No HW is due in midterm week.

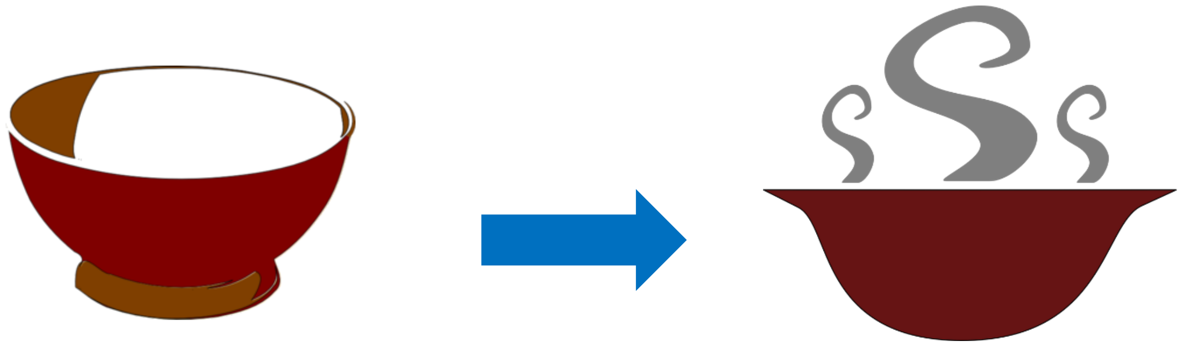
CSE 421

Dijkstra's Algorithm, Divide and Conquer

Shayan Oveis Gharan

Boiling Water Example

Q: Given an empty bowl, how do you make boiling water?



A: Well, I fill it with water, turn on the stove, leave the bowl on the stove for 20 minutes. I have my boiling water.

Q: Now, suppose you have a bowl of water, how do you make boiling water?

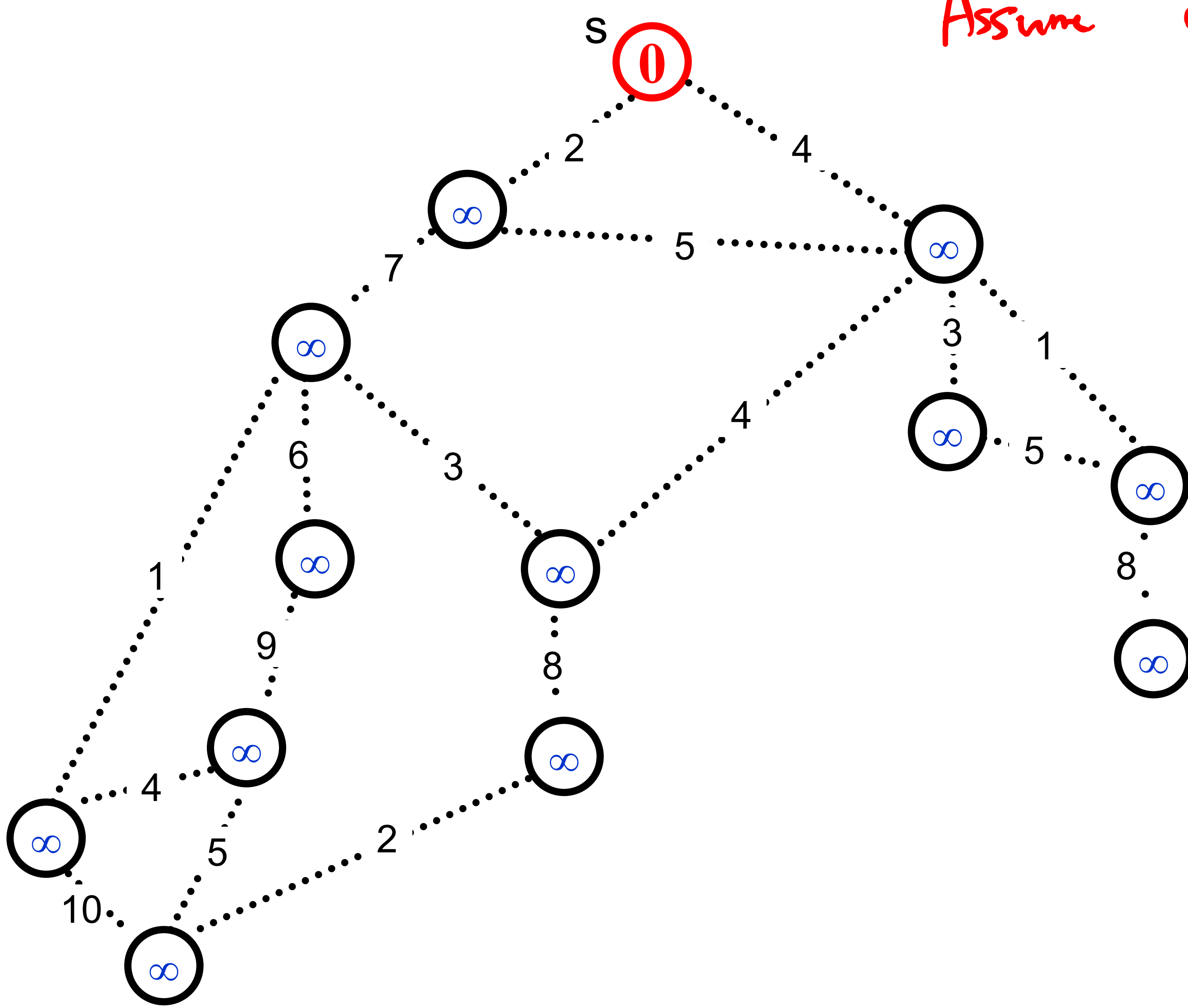
A: First, I pour water away, now
I have an empty bowl and
I have already solved this!

