

CSE 421 Winter 2025

Lecture 21: Linear **Duality** / **Relaxation**
Literal not LP

Owen Boseley

<http://www.cs.uw.edu/421>

throw new
NoPollEVFoundException

Hello.

My name is Owen! I'm uh. An undergrad who is obsessed with teaching and theory.

{Resume = null, Teaching Credentials = null}

I'll be available for questions after lecture! (only 10 min I have crypto
☹)



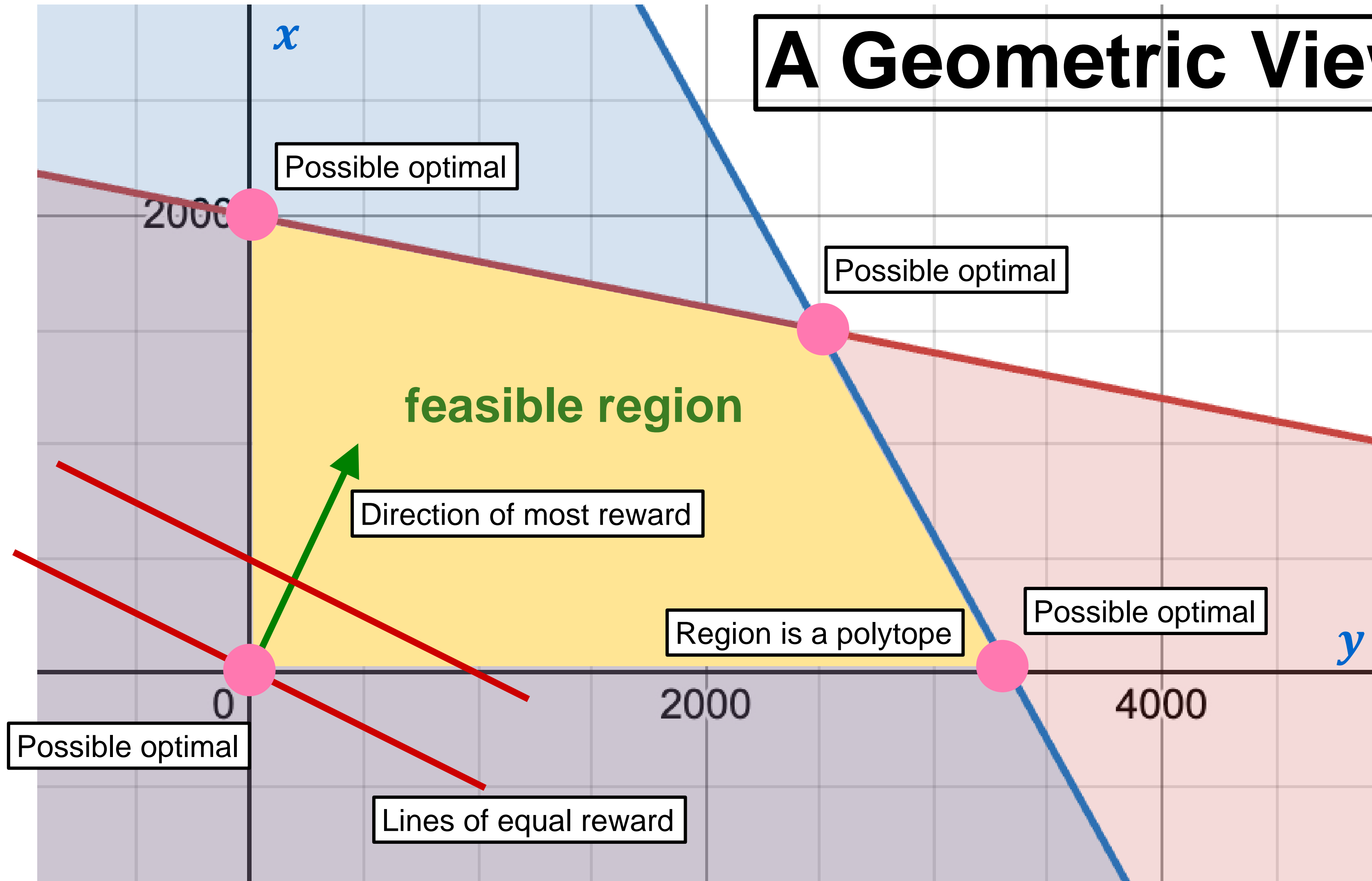
Topic of the Day.

- Linear Programs Review
- Linear Duality
- Boba Buying

A review of LPs

- We looked at a very general class of optimization problems
- We looked at few ways of **understanding** them:
 - **Geometric** View
 - **Linear Algebraic** View
- We looked at ways of **writing** them:
 - **Standard** Form
- We looked at ways of **applying** them to 421:
 - **Max Flow / Min Cut**
 - **Selling Boba.. lul**

A Geometric View



An Algebraic View

$$\begin{array}{ll} \text{maximize} & c_1x_1 + \cdots + c_nx_n \\ \text{subject to} & a_{11}x_1 + \cdots + a_{1n}x_n \leq b_1 \\ & \vdots \\ & a_{m1}x_1 + \cdots + a_{mn}x_n \leq b_m \\ & x_1, \dots, x_n \geq 0 \end{array}$$

(equivalently)

$$\begin{array}{ll} \text{maximize} & c^T x \\ \text{subject to} & Ax \leq b \\ & x \geq 0 \end{array}$$

This is **standard form**:

- maximization
- \leq inequalities with constant RHS
- nonnegative x 's

Now back to Boba...

