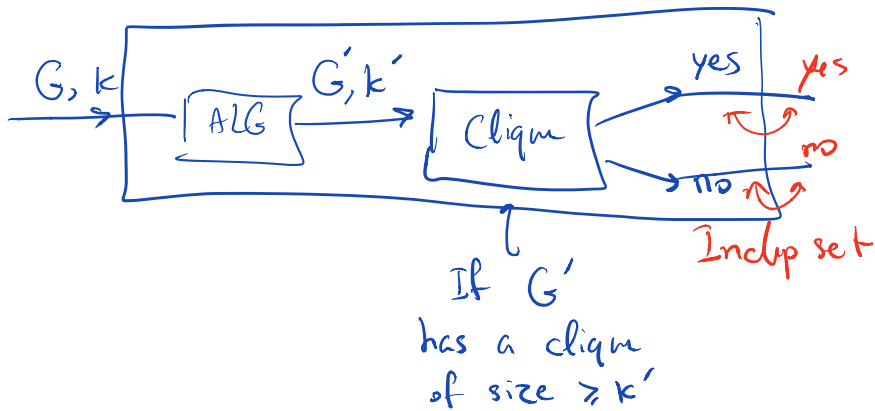


Indep set \leq^p Clique



$\left\{ \begin{array}{l} G' \equiv \text{comp of } G; (u,v) \text{ edge of } G' \text{ iff } (u,v) \text{ not edge of } G. \\ k' = k \end{array} \right.$
 in polytime.

Claim: G has indep set size $\geq k$ iff G' has clique size $\geq k$.

Pf. G has indep set $|S| \geq k \Rightarrow G'$ has clique $|S'| \geq k$.

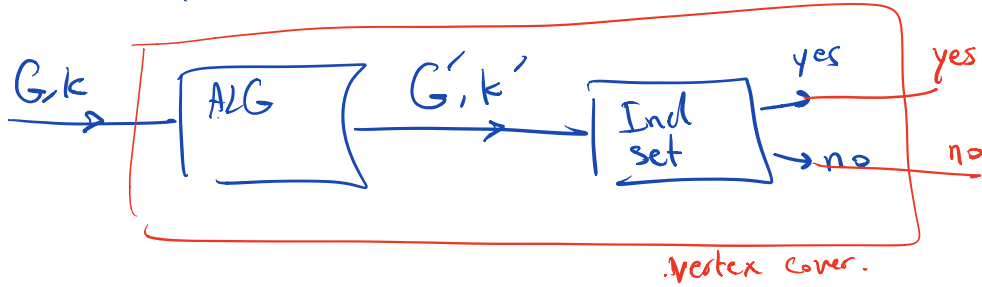
If S indep set then S is a clique in $G' \Rightarrow$ let $S' \leftarrow S$ in G

G' has a clique $|S'| \geq k \Rightarrow G$ has indep set $|S| \geq k$.
 $S = S'$

Vertex Cover \leq^p Indep set

Vertex Cover Given G, k does it have vertex cover S , with

$$|S| \geq k.$$



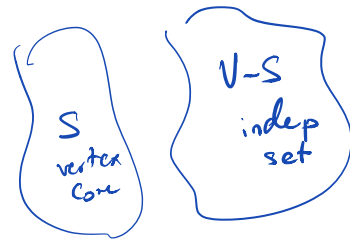
$$G' = G, \quad k' = n - k$$

Claim: S is a vertex cover of G iff $V-S$ is an indep set of G .

Implying G has vertex cover $|S| \leq k \iff G$ has indep set $|S| \geq n - k$

Pf of claim: S vertex cover $\implies V-S$ indep set.

\implies By contradict: if $\exists e \in V-S$ then e is not covered by S so S is not a vertex cover.



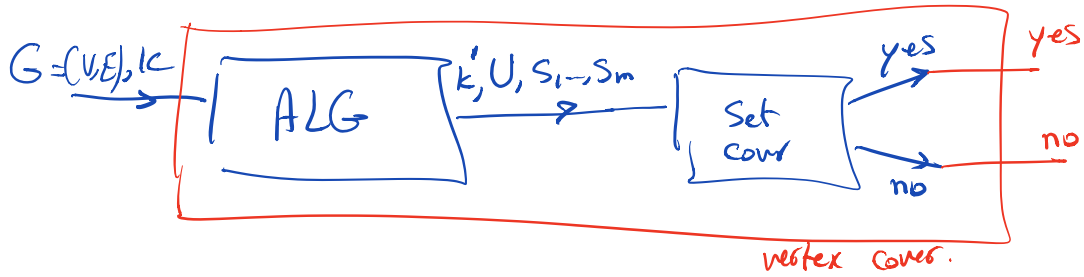
\leftarrow $V-S$ indep set $\implies S$ vertex cover.
By contradiction: If S not a vertex cover $\implies \exists e$ not covered by S . So both endpoints are in $V-S$. So $V-S$ not indep set.

HW8-P1 \min vertex cover is polytime in Bipartite graphs

\implies max indep set is also polytime solvable in Bipartite graphs.

Set Cover: Given ground elements U , $S_1, \dots, S_m \subseteq U$ and integ k , is there $\leq k$ sets to cover all U .

Vertex Cover \leq_p Set Cover



$U = E, k' = k \quad V = \{v_1 - v_n\}$

$S_{v_i} = \{e \in E : e \text{ incident to } v_i\}$

Claim: G has vertex cover $|S| \leq k$ iff $U, k, S_{v_1} - S_{v_n}$ is yes there are $\leq k$ sets to cover U .

Pf: G has vertex cover $|S| \leq k, S \subseteq V$

$\Rightarrow U, k, S_{v_1} - S_{v_n}$ is yes

Just include S_{v_i} for all $v_i \in S$.

$|S| \leq k \Rightarrow$ at most k sets

S covers all edges, and S_{v_i} has edges adj to v_i implies chosen sets cover $U = E$

$U, k, S_{v_1} - S_{v_n}$ is yes $\Rightarrow G$ has vertex cover $|S| \leq k$

$\{S_{v_1} - S_{v_k}\}$ covers all of U . Then $v_1 - v_k$ form a vertex cover of G .