Lesson: Never solve a problem twice!

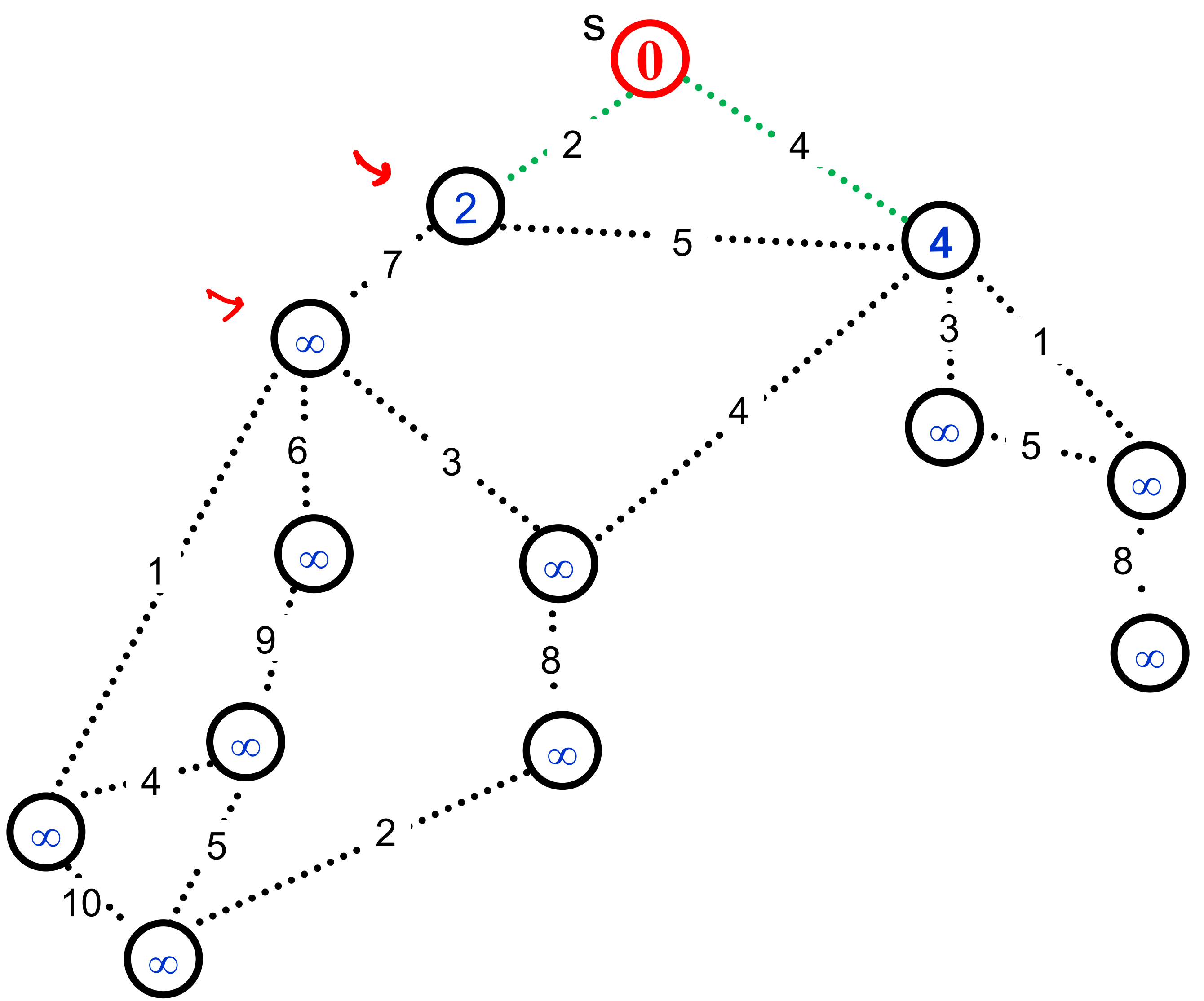
Related Q: Try to do max-spanning tree using min-spanning tree