Setting.

- Glenn and I love boba. So much that we want to buy out the entirety of Nathan's store's stock!
- Our **Goal**: Buy all the tea and boba at minimal cost
 - Cause we are TAs and broke...

Setting.

- Glenn and I love boba. So much that we want to buy out the entirety of Nathan's store's stock!
- Our **Goal**: Buy all the tea and boba at minimal cost
 - Cause we are TAs and broke...
- Our **Strategy**: Offer the shop a rate for each ingredient
 - We offer **x cents / oz of boba**, **y cents / oz of tea**

Setting.

- Glenn and I love boba. So much that we want to buy out the entirety of Nathan's store's stock!
- Our **Goal**: Buy all the tea and boba at minimal cost
 - Cause we are TAs and broke...
- Our **Strategy**: Offer the shop a rate for each ingredient
 - We offer **x cents / oz of boba**, **y cents / oz of tea**
- Our **Constraints**: Nathan wants us to buy through his standard and premium drinks
 - Our rates must behave as if we had bought standard drinks

Claim, solve with an LP.

$$\begin{array}{ll} \text{maximize} & 20s + 40p \\ \text{subject to} & 0.1s + 0.5p \leq 1000 \\ & 0.9s + 0.5p \leq 3000 \\ & s, p \geq 0 \end{array}$$

Original Problem

$$\begin{array}{ll} \text{minimize} & 1000x + 3000y \\ \text{subject to} & 0.1x + 0.9y \geq 20 \\ & 0.5x + 0.5y \geq 40 \\ & x, y \geq 0 \end{array}$$

Our Problem

Try this on your own at home!

Claim, solve with an LP.

$$\begin{array}{ll} \text{minimize} & 1000x + 3000y \\ \text{subject to} & 0.1x + 0.9y \geq 20 \\ & 0.5x + 0.5y \geq 40 \\ & x, y \geq 0 \end{array}$$

$$\begin{array}{ll} \text{maximize} & 20s + 40p \\ \text{subject to} & 0.1s + 0.5p \leq 1000 \\ & 0.9s + 0.5p \leq 3000 \\ & s, p \geq 0 \end{array}$$

Hold on... this feels familiar?

Question: How is this different and similar to the boba problem on Wednesday?

2 min to discuss with those around you...

Claim, solve with an LP.

$$\begin{array}{ll} \text{minimize} & 1000x + 3000y \\ \text{subject to} & 0.1x + 0.9y \geq 20 \\ & 0.5x + 0.5y \geq 40 \\ & x, y \geq 0 \end{array}$$

$$\begin{array}{ll} \text{maximize} & 20s + 40p \\ \text{subject to} & 0.1s + 0.5p \leq 1000 \\ & 0.9s + 0.5p \leq 3000 \\ & s, p \geq 0 \end{array}$$

Claim: This is the “**dual**” problem to our boba problem from Wednesday.

- We are **buying** not selling
- We are **minimizing** not maximizing

But... are the solutions the same?

Short answer, **YES!!!!!!**

Duality Definition

$$\begin{aligned} &\text{minimize } b^T y \\ &\text{subject to } A^T y \geq c \\ &\quad y \geq 0 \end{aligned}$$

The Dual
“Buying Boba”

The Primal
“Selling Boba”

$$\begin{aligned} &\text{maximize } c^T x \\ &\text{subject to } Ax \leq b \\ &\quad x \geq 0 \end{aligned}$$

See slides for other formulation....

Properties of Duals

Properties of Duals

- Property 1: The **Dual** of a **Dual** is the **Primal**.
 - Look at the definition to convince yourself

Properties of Duals

- Property 1: The **Dual** of a **Dual** is the **Primal**.
 - Look at the definition to convince yourself
- Property 2: **Weak Duality**
 - Any feasible **Dual** solution is an upper bound for any feasible **Primal** solution.
 - In some sense, the dual is trying to minimize an upper bound on the primal. Think about the boba problem!

Properties of Duals

- Property 1: The **Dual** of a **Dual** is the **Primal**.
 - Look at the definition to convince yourself
- Property 2: **Weak Duality**
 - Any feasible **Dual** solution is an upper bound for any feasible **Primal** solution.
 - In some sense, the dual is trying to minimize an upper bound on the primal. Think about the boba problem!

Take 2 min to talk about why this might be true of the boba problem!

Properties of Duals

- Property 1: The **Dual** of a **Dual** is the **Primal**.
 - Look at the definition to convince yourself
- Property 2: **Weak Duality**
 - Any feasible **Dual** solution is an upper bound for any feasible **Primal** solution.
 - In some sense, the dual is trying to minimize an upper bound on the primal. Think about the boba problem!
- Property 3: **Strong Duality**
 - The **Optimal Dual = Optimal Primal** (if the primal and dual have a solution).
 - It turns out that they are actually equal (assuming a solution exists).

But... what is a dual? Why care?

- Duals are useful **compliments** to LPs!
- They convey potentially new information and can also be used in creative proofs of seemingly different problems

Something cool?



