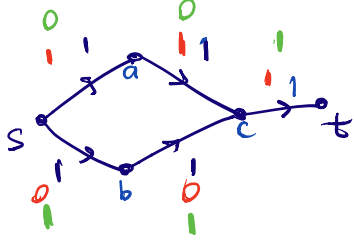
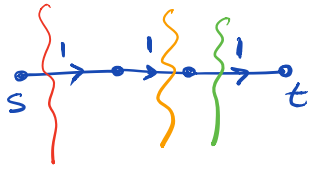


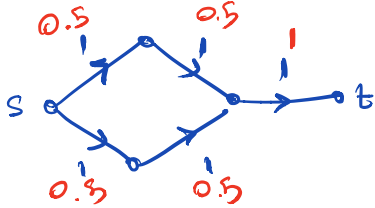
1) Given G , Is max-flow unique? no



2) Is Min $\{s,t\}$ Cut unique? no



3) Is Max flow ^{always} integral? no (assuming ^{all} capacities are integral)



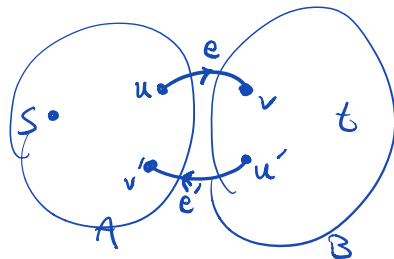
Claim: Supp no more augmentation for $f \Rightarrow \exists$ s-t cut (A,B) s.t. $\text{cap}(A,B) = v(f)$.

$A = \{v : v \text{ is reachable from } s \text{ in } G_f\}$. $B = \bar{A}$

$s \in A$ (since s is reachable from s)

$t \in B$ B/C no more augmentation for f . G

No edge from $A \rightarrow B$ in G_f

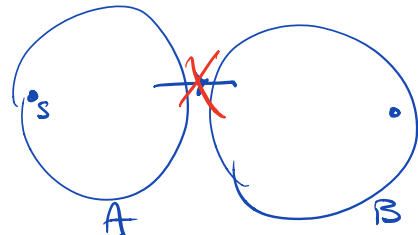


$e \notin G_f \Rightarrow c_e = f(e)$

$e' \notin G_f \Rightarrow f(e') = 0$

[If not $e'^R \in G_f$ and contradiction]

G_f



Flow value Lemma

$$v(f) = \sum_{e \text{ out of } A} f(e) - \sum_{e' \text{ into } A} f(e') = \sum_{e \text{ out of } C_e} f(e) = \text{cap}(A, B).$$

$$\max_{\text{matching}} G = \max_{\text{flow}} H.$$

i) $\max_{\text{matching}} G \leq \max_{\text{flow}} H.$

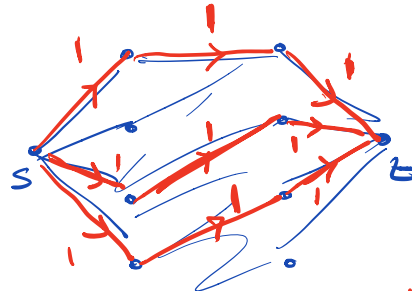
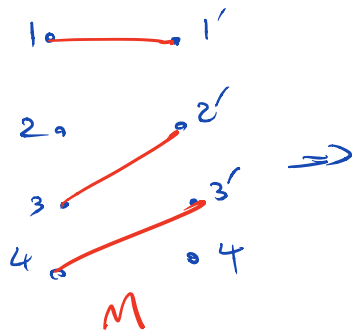
pf.

Technique start with M to be max matching of G .



Construct a flow f st. $v(f) = |M|$

$$\Rightarrow \max_{\text{matching}} G \leq \max_{\text{flow}} H.$$



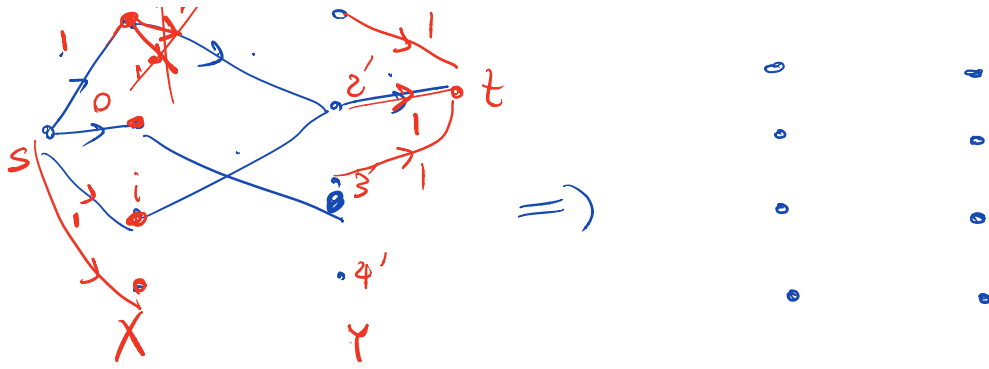
feasible flow: B/C in-flow = out-flow for all except s, t

$$v(f) = |M|.$$

$$\max_{\text{flow}} \leq \max_{\text{matching}}.$$

Say f is an integer max-flow \Rightarrow construct a matching M st. $|M| \geq v(f)$.
(we proved it always exist)





$\Rightarrow f \text{ integ} \Rightarrow f((s, i)) = 0/1$ for all $i \in X$

$f((i', t)) = 0/1$ for all $i' \in Y$

So every $i \in X$ has either 0 or 1 incoming flow
 \Rightarrow either 0 or 1 edges with flow 1 leave i .

\Rightarrow Define $M = (i, i')$ s.t. $f((i, i')) = 1$

This is a matching B/C every vertex in X or Y
 has at most 1 incoming or outgoing flow.

\Rightarrow Also $|M| = v(f)$