

## Approximation Ratio

If I'm trying to find the minimum vertex cover, then to have a 2-approximation, I need to show:

If I'm trying to find the maximum clique, then to have an  $n$ -approximation, I need to show:

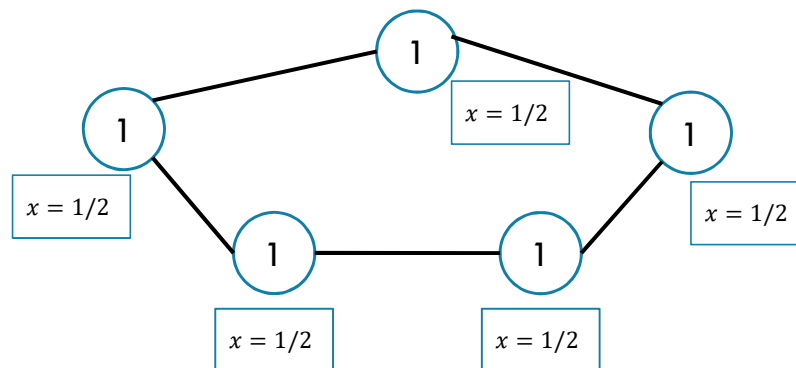
## Summary

Coping with NP-hardness.

1. Understand your problem really well (make sure you're not solving an easy special case).
2. Prove the problem really is NP-hard.
3. Try a band-aid (SAT library, Integer programming library, etc.)
4. Try to find a good-enough exponential time algorithm or an approximation algorithm.

## Non-Bipartite

What if our original graph isn't bipartite?



The LP finds a fractional vertex cover of weight 2.5

There's no "real"/integral VC of weight 2.5. – lightest is weight 3.

There's a "gap" between integral and fractional solutions.

## So, what if the graph isn't bipartite?

Big idea:  
Just round!

If  $x_u \geq \frac{1}{2}$ , round up to 1.

If  $x_u < \frac{1}{2}$ , round down to 0

Two questions – is it a vertex cover? How far are we from the true minimum?

[Pollev.com/robbie](https://pollev.com/robbie)

Minimize  $\sum w(u) \cdot x_u$

Subject to:

$x_u + x_v \geq 1$  for all  $(u, v) \in E$

$0 \leq x_u \leq 1$  for all  $u$ .