Tool in our toolkit

But... what is a dual? Why care?

- **Intuitions** for the dual are sometimes not good... but to help motivate you guys to love duality, let's look at some approaches to understand them!

Intuition #1: Games

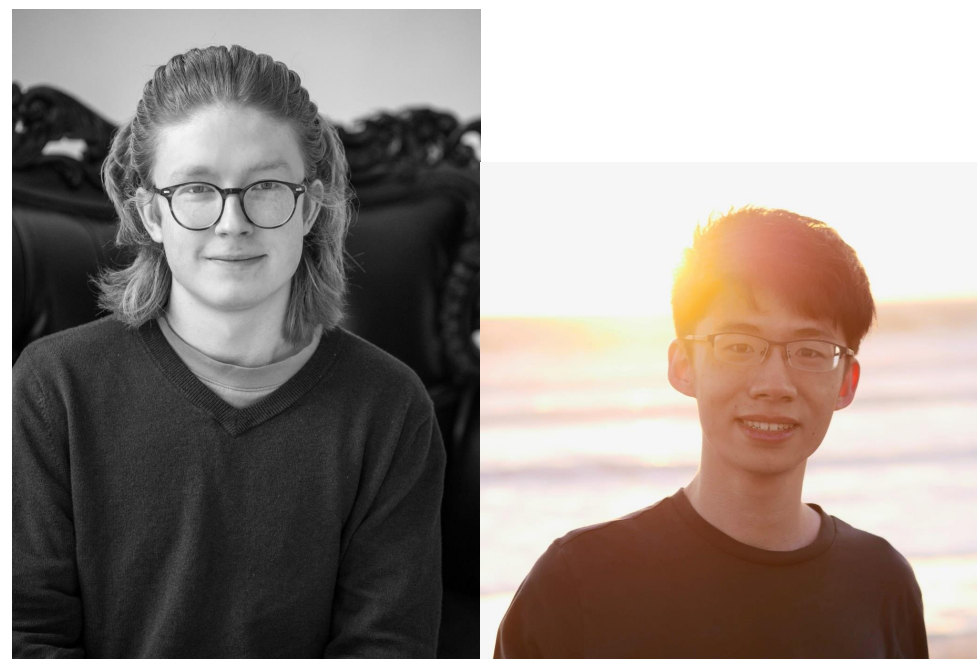
- **Primal:** We want to maximize our reward!



Sell Boba for lots!

VS.

- **Dual:** We want to minimize their reward!



Buy Boba for Cheap!

I'm sure Nathan would sell boba for reasonable prices...

Intuition #2: Geometric

Dual is pushing down!



Primal is pushing up!



There are more concrete geometric intuitions...

More fun things...

Max Flow (Revisited)

Input: A flow network $G = (V, E)$, source s , sink t , and $c : E \rightarrow \mathbb{R}^{\geq 0}$

Goal:

maximize flow out of s

subject to respecting capacities and flow conservation

We want $x_e =$ flow on edge $e \in E$.

maximize $\sum_{e \text{ out of } s} (x_e)$

subject to $0 \leq x_e \leq c(e)$ for all $e \in E$

$\sum_{e \text{ out of } v} (x_e) = \sum_{e \text{ into } v} (x_e)$ for all $v \in V \setminus \{s, t\}$

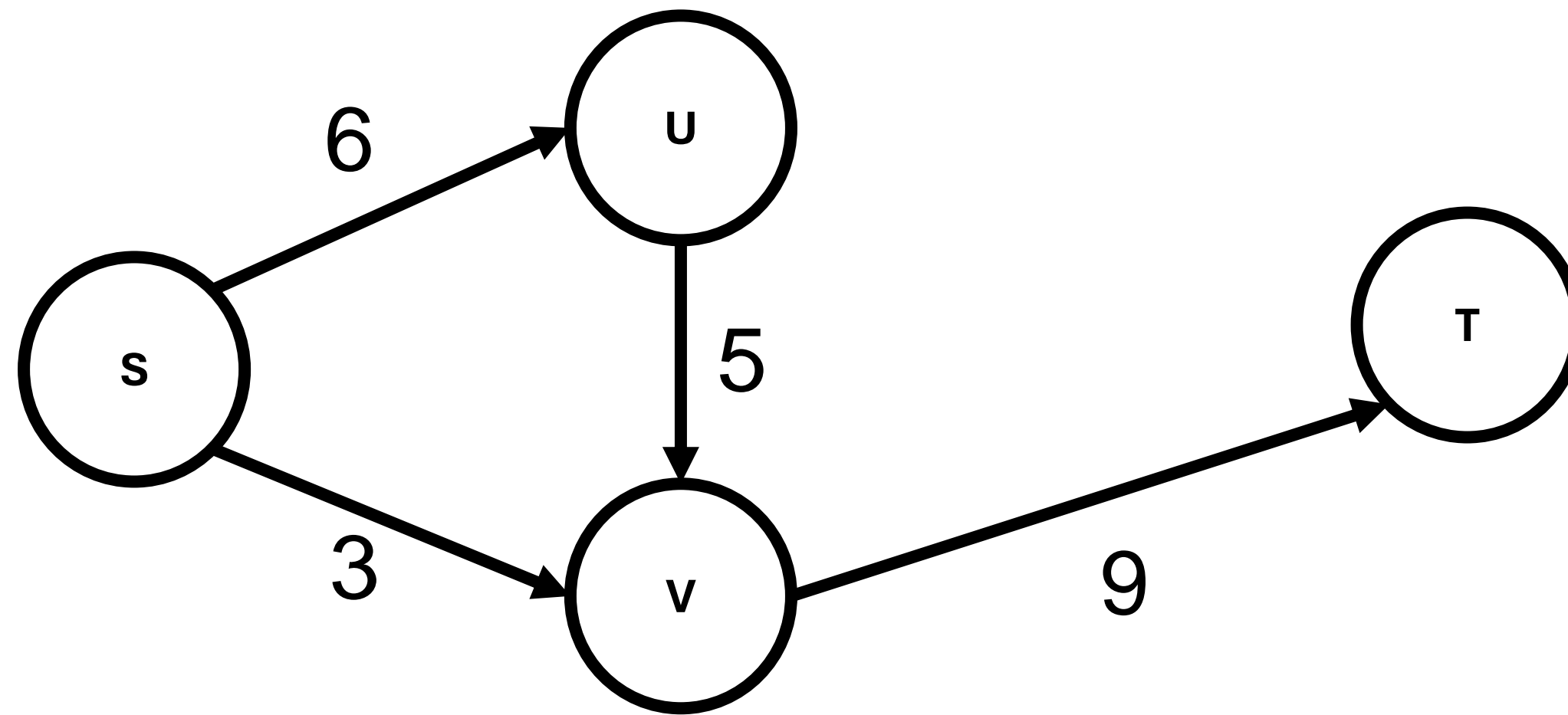
Max Flow (Revisited)

