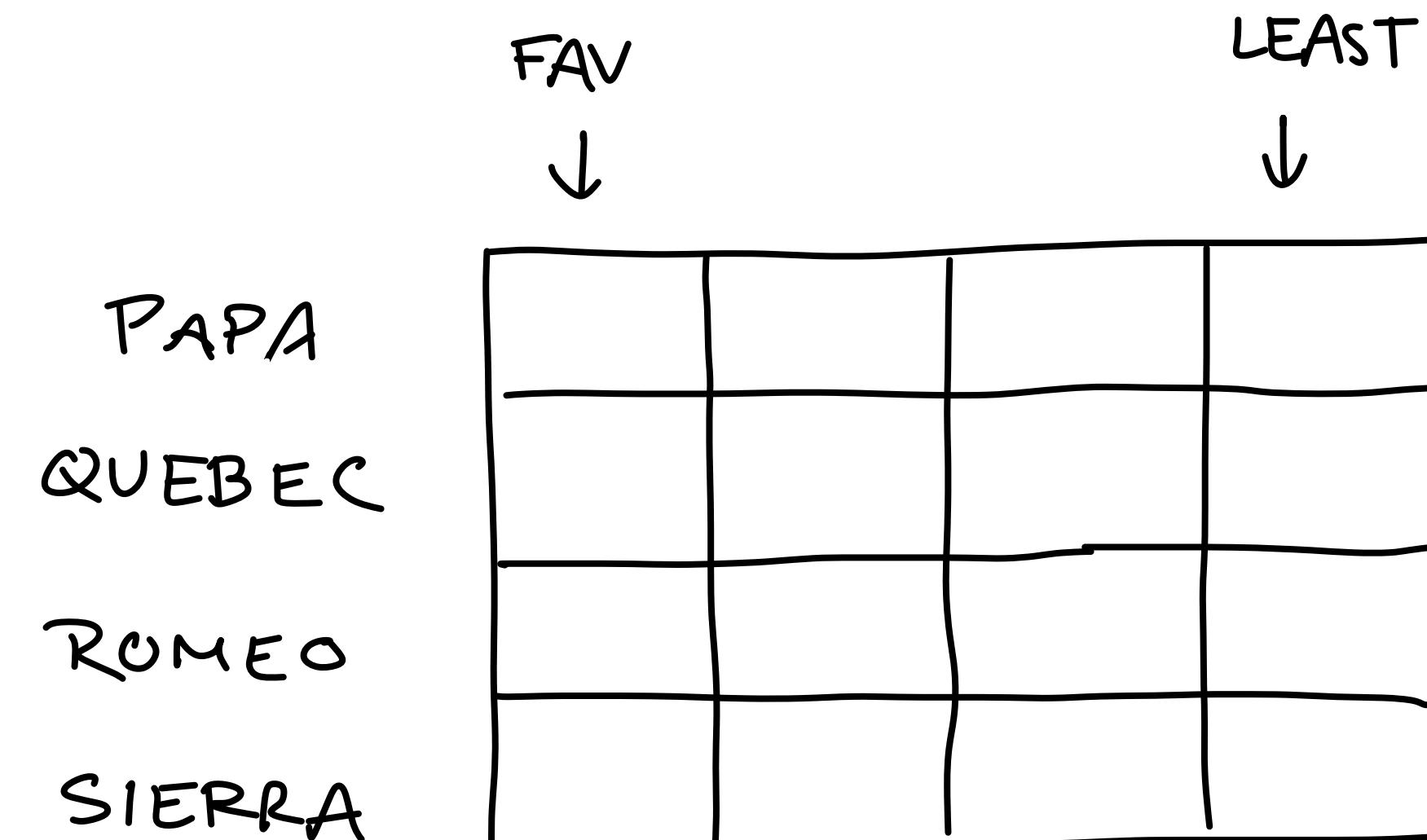
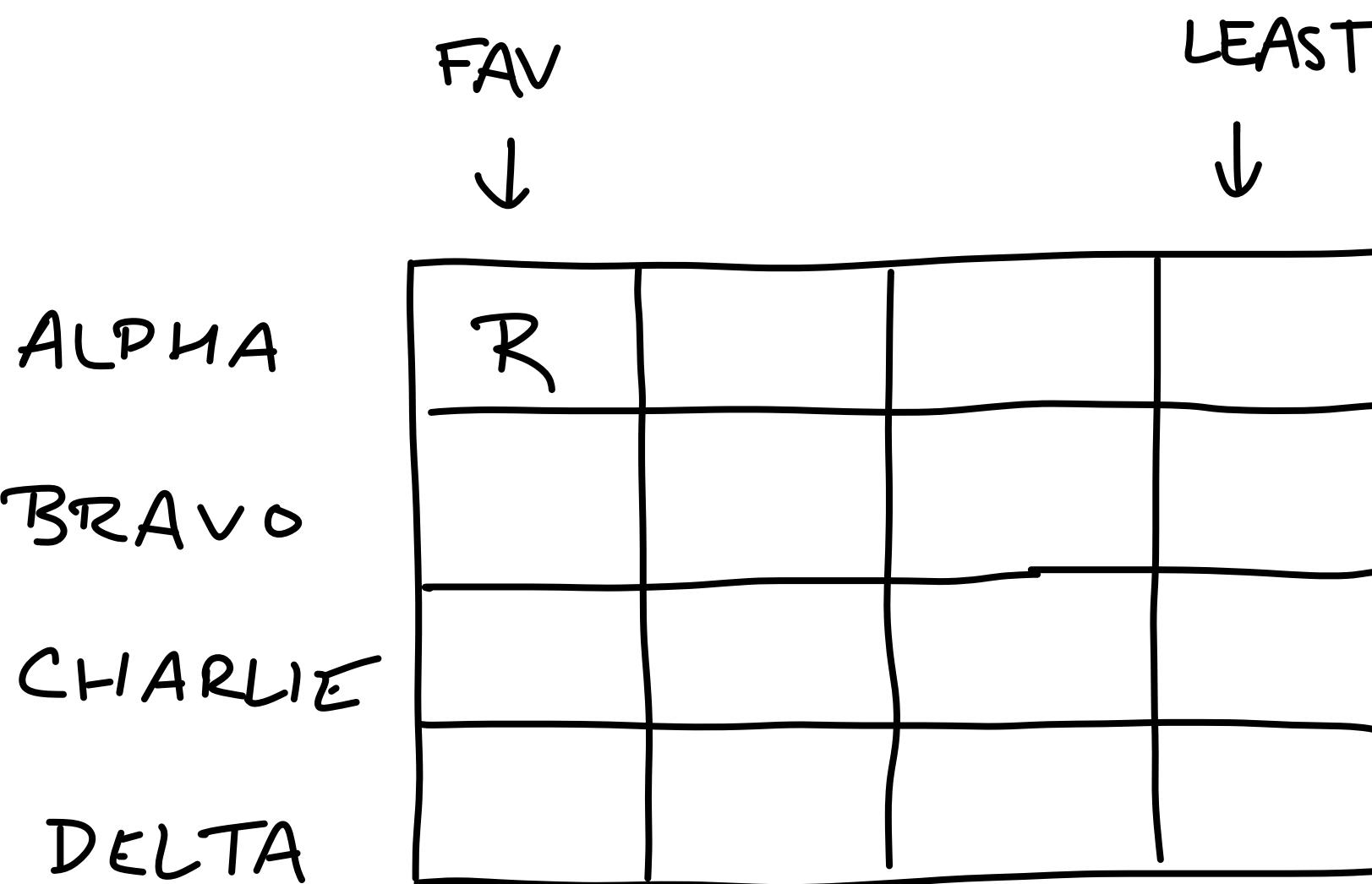
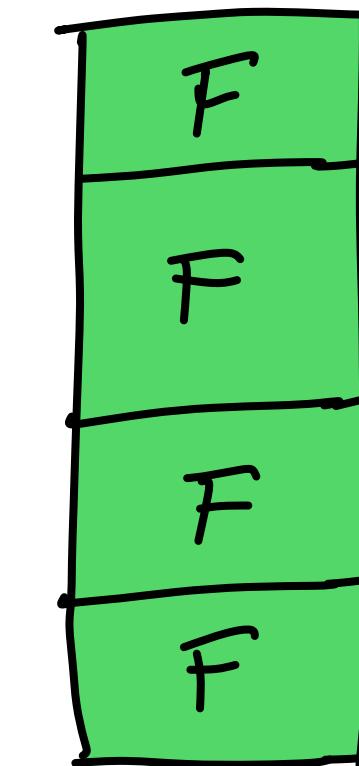
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PAPA QUEBEC ROMEO SIERRA