Dijkstra's Algorithm: Example

Assume $c_e \geq 0$



Dijkstra's Algorithm: Example



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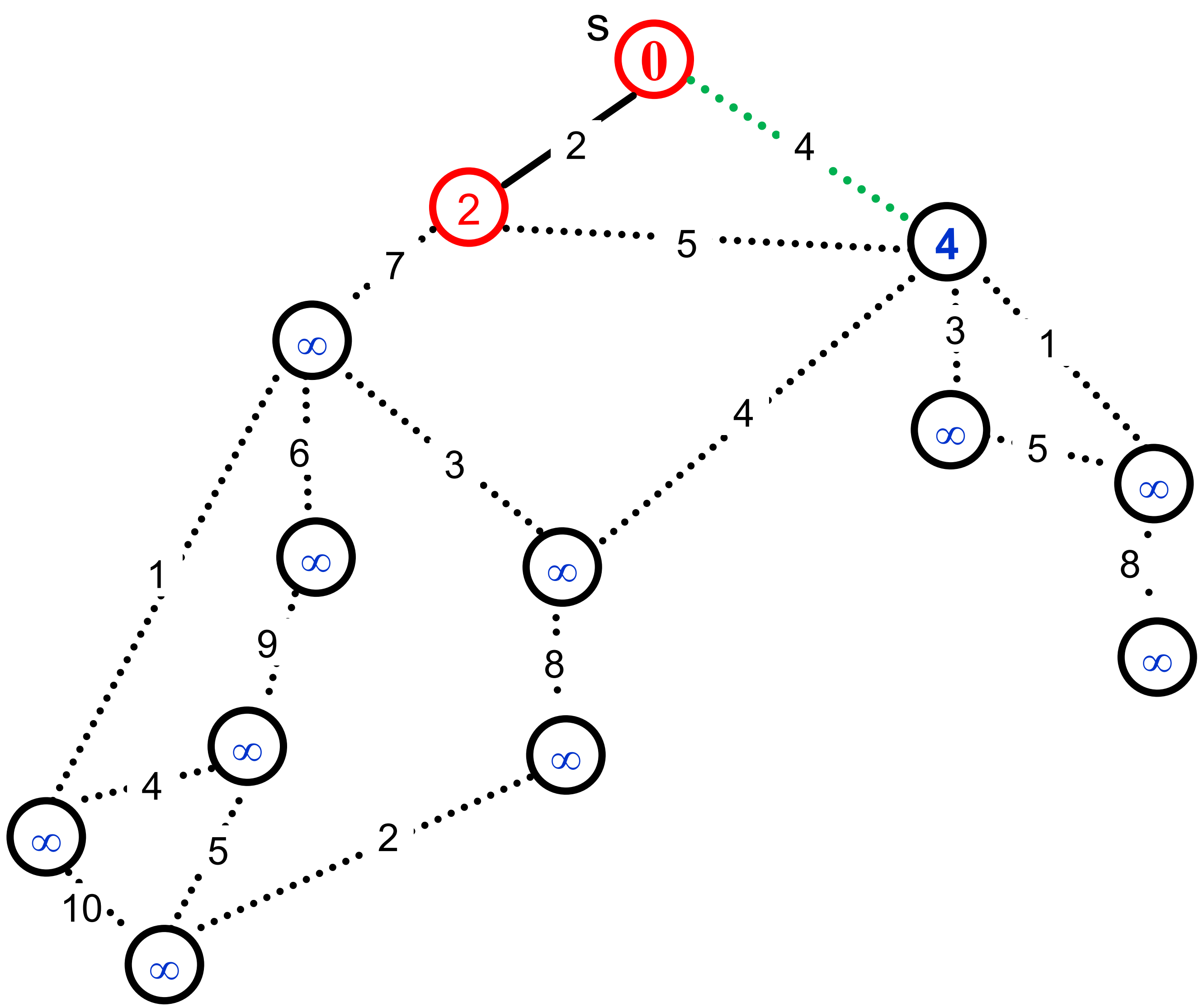