$$\begin{aligned} & \text{maximize} && \sum_{e \text{ out of } s} (x_e) \\ & \text{subject to} && \mathbf{0} \leq x_e \leq c(e) \text{ for all } e \in E \\ & && \sum_{e \text{ out of } v} (x_e) = \sum_{e \text{ into } v} (x_e) \text{ for all } v \in V \setminus \{s, t\} \end{aligned}$$

In standard form:

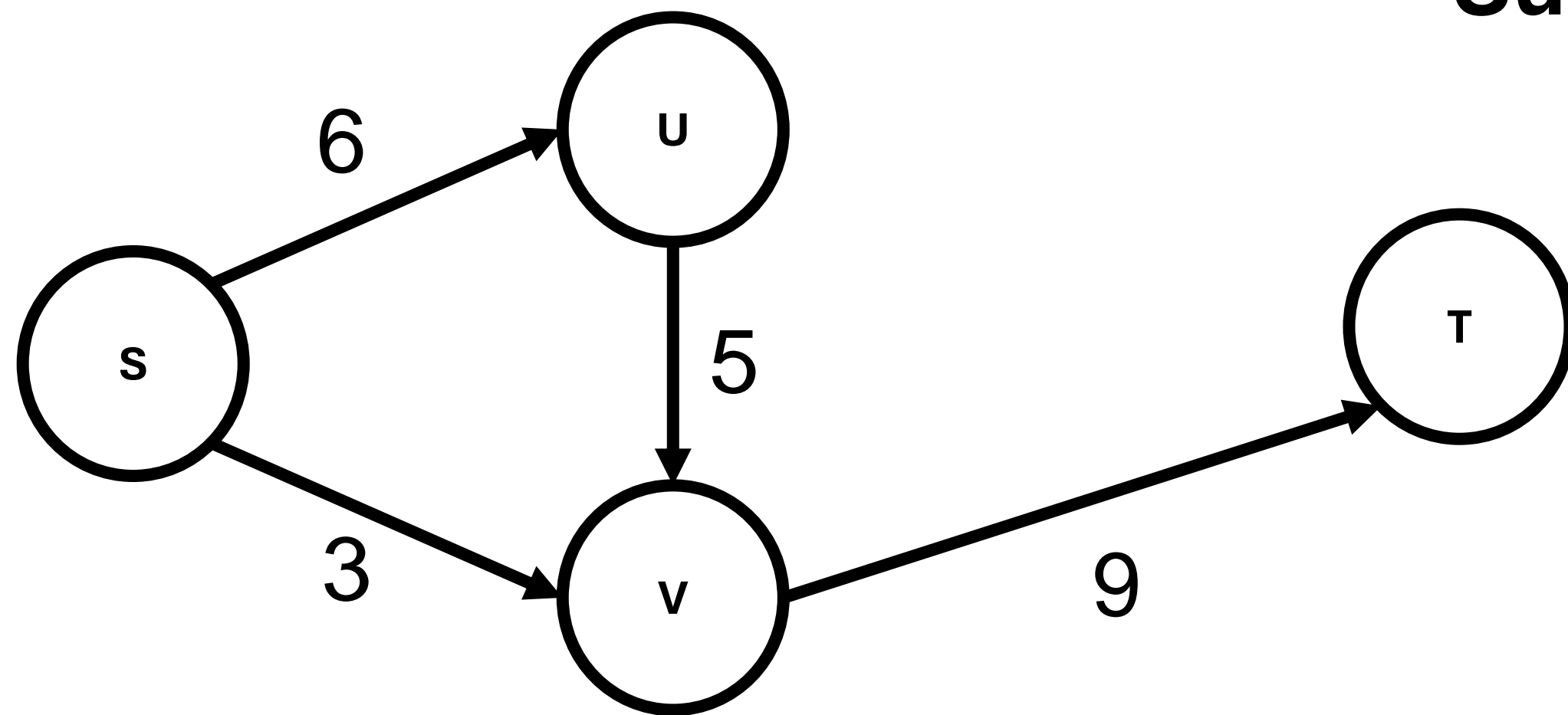
$$\begin{aligned} & \text{maximize} && \sum_{e \text{ out of } s} (x_e) \\ & \text{subject to} && x_e \leq c(e) \text{ for all } e \in E \\ & && \sum_{e \text{ out of } v} (x_e) - \sum_{e \text{ into } v} (x_e) \leq \mathbf{0} \text{ for all } v \in V \setminus \{s, t\} \\ & && \sum_{e \text{ into } v} (x_e) - \sum_{e \text{ out of } v} (x_e) \leq \mathbf{0} \text{ for all } v \in V \setminus \{s, t\} \\ & && x_e \geq \mathbf{0} \text{ for all } e \in E \end{aligned}$$