Gale-Shapley walkthrough

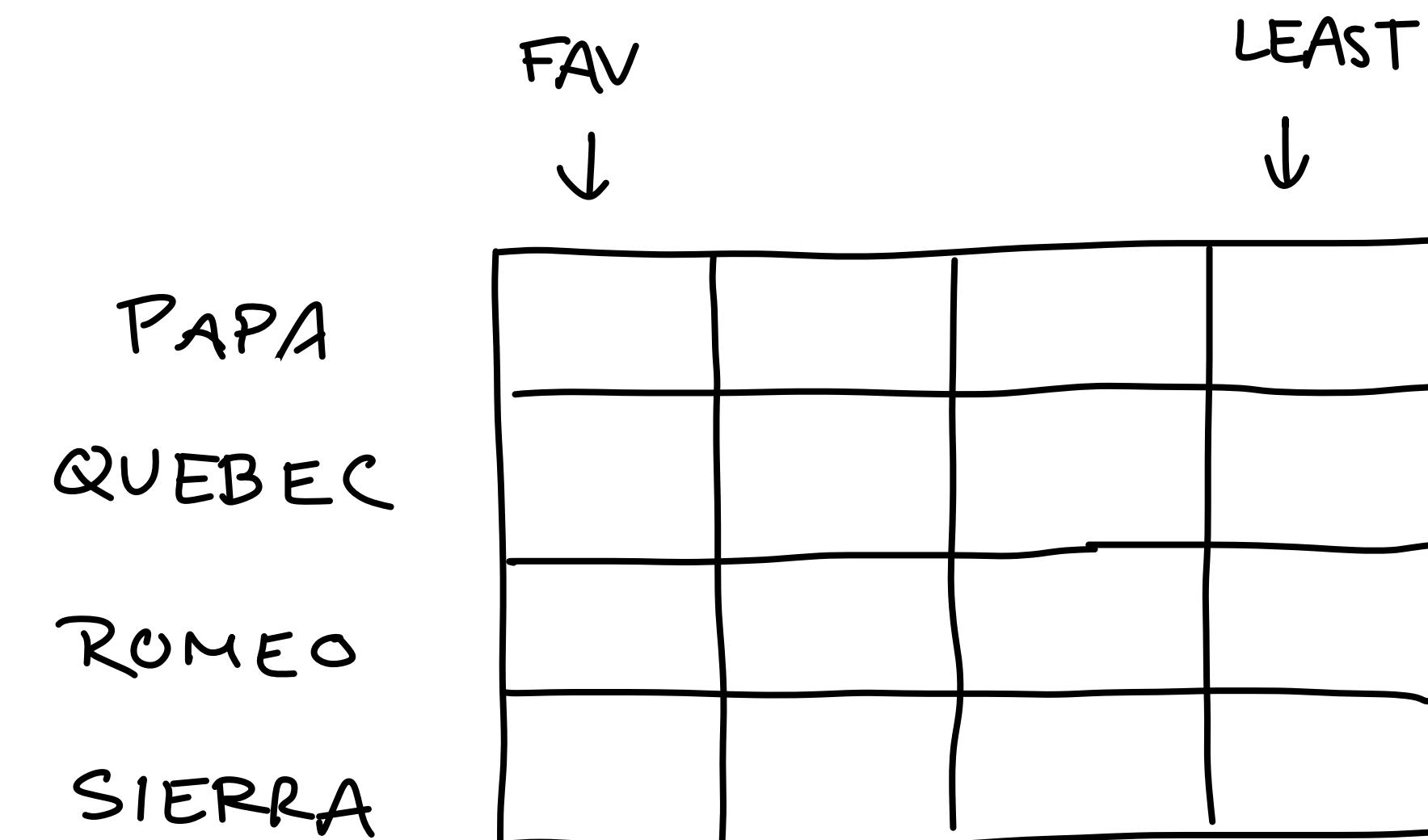
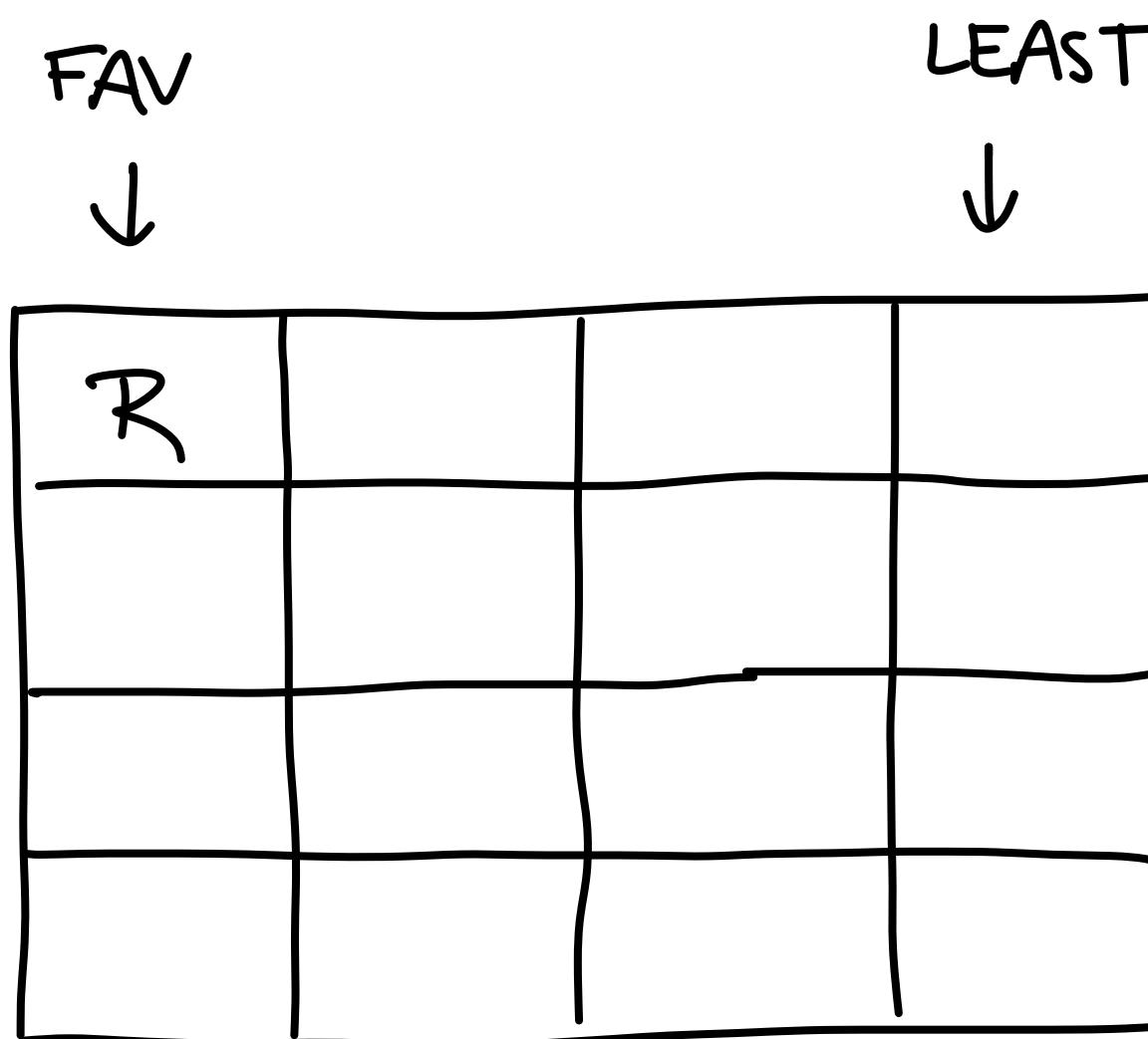
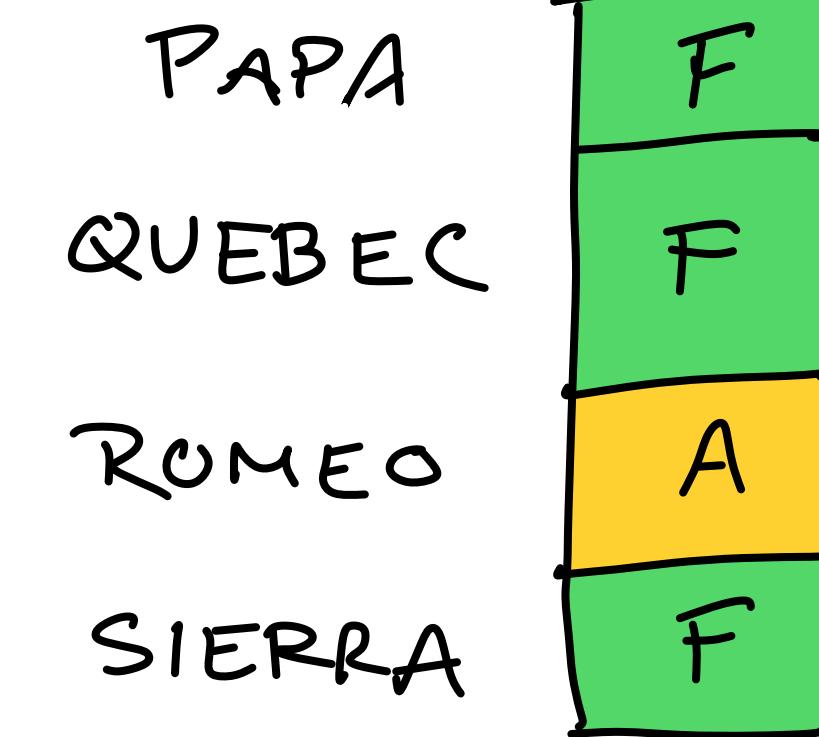
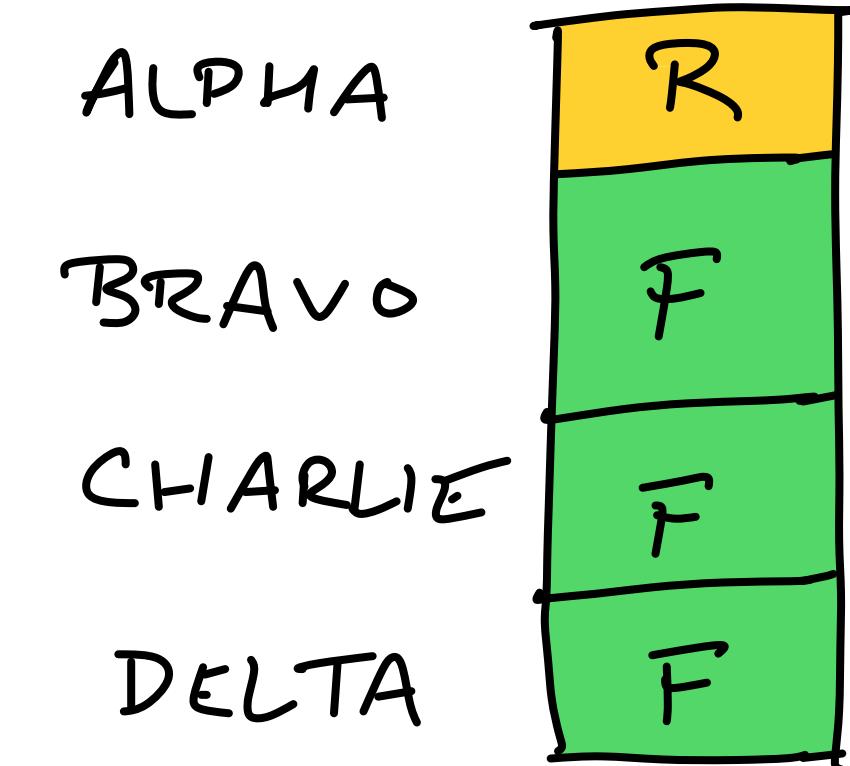
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Current partner:



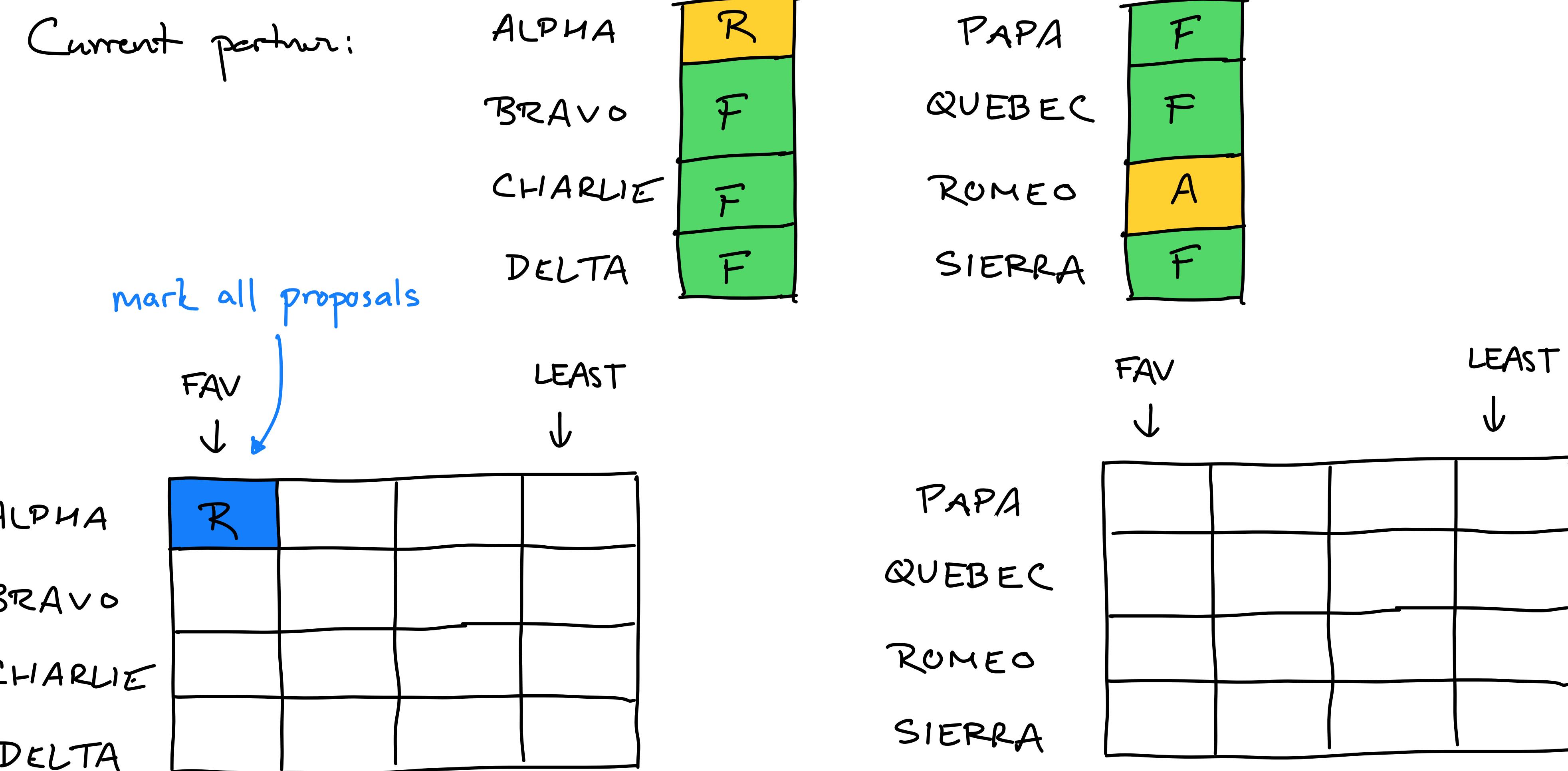
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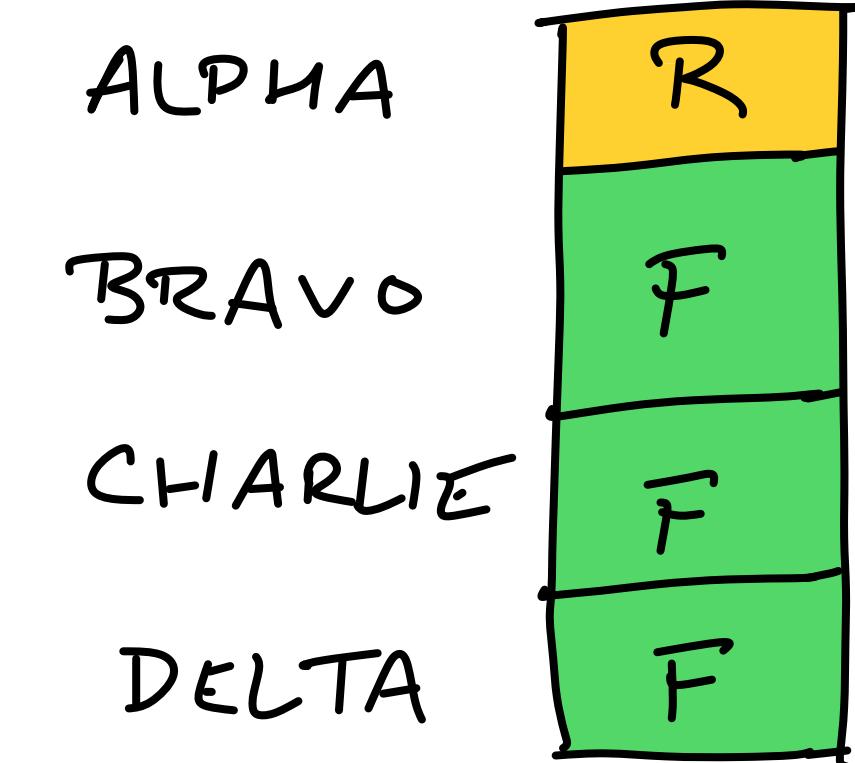
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Current partner:



mark all proposals

FAV



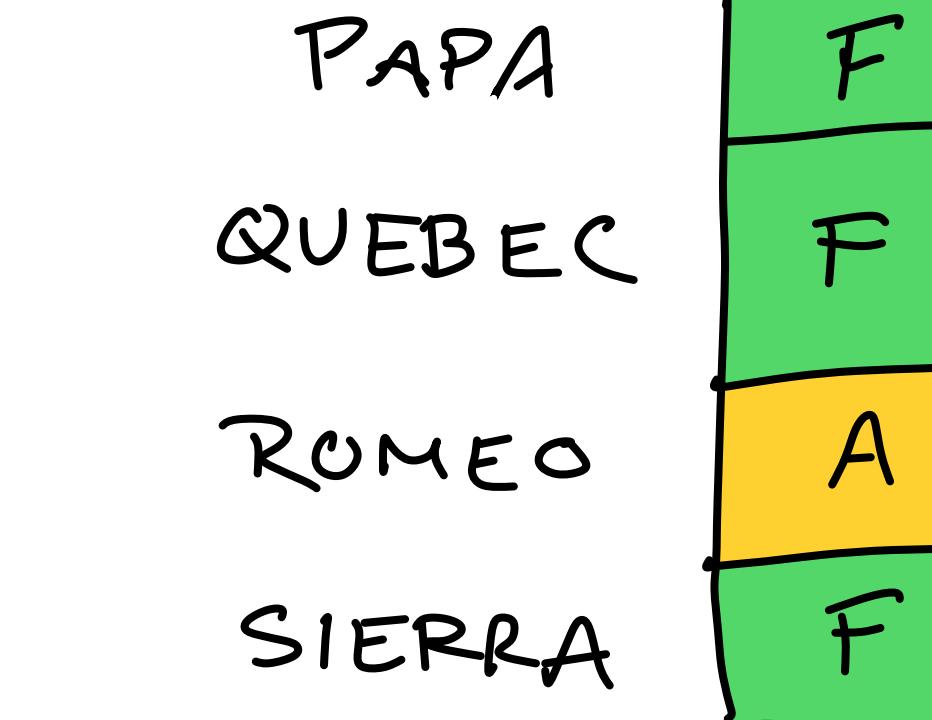
LEAST



ALPHA

R			
Q			

BRAVO



FAV



LEAST



PAPA

CHARLIE

QUEBEC

DELTA

ROMEO

SIERRA

Gale-Shapley walkthrough

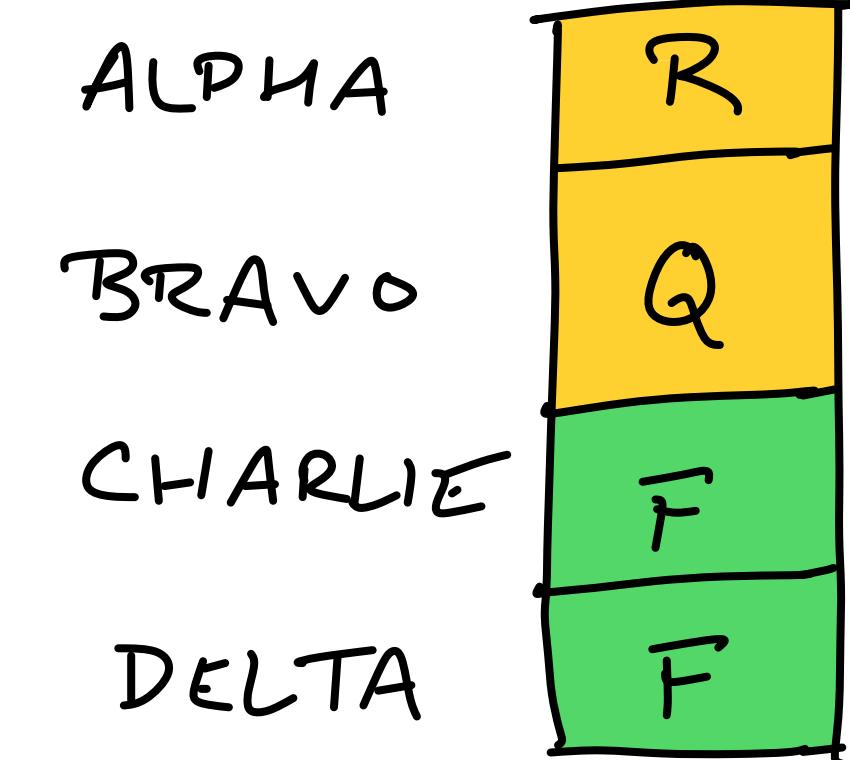
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        replace (p',r) by (p,r) // p now engaged, p' now free
    else
        r rejects p
}

```

Current partner:



mark all proposals

FAV

LEAST

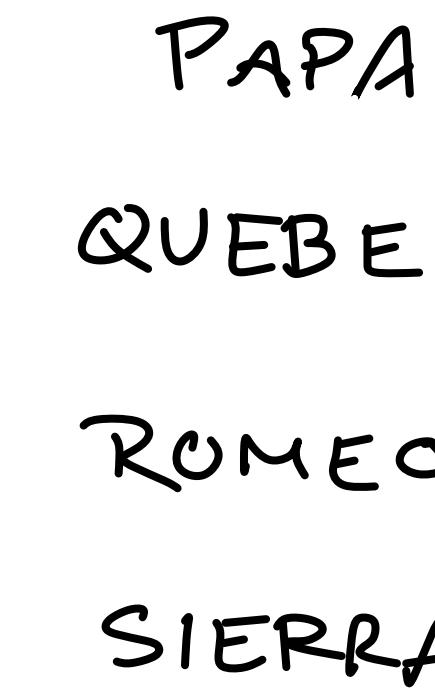
ALPHA

BRAVO

CHARLIE

DELTA

ALPHA	R		
	Q		



FAV

LEAST

PAPA

QUEBEC

ROMEO

SIERRA

Gale-Shapley walkthrough

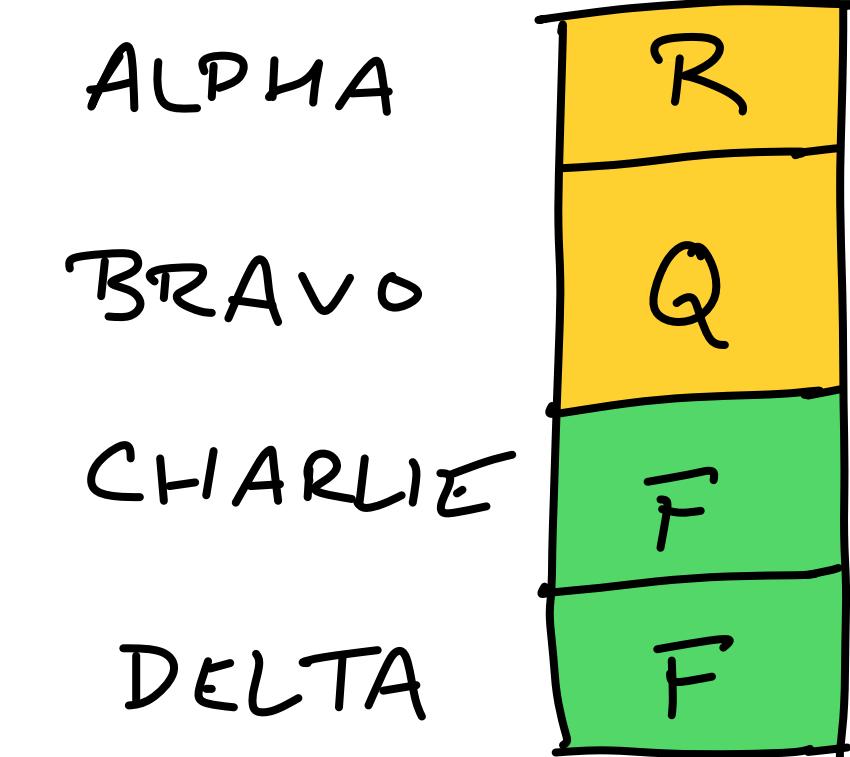
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}

```

Current partner:



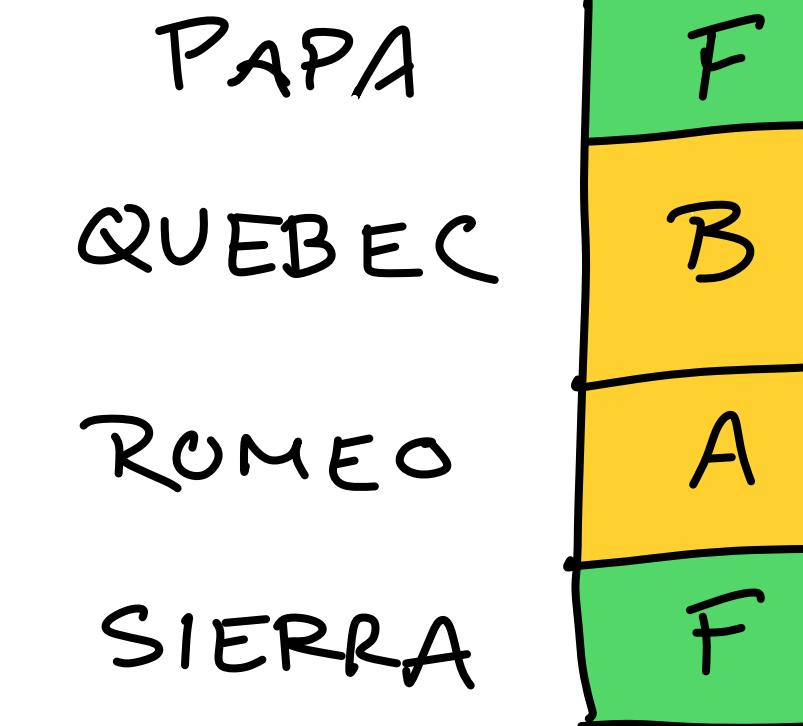
mark all proposals

FAV
↓

LEAST
↓

ALPHA
BRAVO
CHARLIE
DELTA

ALPHA	R			
BRAVO	Q			
CHARLIE	Q			
DELTA				



FAV
↓

LEAST
↓

PAPA
QUEBEC
ROMEO
SIERRA

Gale-Shapley walkthrough

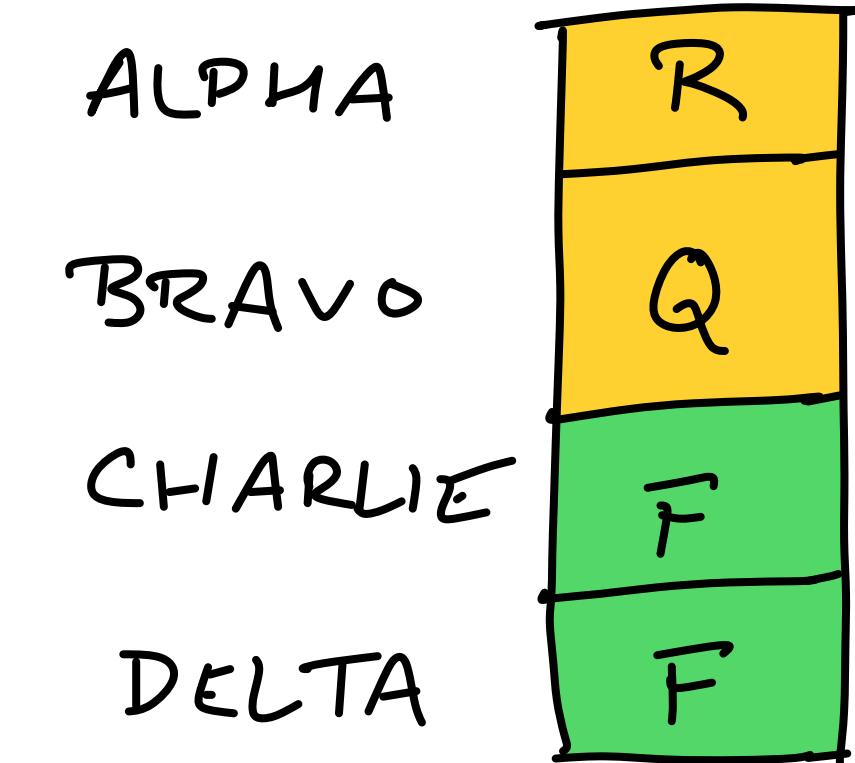
```

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```

Current partner:



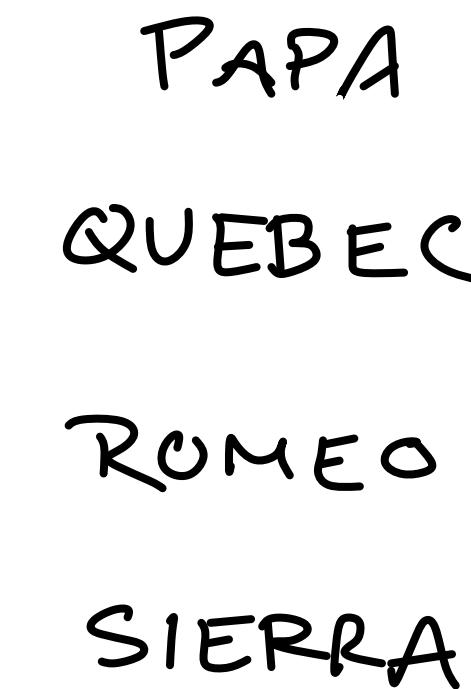
mark all proposals

FAV
↓

LEAST
↓

ALPHA
BRAVO
CHARLIE
DELTA

ALPHA	R			
BRAVO	Q			
CHARLIE	Q			
DELTA				



Who do I prefer:
BRAVO OR CHARLIE?

FAV LEAST
↓ ↓

PAPA
QUEBEC
ROMEO
SIERRA

.			
.			
.			
.			

Gale-Shapley walkthrough

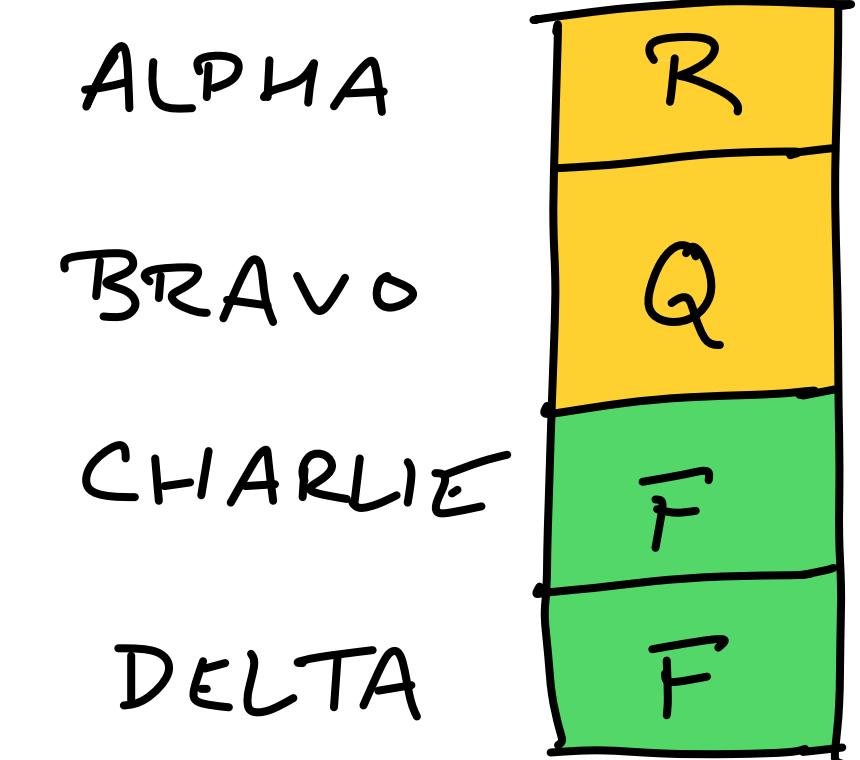
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}

```

Current partner:



mark all proposals

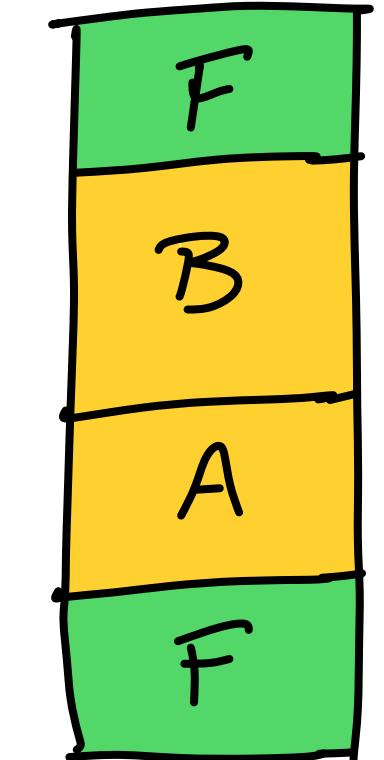
FAV
↓

LEAST
↓

ALPHA
BRAVO
CHARLIE
DELTA

ALPHA	R			
BRAVO	Q			
CHARLIE	Q			
DELTA				

PAPA QUEBEC ROMEO SIERRA



Who do I prefer:
BRAVO OR CHARLIE?