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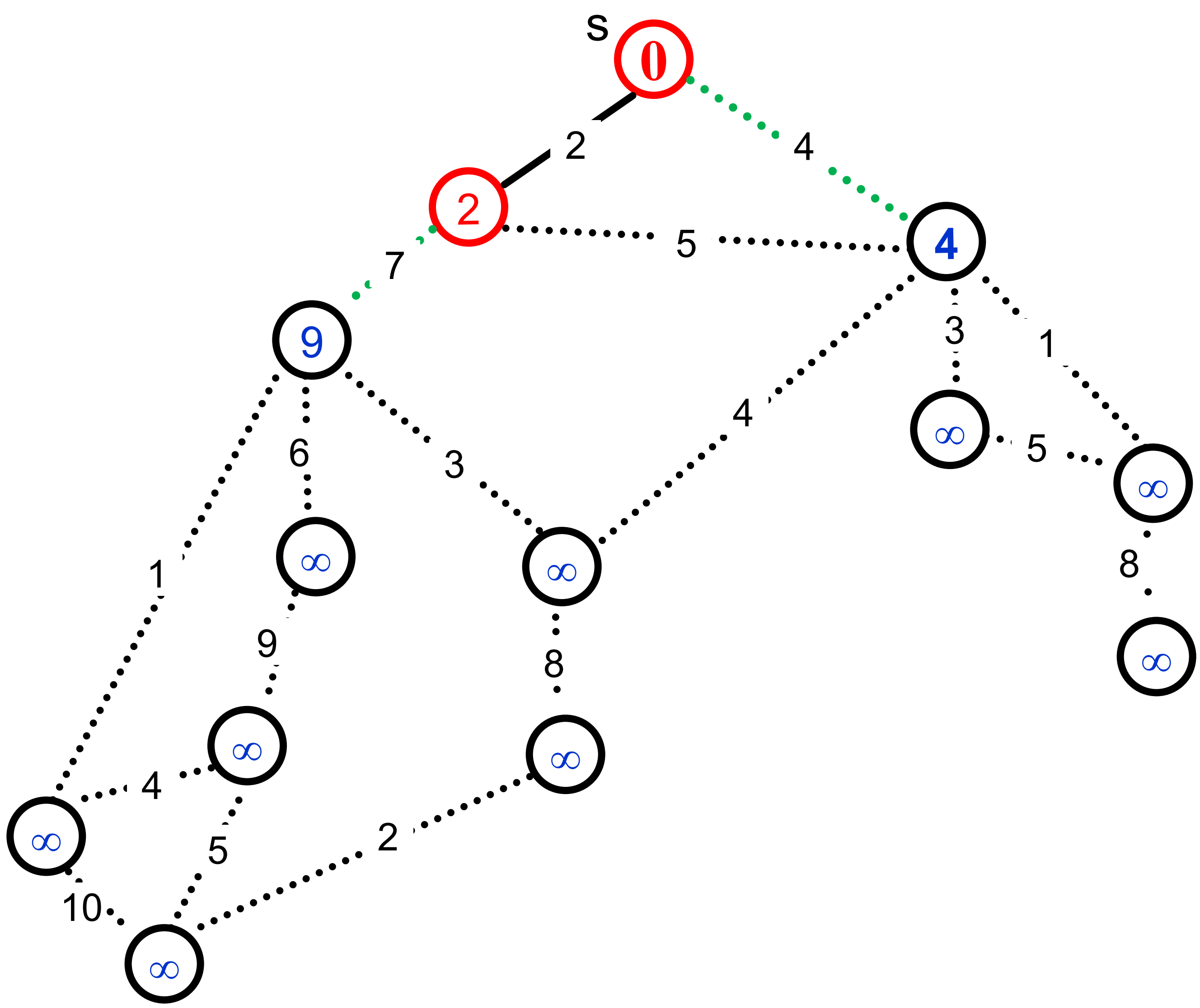