Max Flow Dual



Take 2-3 min to make the LP with those around you

Max Flow Dual



Maximize $x_1 + x_2$

Subject to $x_1 \leq 6, x_2 \leq 3, x_3 \leq 5, x_4 \leq 9$

$$x_1 - x_3 \leq 0$$

$$x_3 - x_1 \leq 0$$

$$x_4 - x_3 - x_2 \leq 0$$

$$x_3 + x_2 - x_4 \leq 0$$

$$\sum_{e \text{ out of } s} (x_e)$$

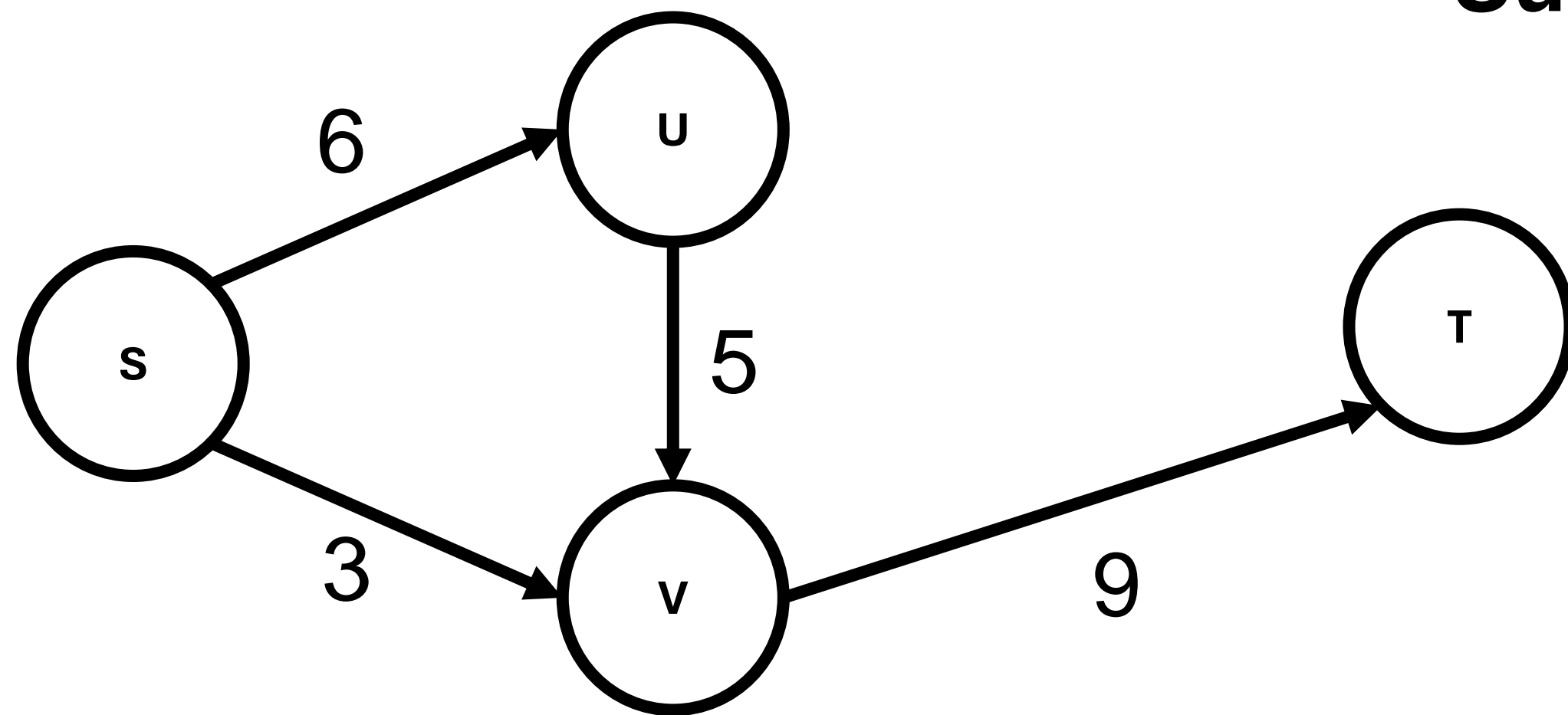
$$x_e \leq c(e) \text{ for all } e \in E$$