FAV LEAST
↓ ↓

PAPA
QUEBEC
ROMEO
SIERRA

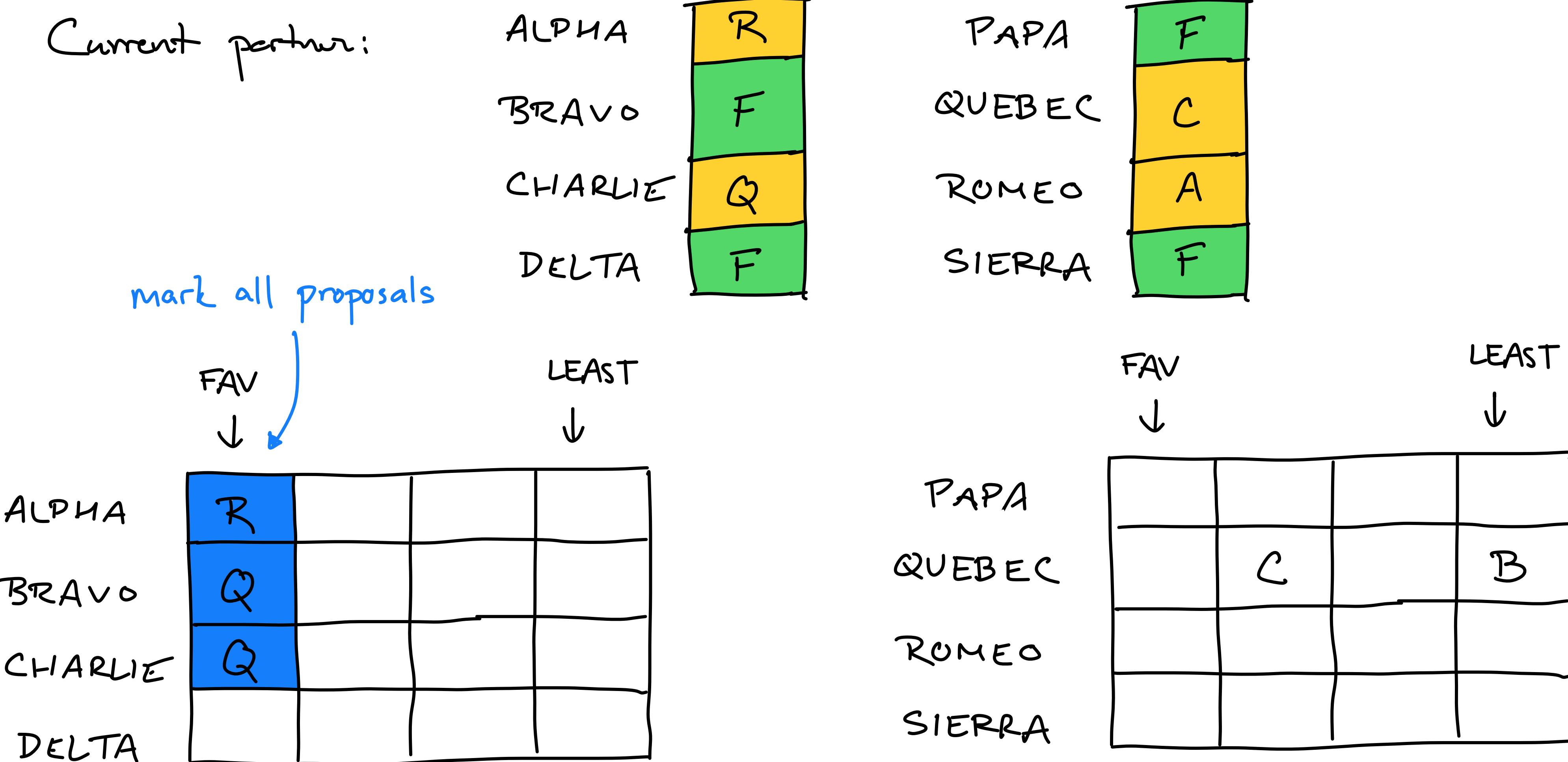
.			
.	C		B
.			
.			

Gale-Shapley walkthrough

```

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```



Gale-Shapley walkthrough

```

Initialize each person to be free.

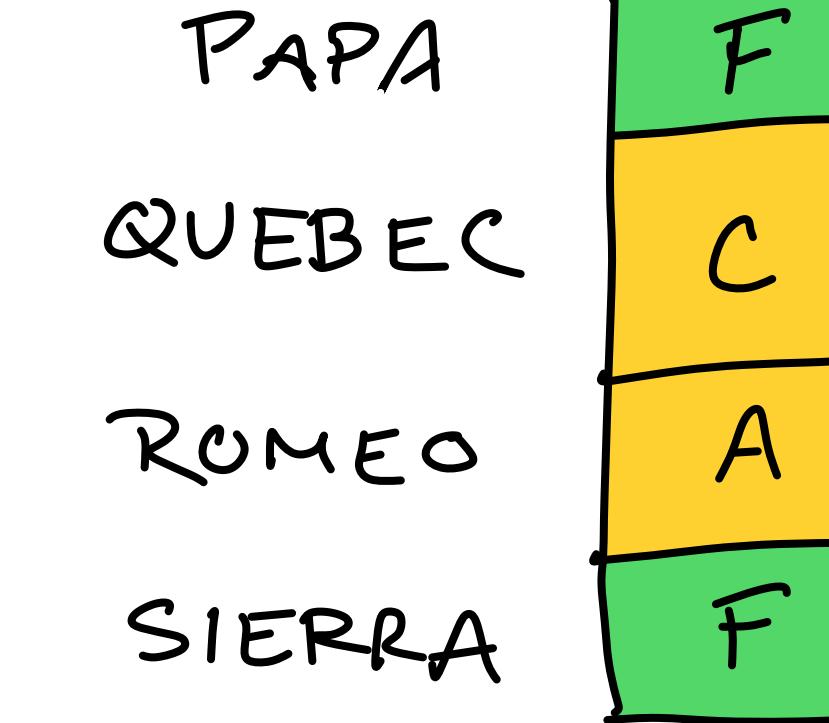
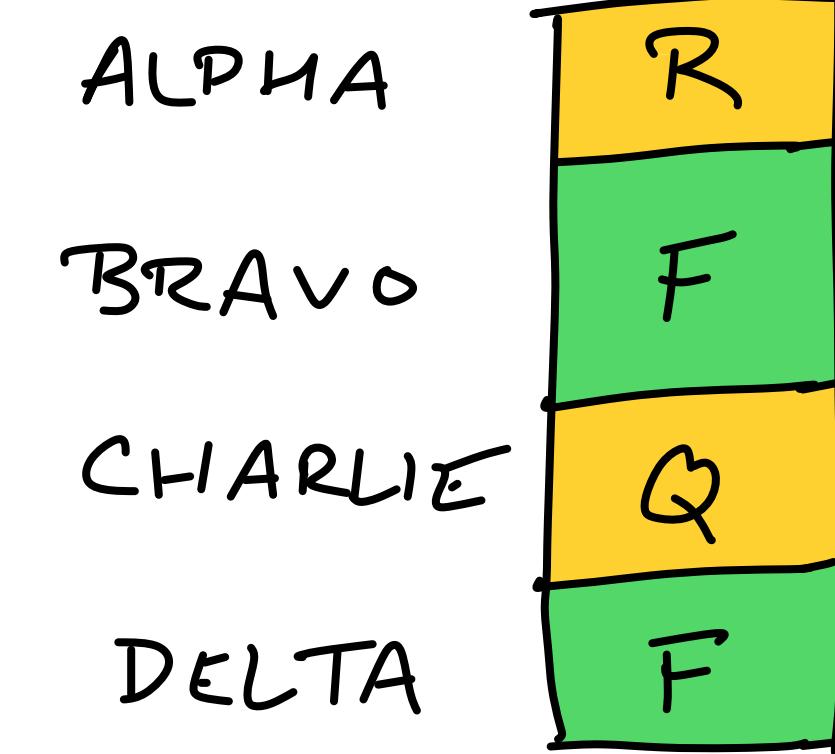
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```

Pick the next free proposer

How to pick?

Current partner:



mark all proposals

FAV

LEAST

ALPHA	R			
BRAVO	Q			
CHARLIE	Q			
DELTA				

FAV

LEAST

PAPA			
QUEBEC	C		B
ROMEO			
SIERRA			

Gale-Shapley walkthrough

```

Initialize each person to be free.

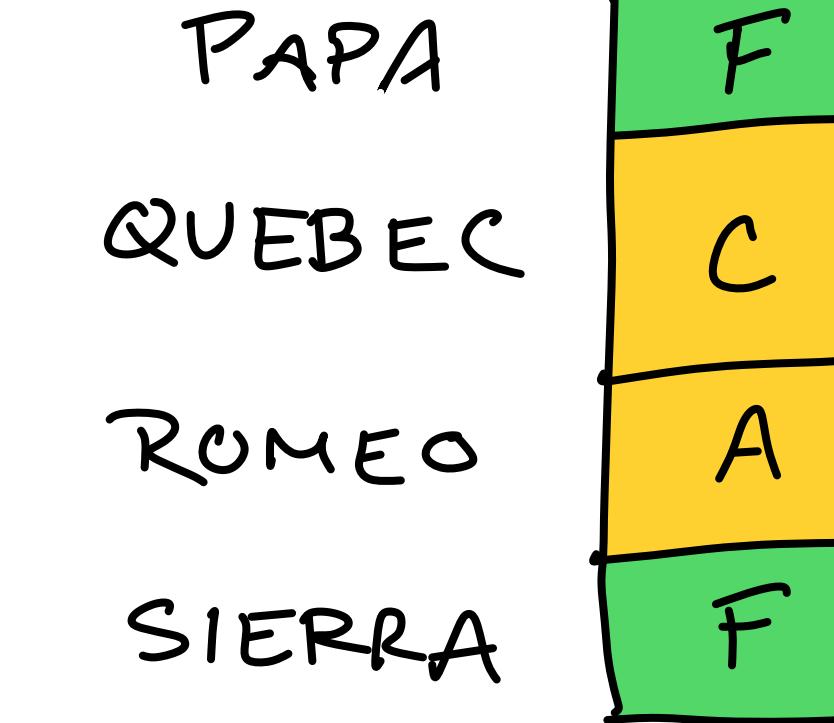
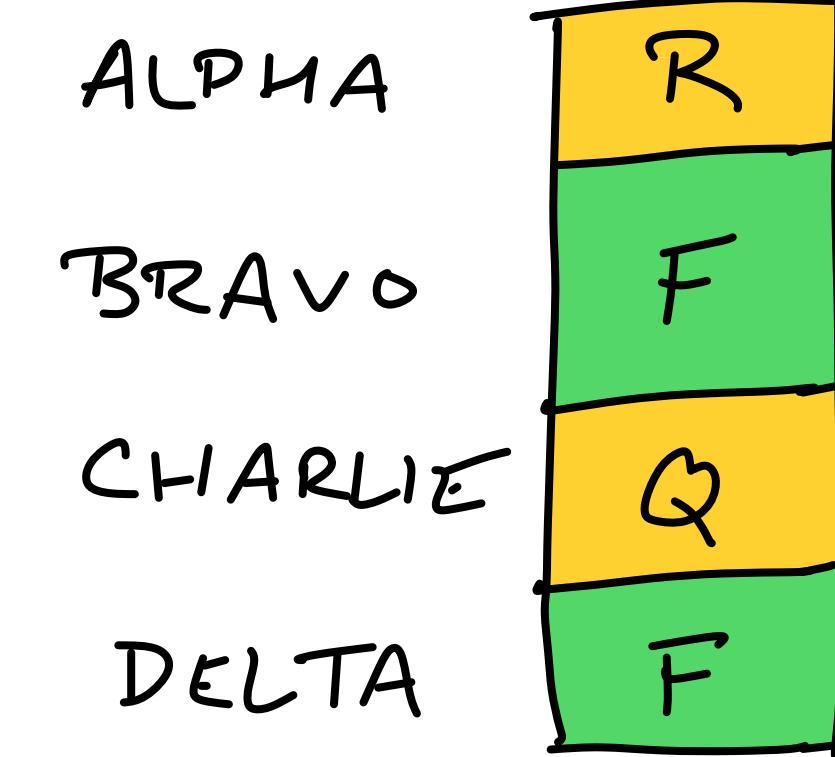
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}

```

Pick the next free proposer

How to pick?