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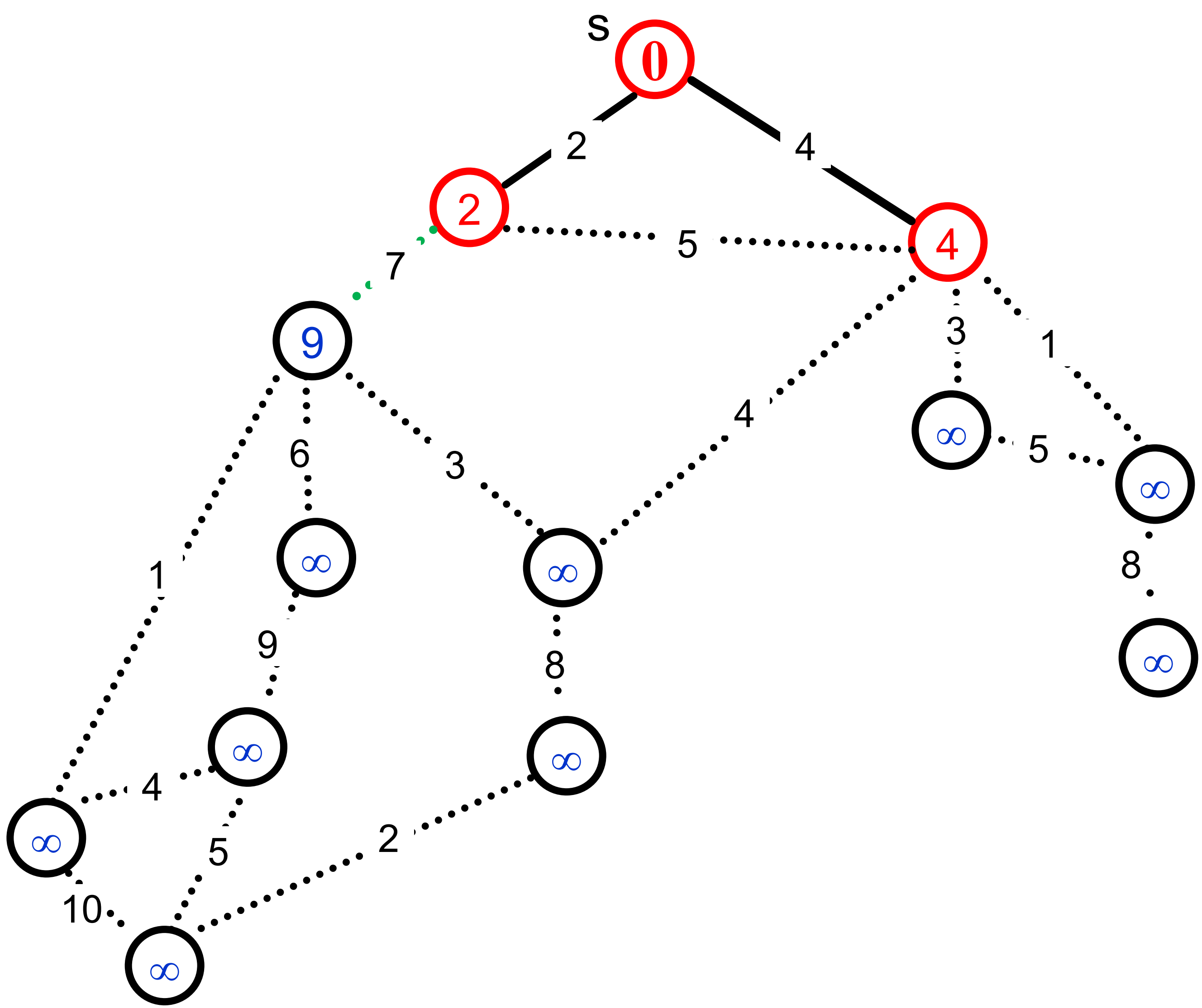