$$\sum_{e \text{ out of } v} (x_e) - \sum_{e \text{ into } v} (x_e) \leq 0 \text{ for all } v \in V \setminus \{s, t\}$$

$$\sum_{e \text{ into } v} (x_e) - \sum_{e \text{ out of } v} (x_e) \leq 0 \text{ for all } v \in V \setminus \{s, t\}$$

$$x_e \geq 0 \text{ for all } e \in E$$

Max Flow Dual



Minimize $6y_1 + 3y_2 + 5y_3 + 9y_4$

Subject to $y_1 + y_5 - y_6 \geq 1$

$y_2 + y_8 - y_7 \geq 1$

$y_3 + y_6 + y_8 - y_5 - y_7 \geq 0$

$y_4 + y_7 - y_8 \geq 0$

$\sum_{e \text{ out of } s} (x_e)$

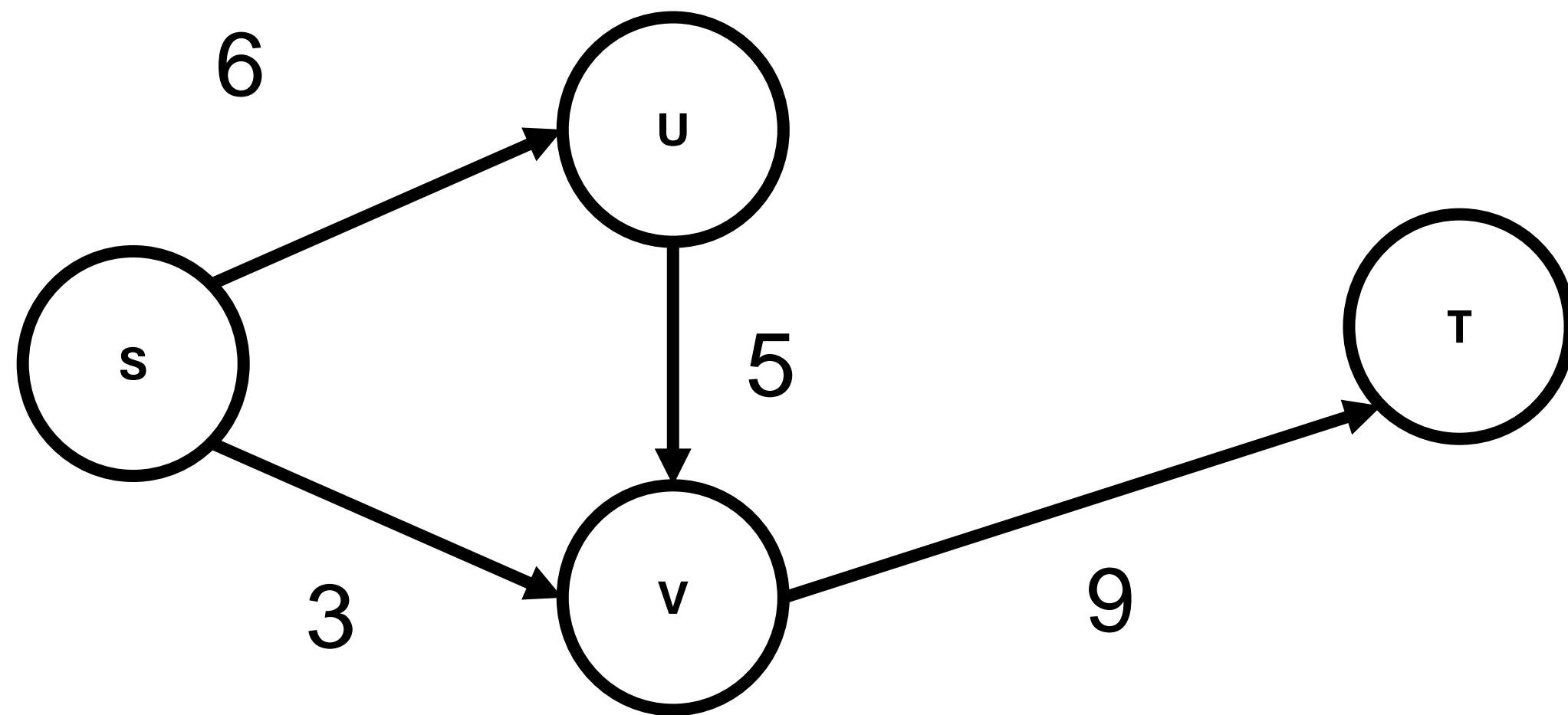
$x_e \leq c(e)$ for all $e \in E$

$\sum_{e \text{ out of } v} (x_e) - \sum_{e \text{ into } v} (x_e) \leq 0$ for all $v \in V \setminus \{s, t\}$

$\sum_{e \text{ into } v} (x_e) - \sum_{e \text{ out of } v} (x_e) \leq 0$ for all $v \in V \setminus \{s, t\}$