Current partner:



mark all proposals

FAV

LEAST

ALPHA	R			
BRAVO	Q			
CHARLIE	Q			
DELTA	P			

FAV

LEAST

PAPA			
QUEBEC		C	B
ROMEO			
SIERRA			

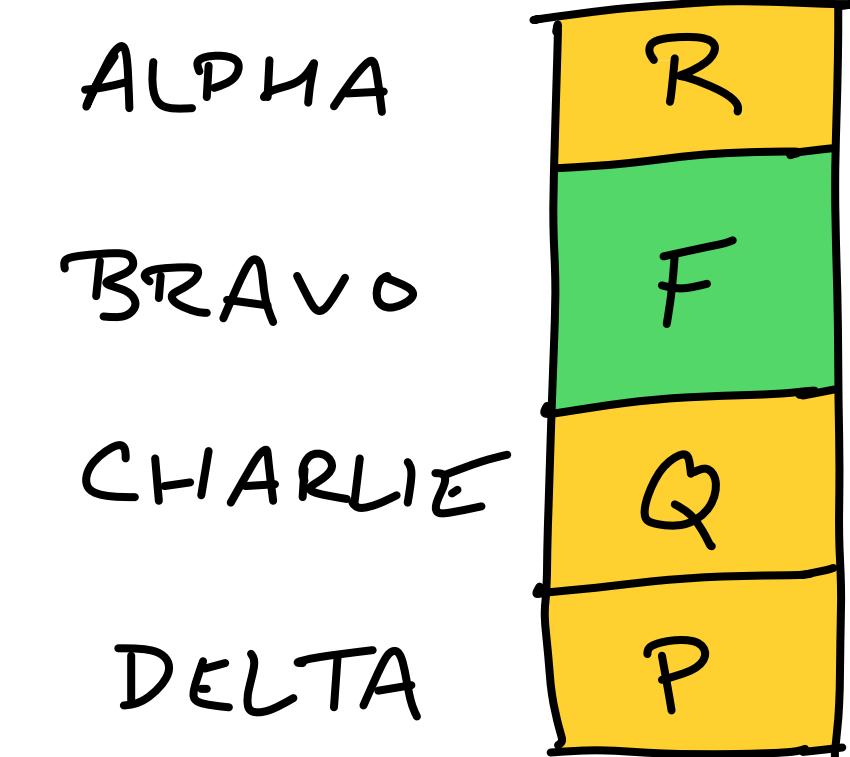
Gale-Shapley walkthrough

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Current partner:



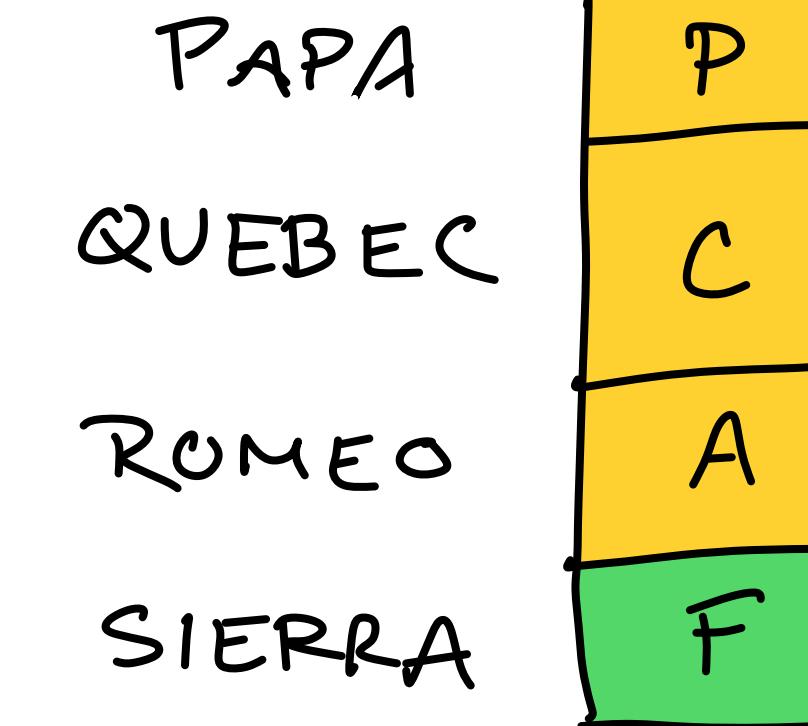
mark all proposals

FAV
↓

LEAST
↓

ALPHA
BRAVO
CHARLIE
DELTA

R			
Q			
Q			
P			



FAV
↓

LEAST
↓

PAPA
QUEBEC
ROMEO
SIERRA

	C		B

Gale-Shapley walkthrough

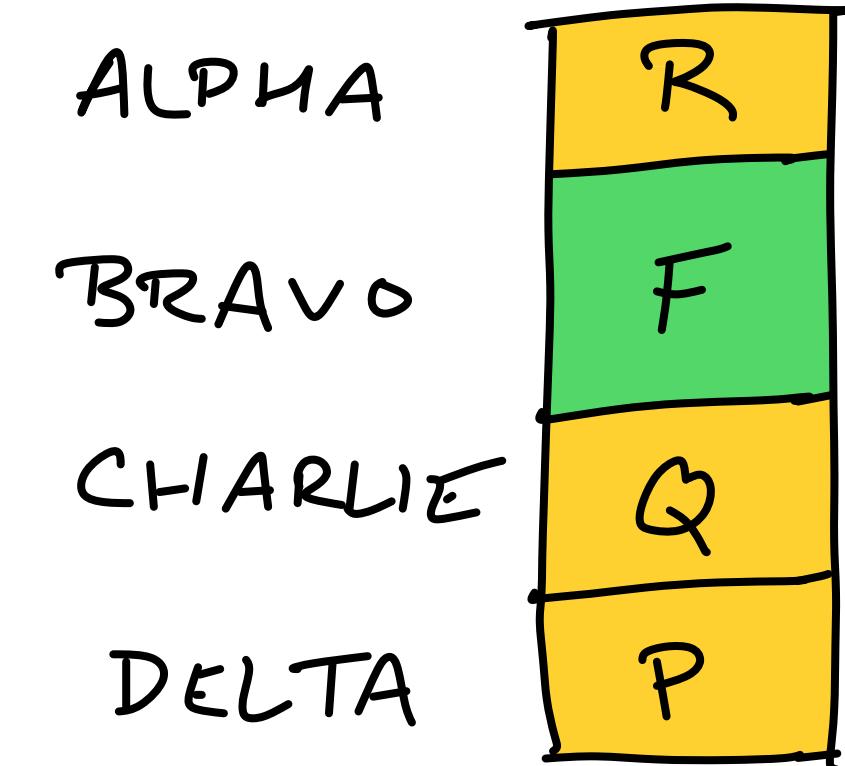
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Current partner:



mark all proposals

FAV



LEAST



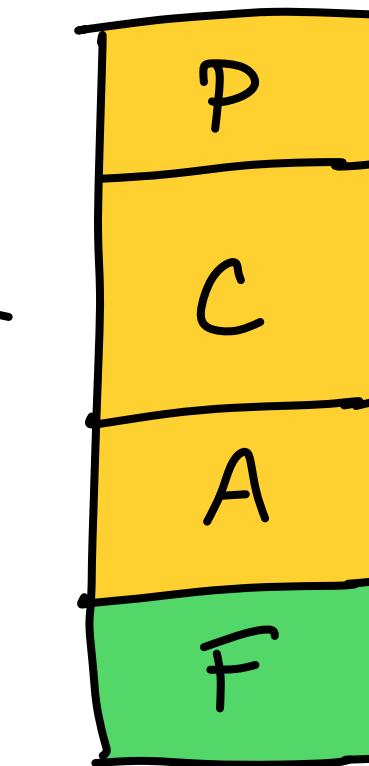
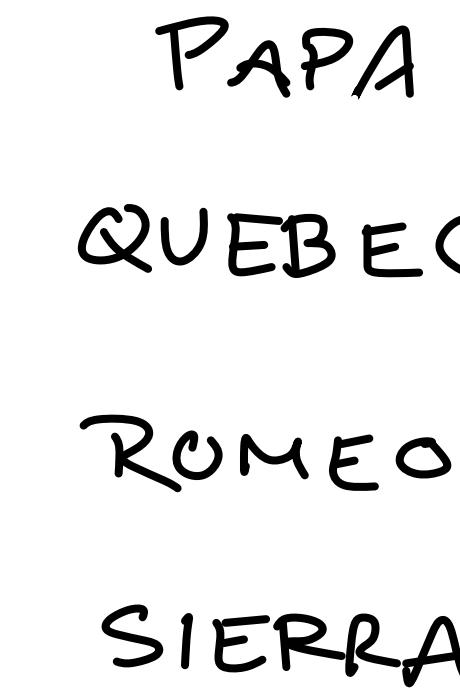
ALPHA

BRAVO

CHARLIE

DELTA

R			
Q	P		
Q			
P			



FAV



LEAST



PAPA

QUEBEC

ROMEO

SIERRA

	C		B

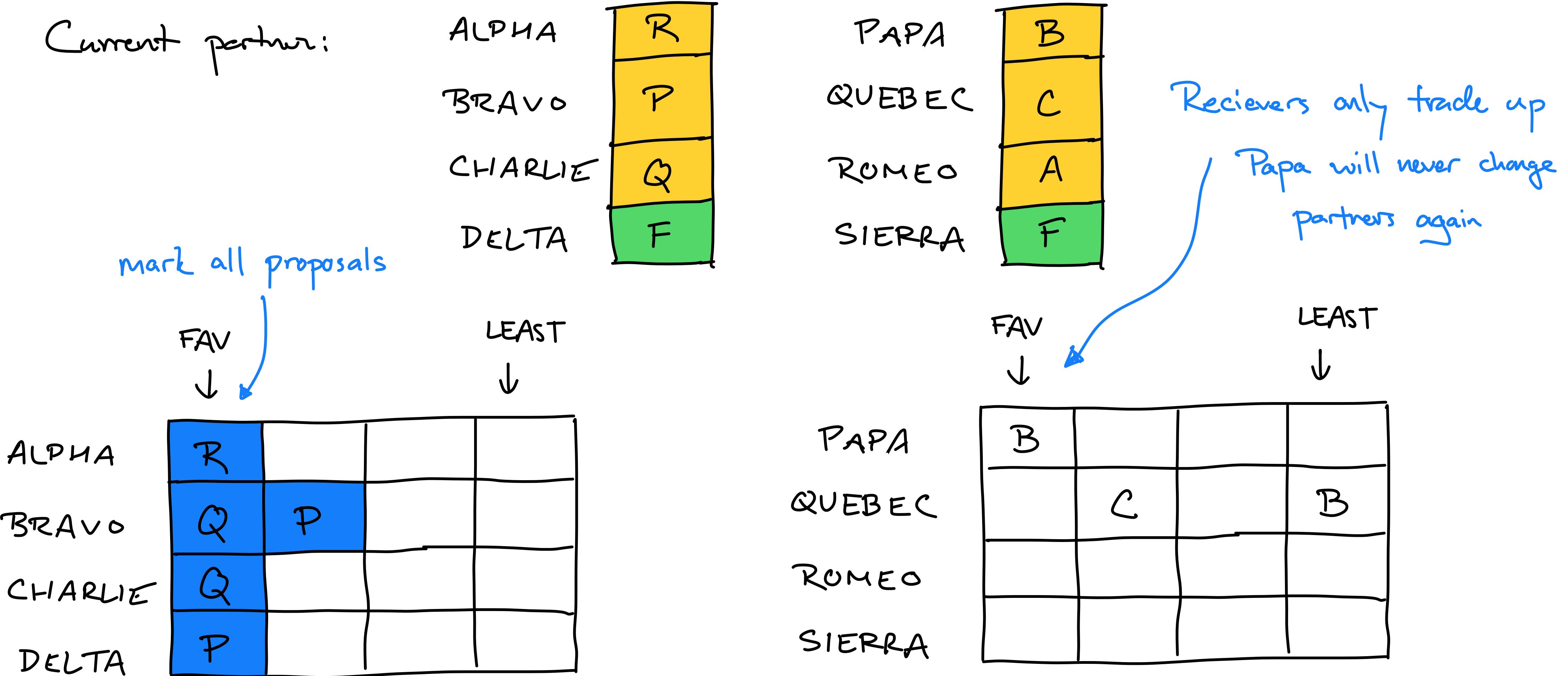
Gale-Shapley walkthrough

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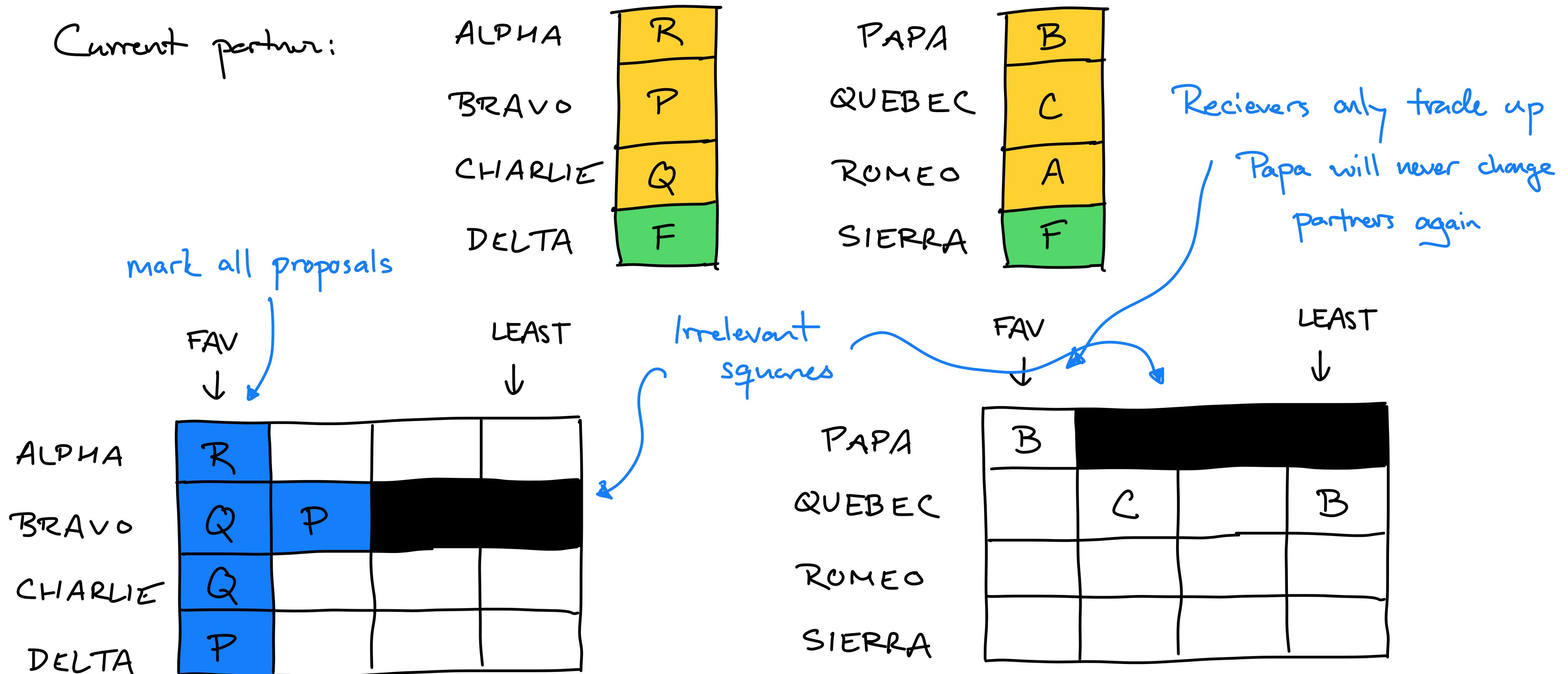