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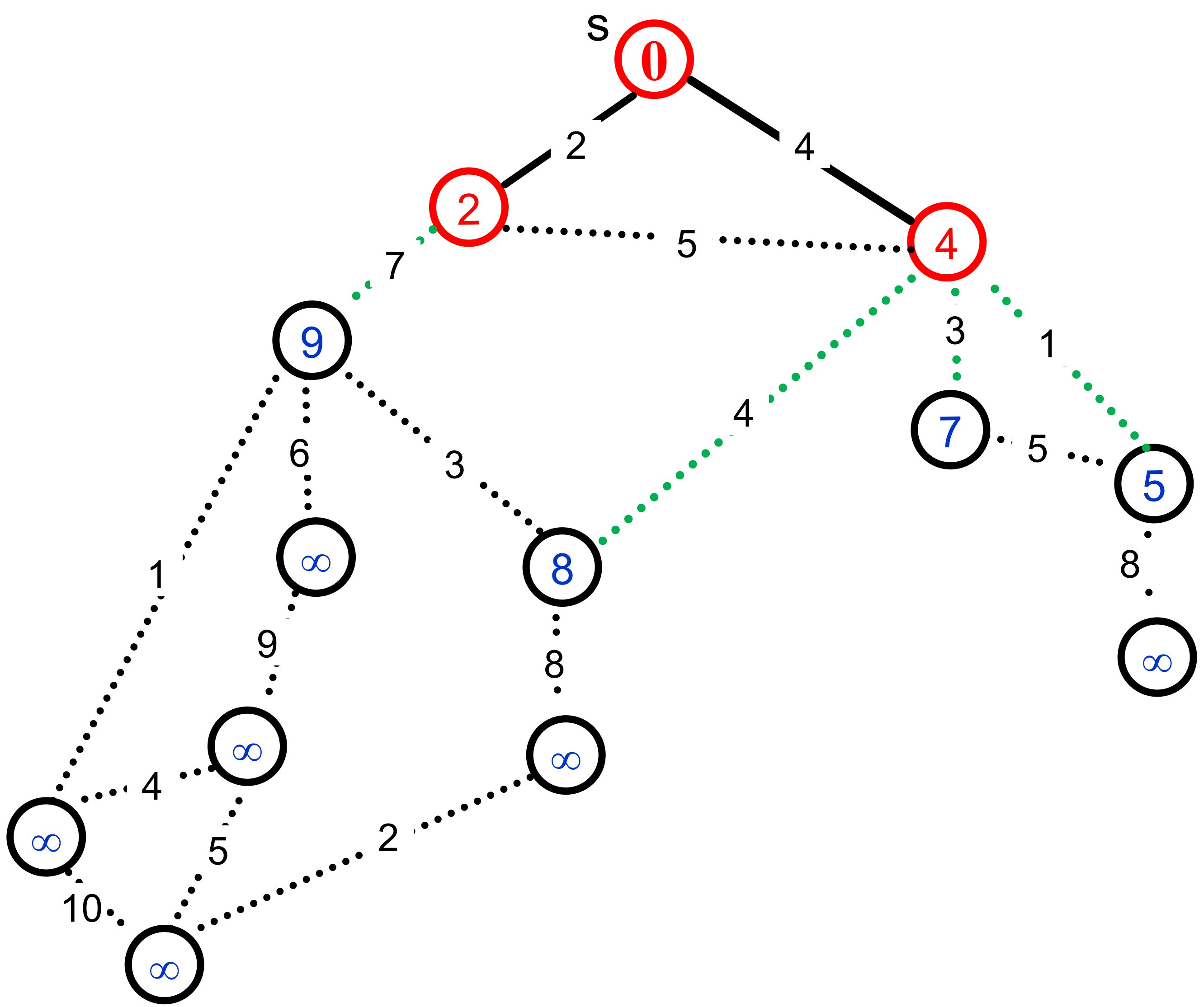