$x_e \geq 0$ for all $e \in E$

Max Flow Dual



Optimal Solution:

$$y_1 = 0 \quad y_2 = 1 \quad y_3 = 1 \quad y_4 = 0$$
$$y_5 = 1 \quad y_6 = 0 \quad y_7 = 0 \quad y_8 = 0$$

$$\sum_{e \text{ out of } s} (x_e)$$

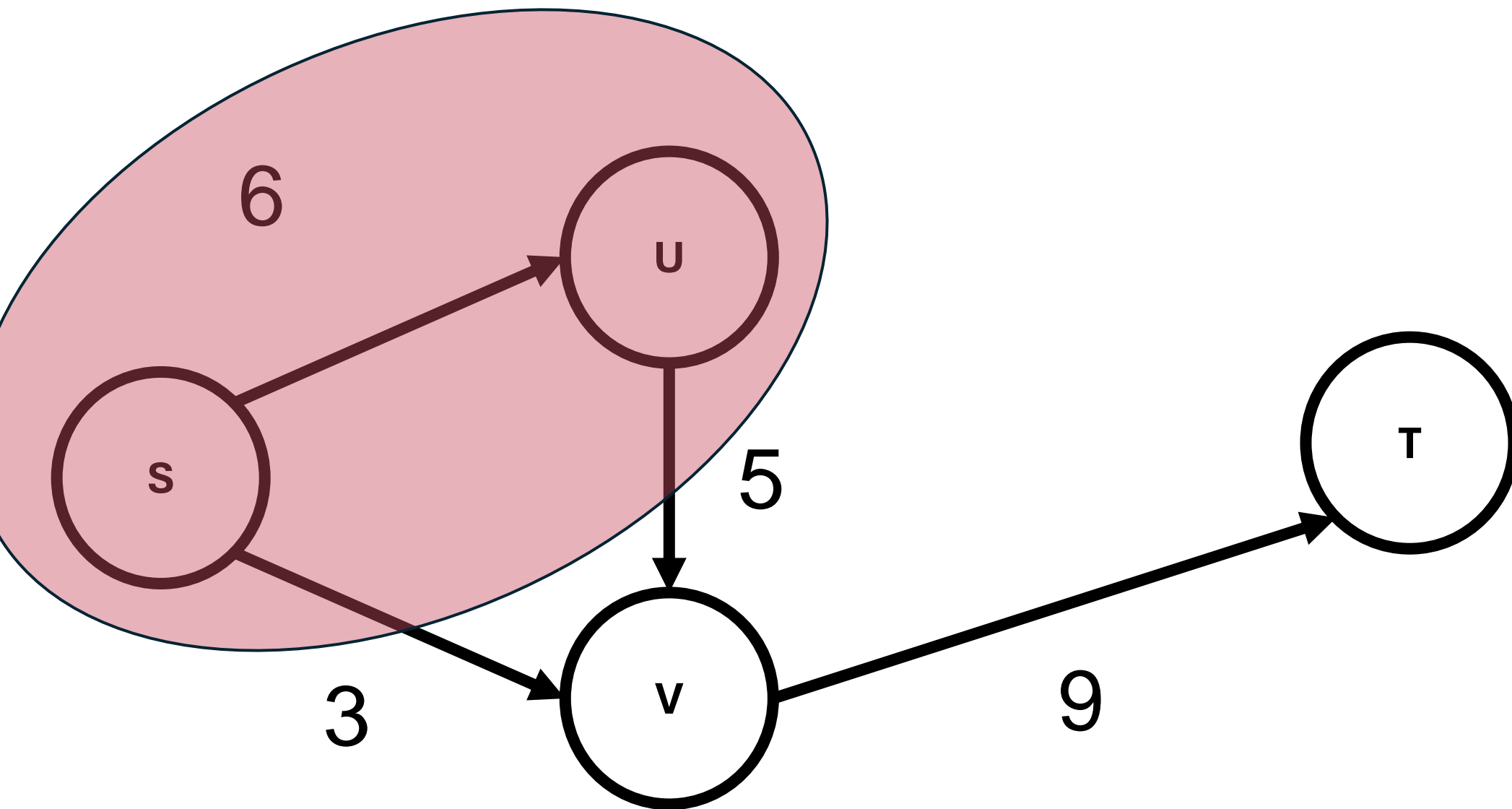
$$x_e \leq c(e) \text{ for all } e \in E$$

$$\sum_{e \text{ out of } v} (x_e) - \sum_{e \text{ into } v} (x_e) \leq 0 \text{ for all } v \in V \setminus \{s, t\}$$

$$\sum_{e \text{ into } v} (x_e) - \sum_{e \text{ out of } v} (x_e) \leq 0 \text{ for all } v \in V \setminus \{s, t\}$$

$$x_e \geq 0 \text{ for all } e \in E$$

Max Flow Dual



This is just the minimum cut!