Gale-Shapley walkthrough

```

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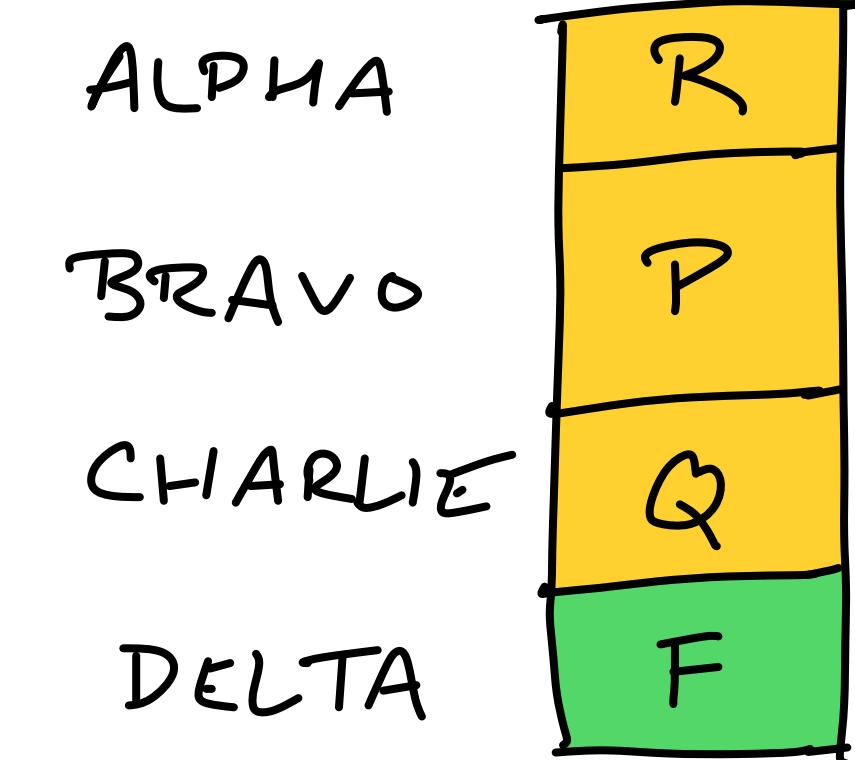
Gale-Shapley walkthrough

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Current partner:



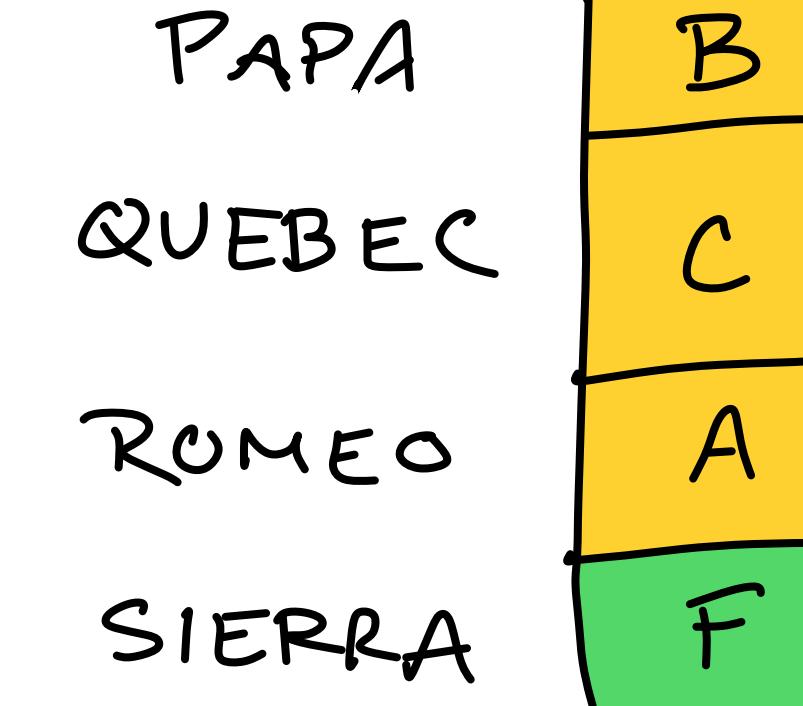
mark all proposals

FAV

LEAST

ALPHA

ALPHA	R			
BRAVO	Q	P		
CHARLIE	Q			
DELTA	P	R		



FAV

LEAST

PAPA

QUEBEC

ROMEO

SIERRA

B			
	C		B

Gale-Shapley walkthrough

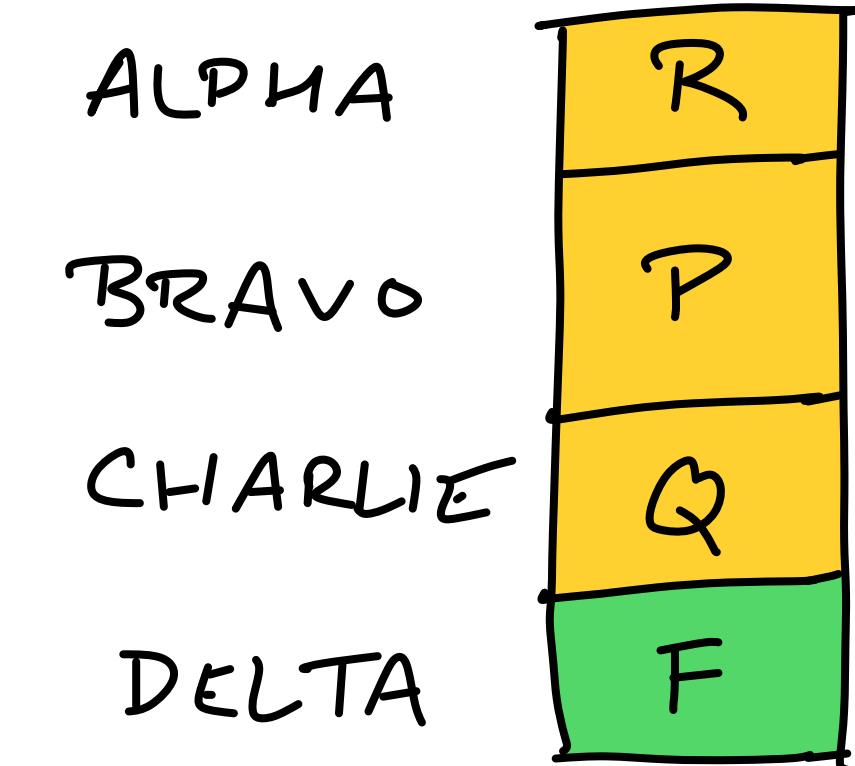
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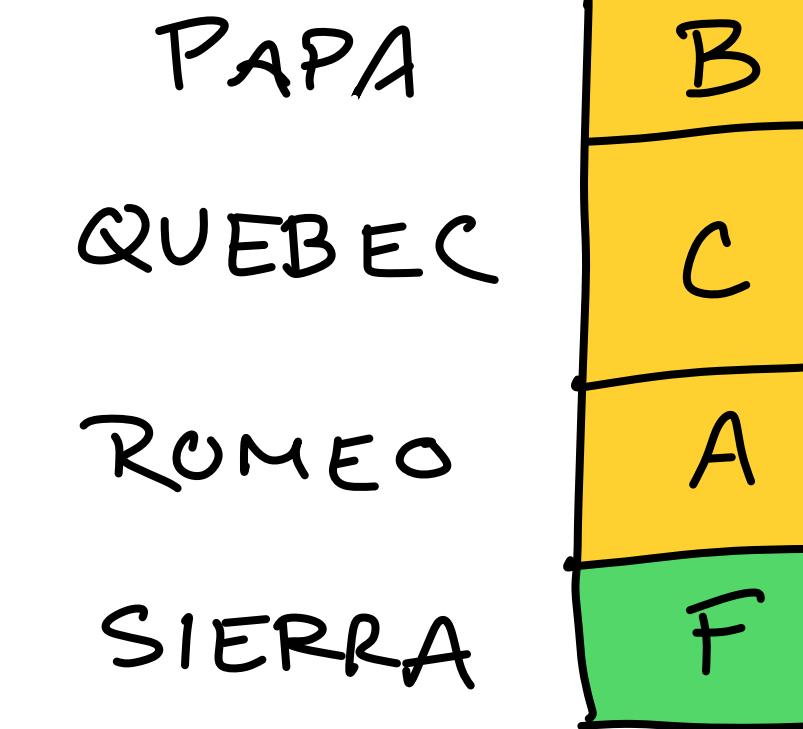
mark all proposals

FAV
 ↓

LEAST
 ↓

ALPHA
 BRAVO
 CHARLIE
 DELTA

ALPHA	R			
BRAVO	Q	P		
CHARLIE	Q			
DELTA	P	R		



FAV
 ↓

LEAST
 ↓

PAPA
 QUEBEC
 ROMEO
 SIERRA

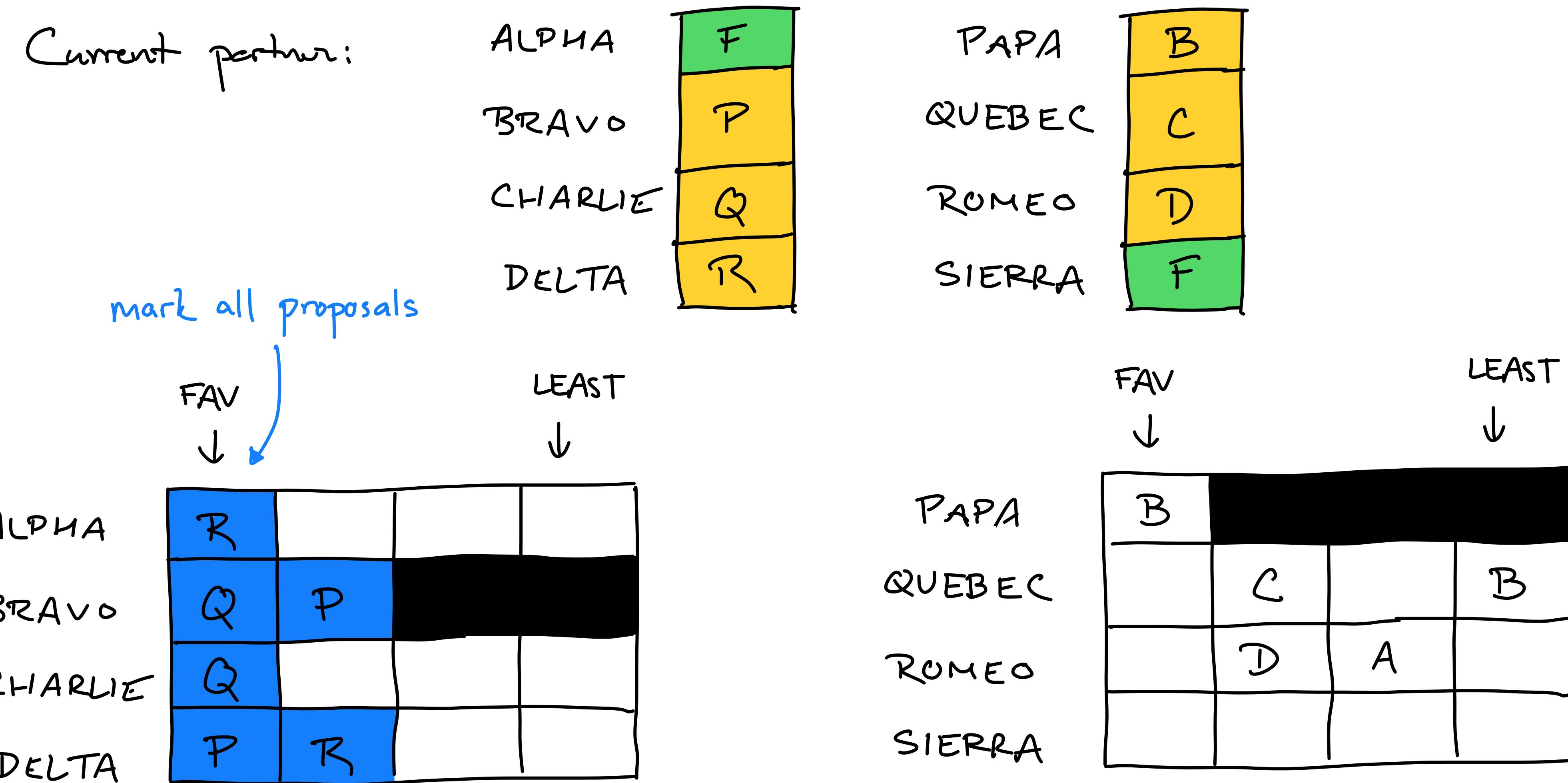
PAPA	B			
QUEBEC		C		B
ROMEO		D	A	
SIERRA				

Gale-Shapley walkthrough

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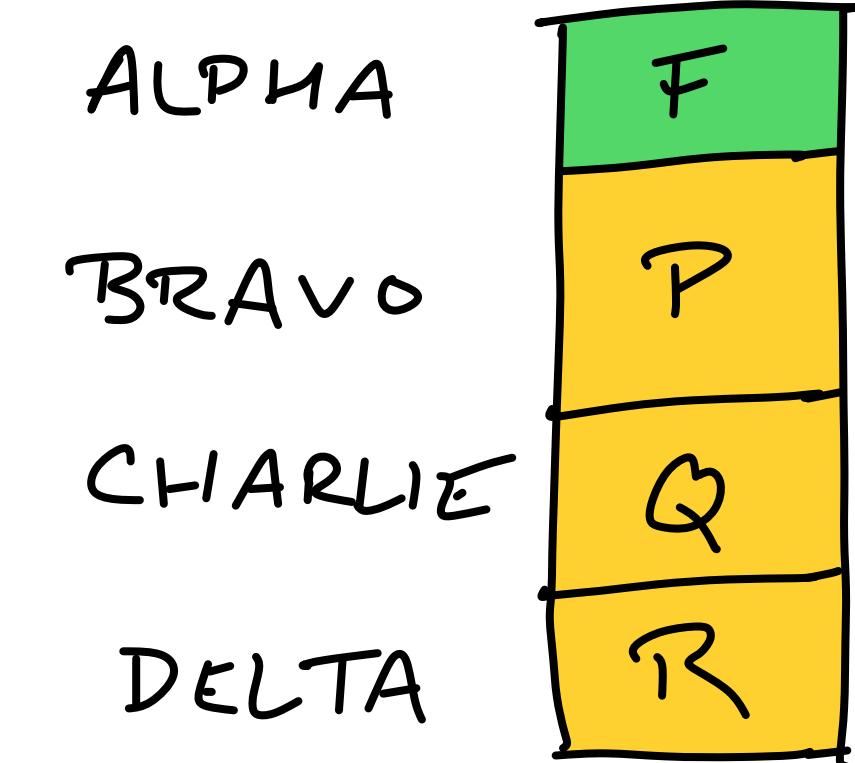
Gale-Shapley walkthrough

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Current partner:



mark all proposals

FAV

LEAST

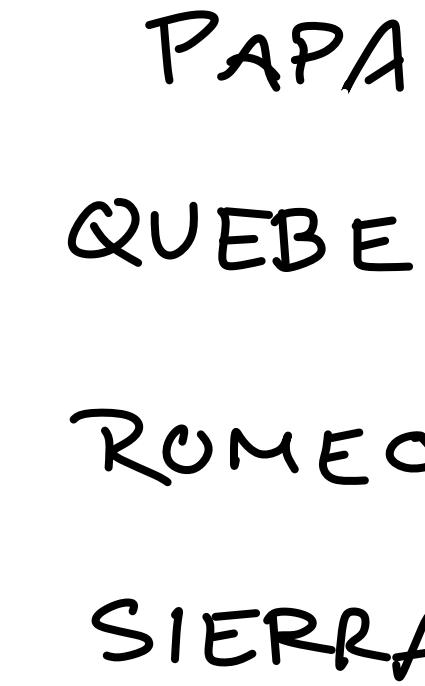
ALPHA

R	S		
Q	P		
Q			
P	R		

BRAVO

CHARLIE

DELTA



FAV

LEAST

PAPA

QUEBEC

