

# CSE 421 Lecture 11

## 1 Master Theorem

**Lemma 1.** Let  $S = 1 + x + x^2 + \dots + x^d$ . Then, we have

$$S = \begin{cases} \Theta_x(1) & \text{if } x < 1, \\ d + 1 & \text{if } x = 1, \\ \Theta_x(x^d) & \text{if } x > 1 \end{cases}$$

where  $\Theta_x$  means the constant depends on  $x$ .

*Proof.* Let  $S = 1 + x + x^2 + \dots + x^d$ . Then, we have

$$\begin{aligned} xS &= x + x^2 + \dots + x^{d+1}, \\ xS - S &= x^{d+1} - 1. \end{aligned}$$

Hence, we have

$$S = \frac{x^{d+1} - 1}{x - 1}.$$

If  $x < 1$ , we have  $S = \frac{1-x^{d+1}}{1-x}$ . Note that  $1 \leq S \leq \frac{1}{1-x}$  where both left and right are independent to  $d$ , we have  $S = \Theta_x(1)$ .

If  $x = 1$ , we have  $S = d + 1$  by the definition of  $S$ .

If  $x > 1$ , we have  $x^d \leq S \leq \frac{x}{x-1} \cdot x^d$ . Hence, we have  $S = \Theta_x(x^d)$ . □

**Theorem 2.** Given  $a \geq 1, b > 1, c > 0$  and  $k \geq 0$ . Suppose that  $T(n) = aT(\frac{n}{b}) + cn^k$  for all  $n \geq b$ , then

1. If  $a < b^k$ , then  $T(n) = \Theta(n^k)$ .
2. If  $a = b^k$ , then  $T(n) = \Theta(n^k \log n)$ .
3. If  $a > b^k$ , then  $T(n) = \Theta(n^{\log_b a})$ .

*Proof.* As we argued in the slide, we have

$$T(n) = cn^k \sum_{i=0}^{\log_b n} \left(\frac{a}{b^k}\right)^i.$$

If  $a < b^k$ , the previous lemma shows

$$T(n) = cn^k \Theta(1) = \Theta(n^k).$$

If  $a = b^k$ , the previous lemma shows

$$T(n) = cn^k \Theta(\log_b n) = \Theta(n^k \log n).$$

If  $a > b^k$ , the previous lemma shows

$$T(n) = cn^k \Theta\left(\frac{a}{b^k}\right)^{\log_b n}.$$

Note that  $b^{\log_b n} = n$  and hence

$$T(n) = \Theta\left(n^k \frac{a^{\log_b n}}{n^k}\right) = \Theta(a^{\log_b n}) = \Theta((b^{\log_b a})^{\log_b n}) = \Theta((b^{\log_b n})^{\log_b a}) = \Theta(n^{\log_b a}).$$

□