Optimal Solution:

$$y_1 = 0 \quad y_2 = 1 \quad y_3 = 1 \quad y_4 = 0$$

$$y_5 = 1 \quad y_6 = 0 \quad y_7 = 0 \quad y_8 = 0$$

$$\sum_{e \text{ out of } s} (x_e)$$

$$x_e \leq c(e) \text{ for all } e \in E$$

$$\sum_{e \text{ out of } v} (x_e) - \sum_{e \text{ into } v} (x_e) \leq 0 \text{ for all } v \in V \setminus \{s, t\}$$

$$\sum_{e \text{ into } v} (x_e) - \sum_{e \text{ out of } v} (x_e) \leq 0 \text{ for all } v \in V \setminus \{s, t\}$$

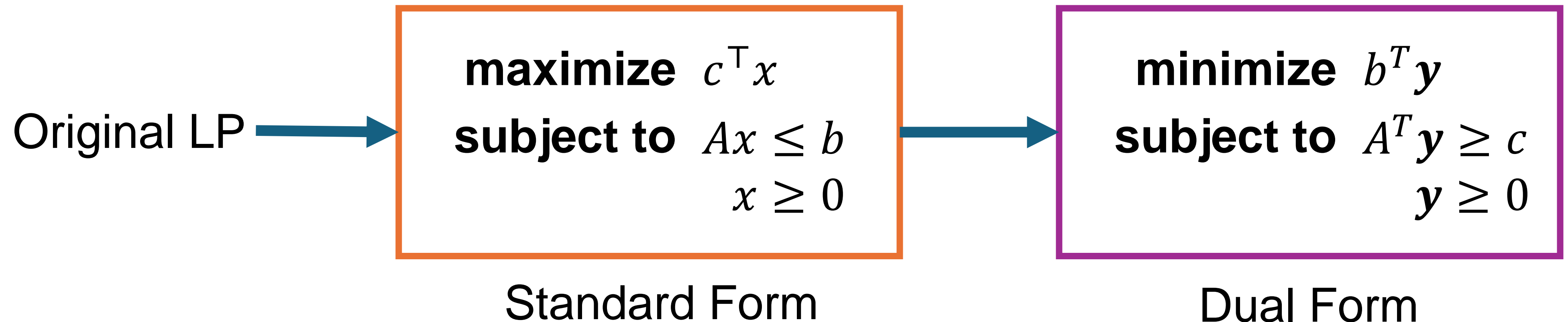
$$x_e \geq 0 \text{ for all } e \in E$$

Claim:

The LP Dual of Maximum Flow is the Minimum Cut!

(exam) Tips to compute duality

- Convert the primal into standard form first!
- Use the formula!
 - I recommend the linear algebra way because it is easier to remember (just transpose, swap, swap)



(exam) **Tips to compute duality**

- Convert the primal into standard form first!
- Use the formula!
 - I recommend the linear algebra way because it is easier to remember (just transpose, swap, swap)
- Remember that the variables of your dual are fundamentally **not** the same as your primal.

Boba Problem: Amount of drinks vs the rate of ingredients

(exam) **Tips to compute duality**

- Convert the primal into standard form first!
- Use the formula!
 - I recommend the linear algebra way because it is easier to remember (just transpose, swap, swap)
- Remember that the variables of your dual are fundamentally **not** the same as your primal.
- Remember the dual of an LP is its own LP, meaning it can be converted to standard form as well!

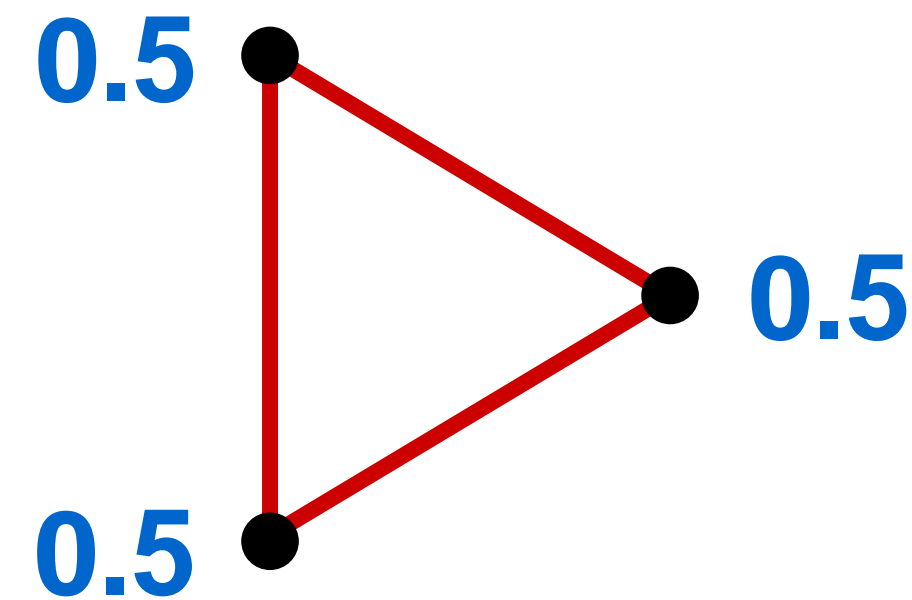
Sneak peek: Vertex Cover?

LP Relaxation: Instead of $x_v = 0$ or 1 (out/in), have $0 \leq x_v \leq 1$

$$\begin{aligned} &\text{minimize } \sum_{v \in V} (x_v) \\ &\text{subject to } x_u + x_v \geq 1 \text{ for all edges } (u, v) \in E \\ &\quad \quad \quad 0 \leq x_v \leq 1 \text{ for all vertices } v \in V \end{aligned}$$

Might give “**fractional solutions**”:

LP optimum = 1.5, true optimum = 2





Back to your regularly scheduled program....

Have a good weekend.