ROMEO

SIERRA

B			
	C		B
	D	A	

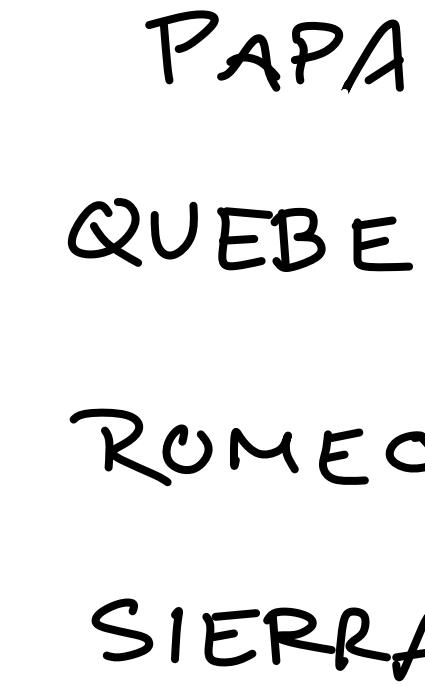
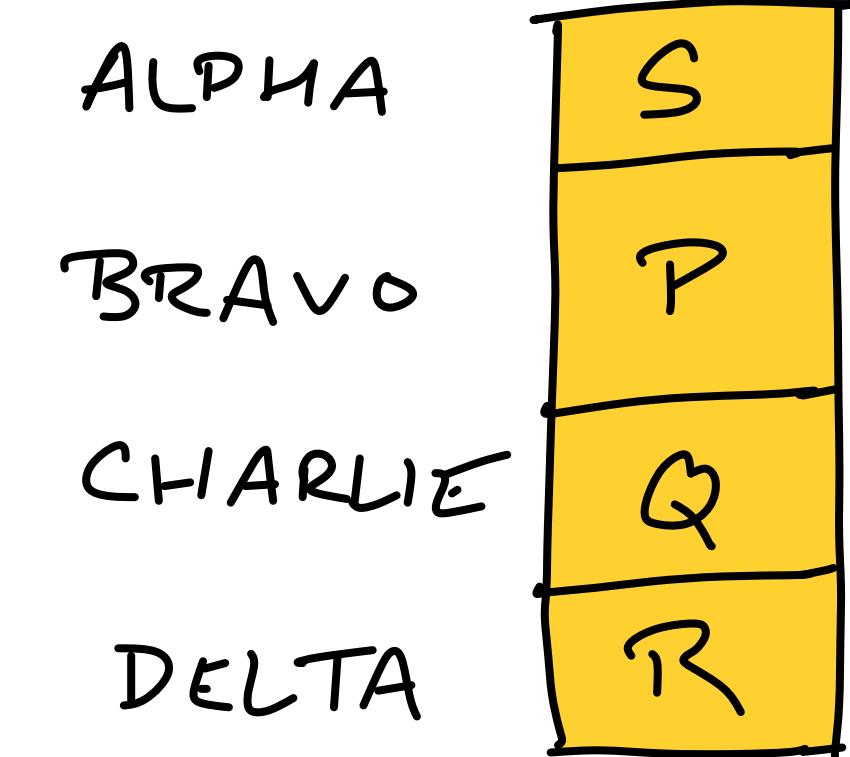
Gale-Shapley walkthrough

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```

Current partner:



mark all proposals

FAV

LEAST

ALPHA

	R	S		
	Q	P		
	Q			
	P	R		

BRAVO

PAPA

CHARLIE

QUEBEC

DELTA

ROMEO

SIERRA

FAV

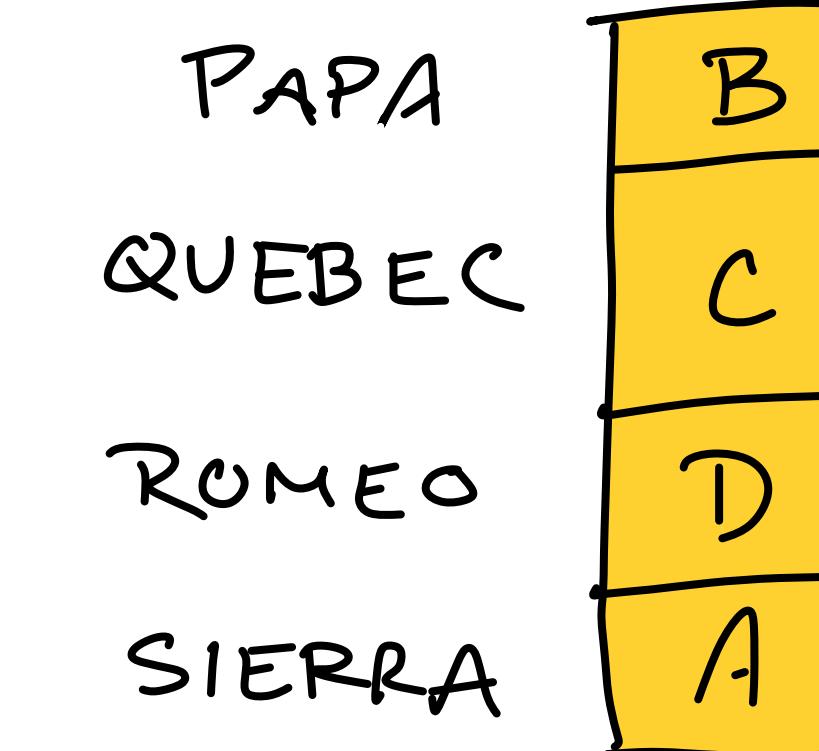
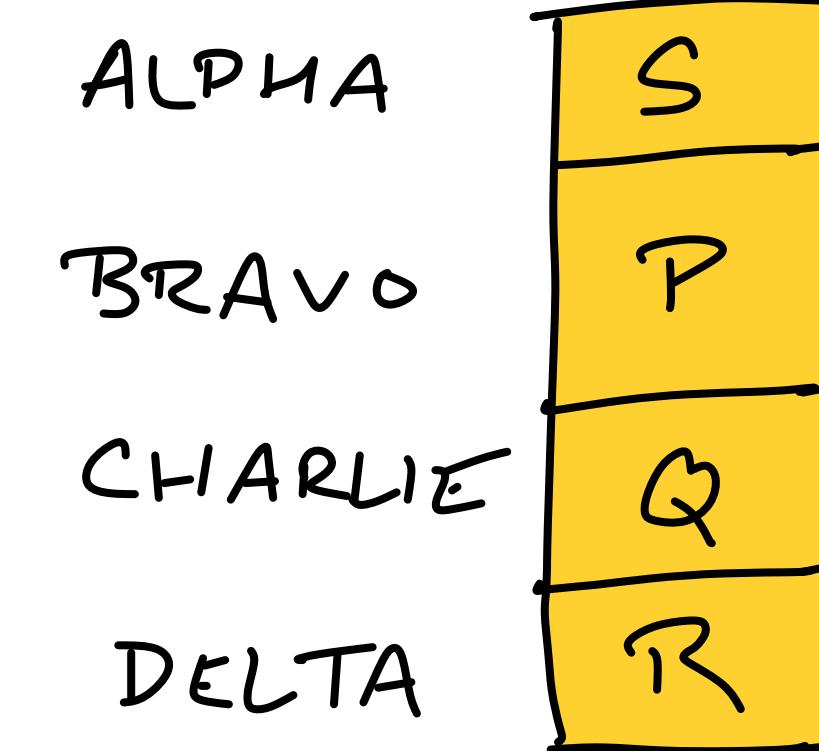
LEAST

B			
	C		B
	D	A	

Gale-Shapley walkthrough

no free proposers.
 Alg terminates and everyone
 is matched.

Current partner:



mark all proposals

FAV

LEAST

ALPHA
 BRAVO
 CHARLIE
 DELTA

	R	S		
	Q	P		
	Q			
	P	R		

FAV

LEAST

PAPA
 QUEBEC
 ROMEO
 SIERRA

B			
	C		B
	D	A	

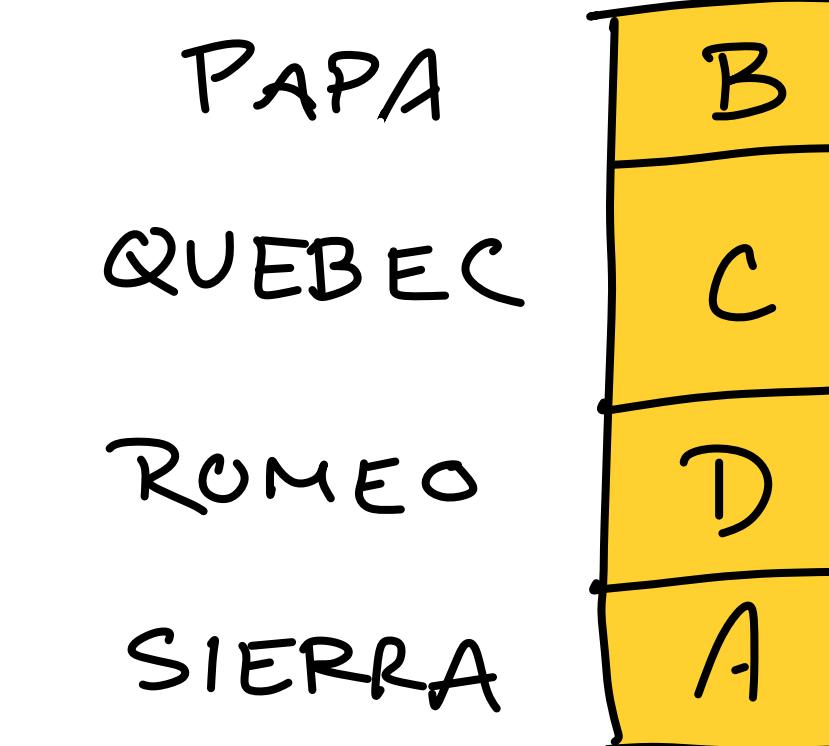
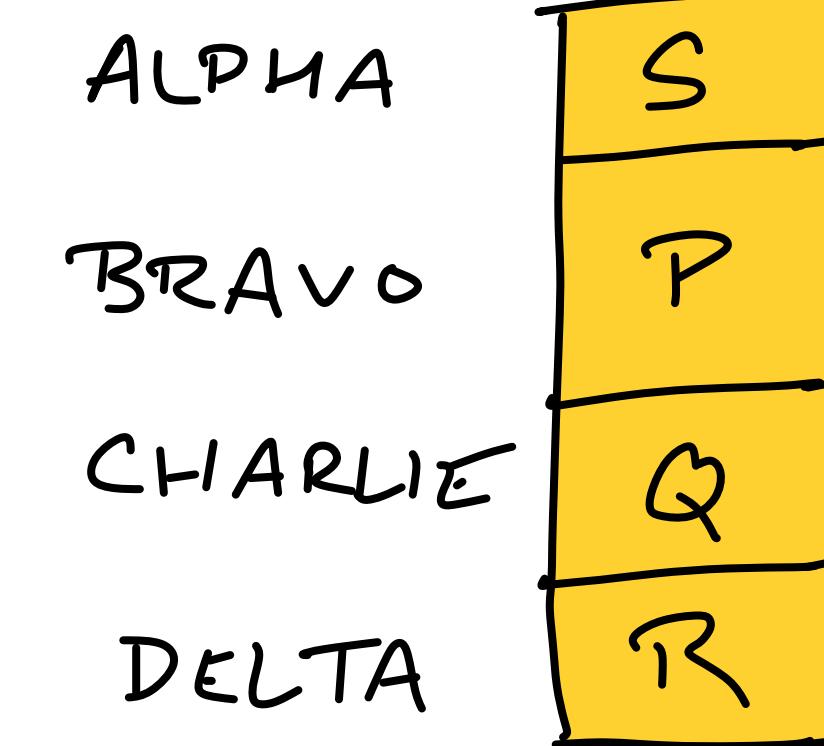
check out how
 empty the receiver
 preference matrix is.

never even
 considered

Gale-Shapley walkthrough

no free proposers.
 Alg terminates and everyone
 is matched.

Current partner:



mark all proposals

FAV



LEAST



ALPHA

	R	S		
	Q	P		
	Q			
	P	R		

BRAVO

PAPA

CHARLIE

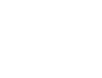
QUEBEC

DELTA

ROMEO

SIERRA

FAV



LEAST



B			
	C		B
	D	A	

Gale-Shapley walkthrough

