

# CSE 421 Section 10

**Final Review**

# Announcements & Reminders

- **HW8**
  - Was due yesterday, Wednesday 12/4
- **Final review** with Professor Beame: **Sunday, 12/8 @ TBA on Zoom**
  - He will go over the practice final, so try it before the session if you can
- The **final exam** is on **Monday, 12/9 @ 2:30-4:45 @ CSE2 G20**
  - If you are sick the day of the exam, let us know and we will schedule a makeup
- **Course evaluations** are due **Sunday, 12/8 @ 11:59pm**
  - **Section evaluations** are due **Monday, 12/9 @ 11:59pm**

# Final exam format

- Similar to midterm exam, but longer
  - A sample final is available on Ed
- 135 minutes
- You will be given a standard reference sheet
  - Is expanded from the midterm, attached to sample final on Ed
- You may bring one sheet of double sided 8.5x11” paper containing your own handwritten notes.
  - Must write name, student number, and UW NetID
  - Must turn in with exam
  - If you want to access your midterm notes sheet, go to Prof. Beame’s OH

# Today's plan

1. (35 min) 6 stations around the room with practice problems  
(focused on second half of course, but exam is cumulative)
  - Station 1: Short answer
  - Station 2: Dynamic programming\*
  - Station 3: Network flow
  - Station 4: Linear programming\*
  - Station 5: Reduction
  - Station 6: Bonus problem
1. (10 min) Go over some of these problems

*\*the problem at this station was an extra problem on a previous section handout*

# Problems

**1**

**2**

**3**

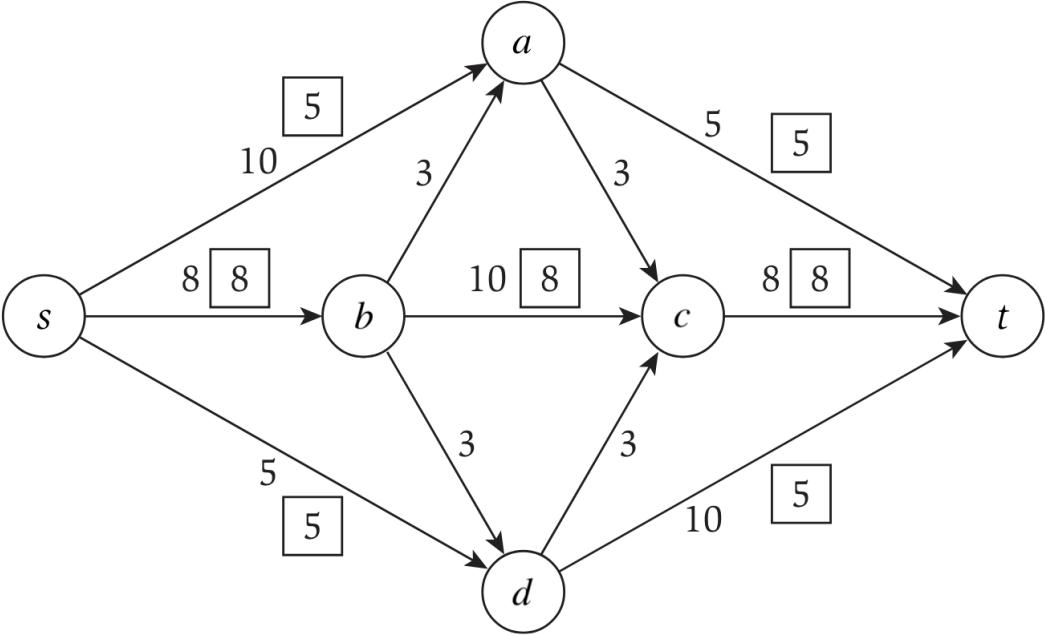
**4**

**5**

**6**

# Problem 1 – Short answer

In the network flow below, is the depicted flow a maximum flow?



## Problem 1 – Short answer

Recall Interval Scheduling: Given a collection of intervals and an integer  $k$ , determine if the collection contains at least  $k$  nonoverlapping intervals.

- i. Does Interval Scheduling  $\leq_p$  Vertex Cover?

## Problem 1 – Short answer

Recall Interval Scheduling: Given a collection of intervals and an integer  $k$ , determine if the collection contains at least  $k$  nonoverlapping intervals.

ii. Does Independent Set  $\leq_p$  Interval Scheduling ?



# Problem 1 – Short answer

A greedy attempt at Set Cover is:

```
while there exists an uncovered object do  
    choose a set that covers the most number of still-uncovered objects
```

Suppose you are given an instance where every set contains exactly 2 elements. Then this algorithm returns a set cover that is at most a factor 2 larger than the minimum.

## Problem 2 – Dynamic programming

Given two strings,  $s = s_1, \dots, s_m$  with length  $m$  and  $t = t_1, \dots, t_n$  with length  $n$ , find the length of their longest common subsequence.

## Problem 3 – Network flows

The bank has  $C_j$  of currency  $j$ , and the exchange rate is  $R_j$  of currency  $j$  for every 1 Franc. Trader  $i$  has  $T_i$  Francs to convert and is willing to convert between  $L_{ij}$  and  $H_{ij}$  of their Francs to currency  $j$ . Determine if the bank can satisfy all requests, and if so, how to maximize the amount of Francs it collects.

## Problem 4 – Linear programming

There are  $k$  groups and  $m_i$  voters in group  $i$ , of which  $a_i$  are already voting for you. If you spend \$1000 advertising issue  $j$ , then  $d_{ij}$  more voters in group  $i$  will vote for you. Determine the minimum spending so that at least half of each group votes for you.

## Problem 5 – Reduction

A Hamiltonian path/cycle is a path/cycle that visits every vertex exactly once. Suppose that HamiltonianPath is NP-hard. Show that HamiltonianCycle is NP-hard.

## Problem 6 – Bonus problem

In a country with  $n$  states and  $p_i$  people voting in state  $i$ , the winner of state  $i$  receives  $v_i$  electoral college votes. In a two-candidate election with no state-level ties, determine the minimum percent of the total popular vote necessary to win at least  $V = \lfloor (\sum_i v_i) / 2 \rfloor + 1$  electoral votes.