

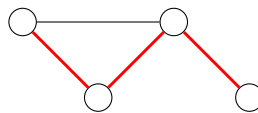
P1) Write the following LP in Standard Form:

$$\begin{aligned} \min \quad & 2x_1 - x_2 \\ \text{s.t.}, \quad & x_1 - x_3 = 4 \\ & 2x_2 - x_3 \geq 5 \\ & x_1, x_3 \geq 0. \end{aligned}$$

P2) Write the dual of the following program:

$$\begin{aligned} \max \quad & 2x_1 + x_2 \\ \text{s.t.}, \quad & x_1 + x_2 \leq 5 \\ & x_1 - 2x_2 \leq 2 \\ & x_1, x_2 \geq 0 \end{aligned}$$

- P3) Write the dual of the LP relaxation of the min vertex cover then turn it into the standard form.
- P4) A Hamiltonian cycle in a directed graph with  $n$  vertices is a directed cycle of length  $n$ , i.e., it is a cycle that visits all vertices exactly once and returns back to the starting point. A directed Hamiltonian path in a graph with  $n$  vertices is a path of length  $n - 1$ , i.e., it is a path that visits all vertices of the graph exactly once. For example, the following graph has a Hamiltonian path marked in red but no Hamiltonian cycle.



Hamiltonian-cycle problem is defined as follows: Given a graph  $G = (V, E)$ , does it have a Hamiltonian cycle?

Hamiltonian-path problem is defined as follows: Given a graph  $G = (V, E)$ , does it have a Hamiltonian path?

Prove that  $\text{Hamil-path} \leq_P \text{Hamil-cycle}$ . In other words, suppose we have a program  $A$  that solves the Hamiltonian cycle problem. Design a polynomial time algorithm that on an input graph  $G$  uses  $A$  only polynomial number of times and in polynomial time returns the solution to the Hamiltonian path problem on the given graph  $G$ .