```

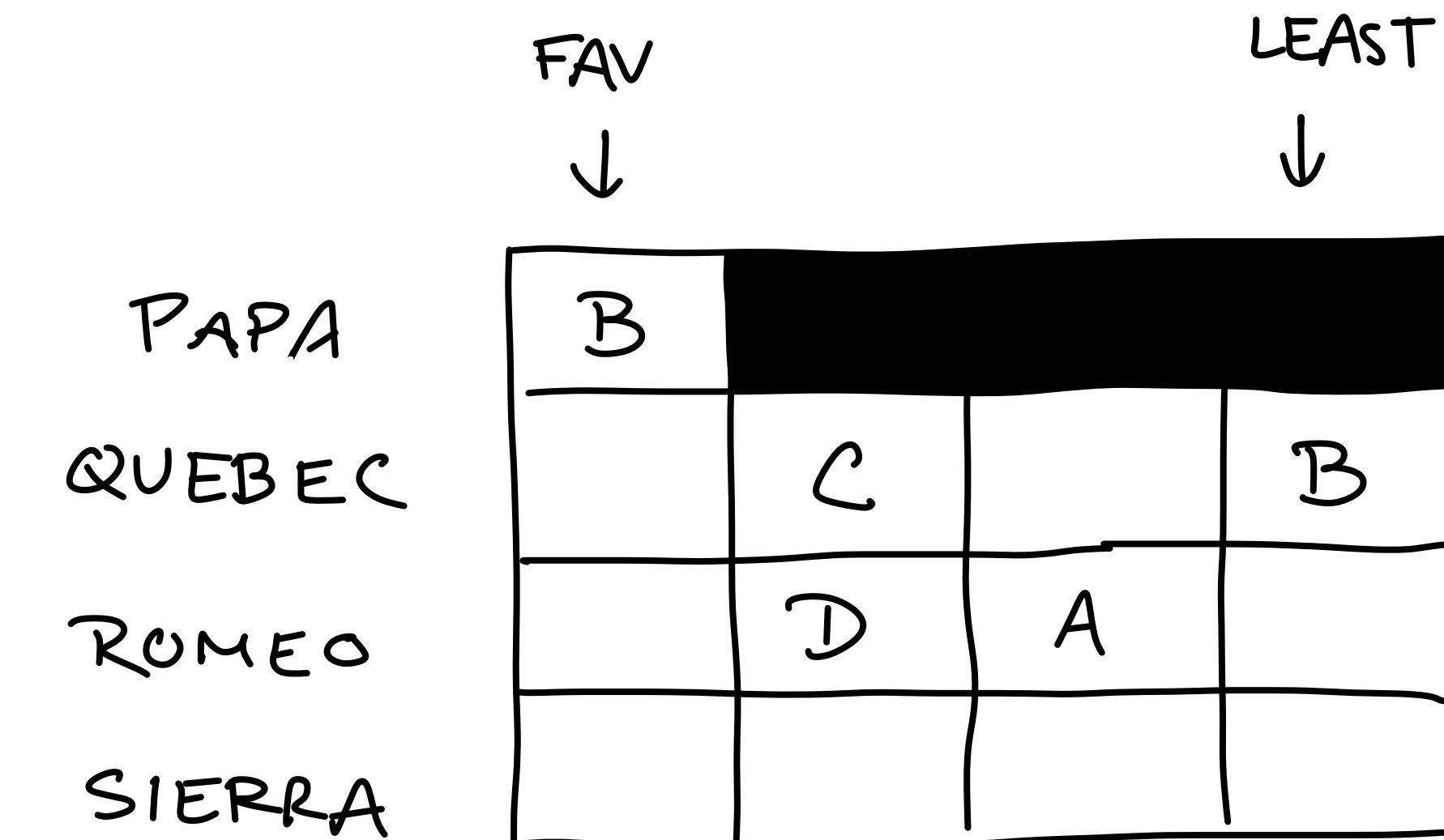
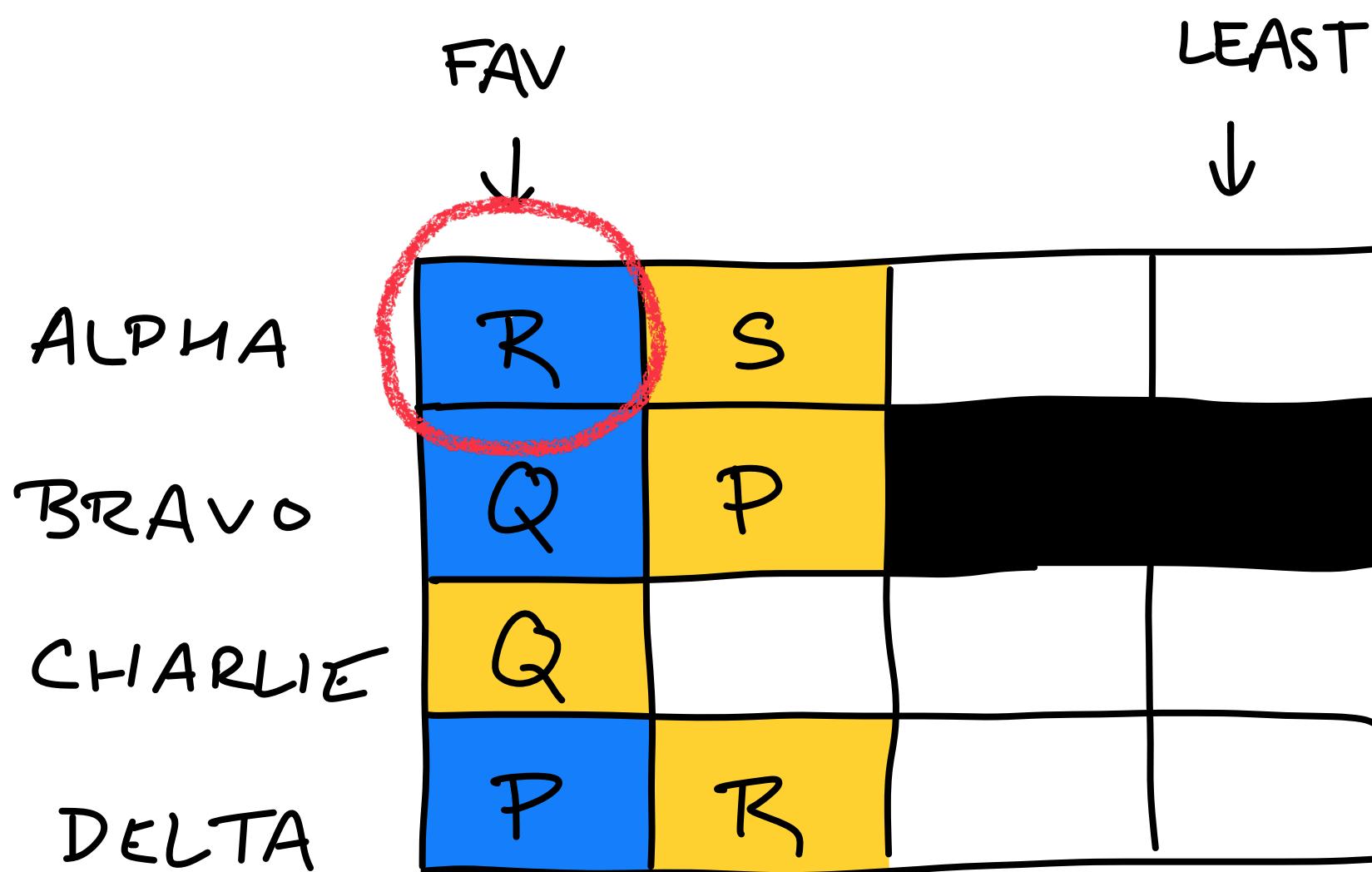
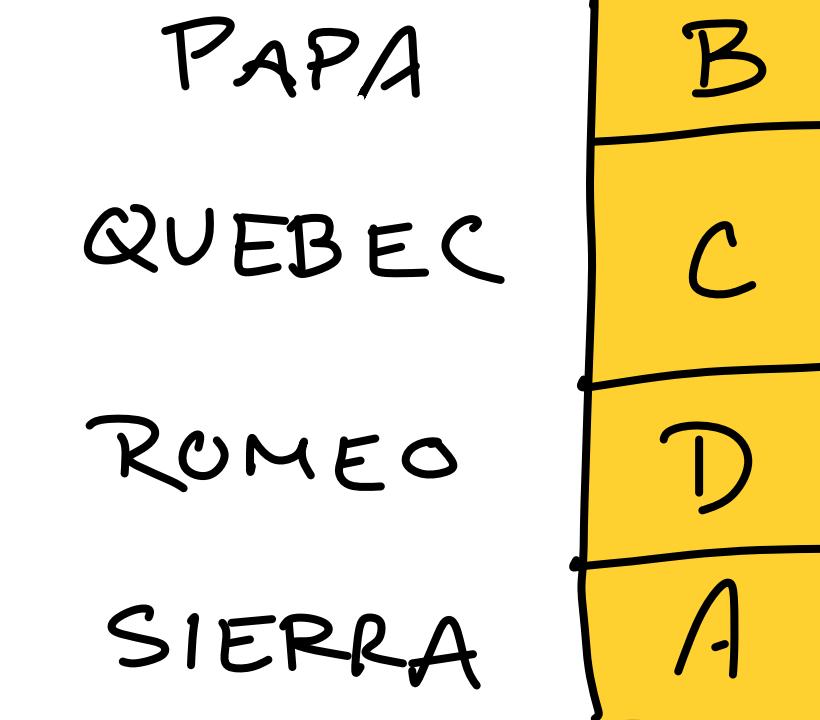
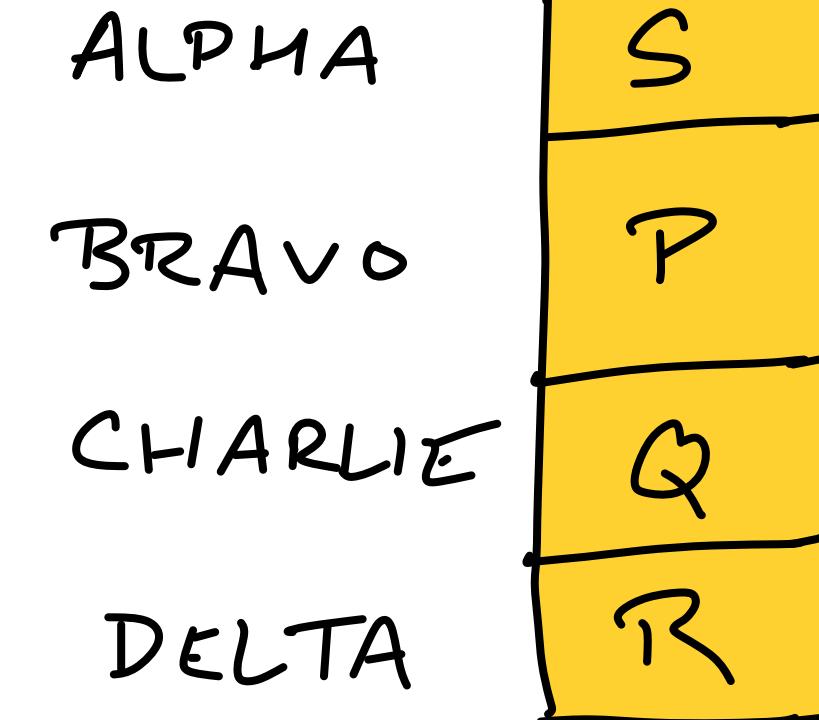
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}

```

Current partner:

Is (A, R)
stable?



Gale-Shapley walkthrough

```

Initialize each person to be free.

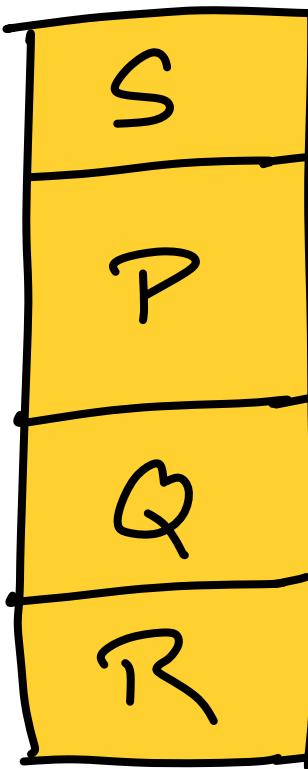
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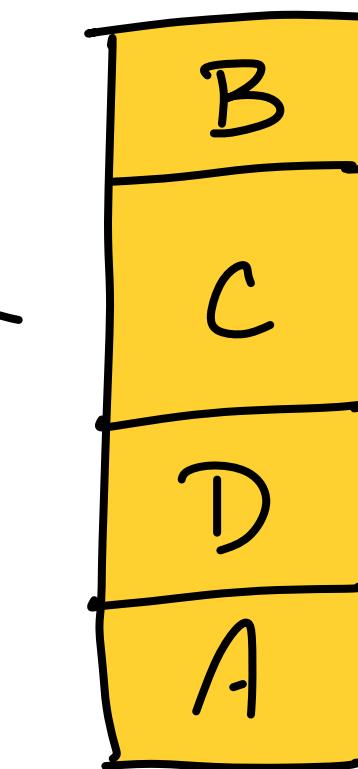
Current partner:

Is (A,Q)
stable?

ALPHA
BRAVO
CHARLIE
DELTA



PAPA
QUEBEC
ROMEO
SIERRA



	FAV	LEAST	
ALPHA	R	S	
BRAVO	Q	P	
CHARLIE	Q		
DELTA	P	R	

	FAV	LEAST	
PAPA	B		
QUEBEC	C		B
ROMEO	D	A	
SIERRA			

Gale-Shapley walkthrough

```

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```

Current partner:

Is (A, P)
stable?