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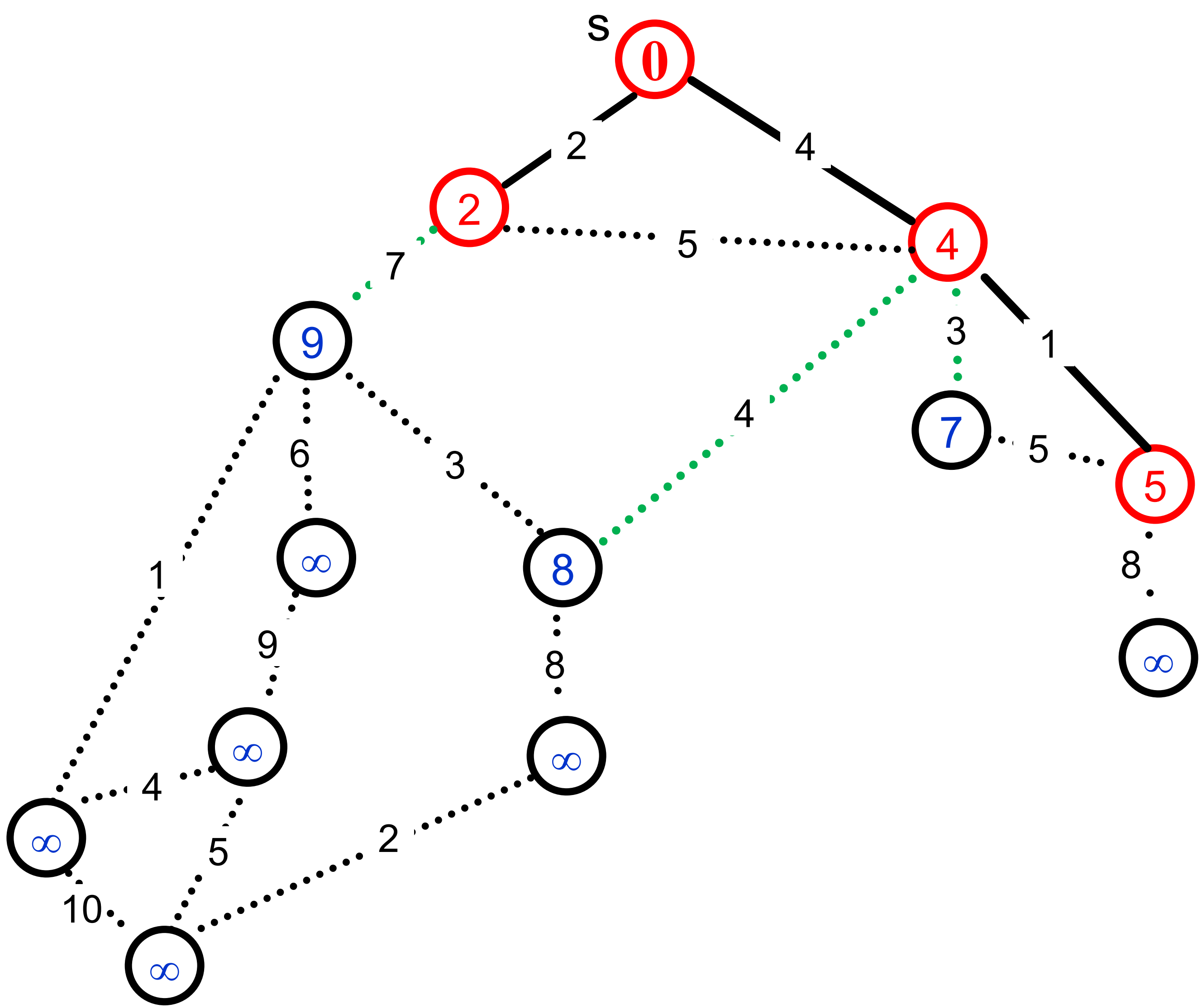

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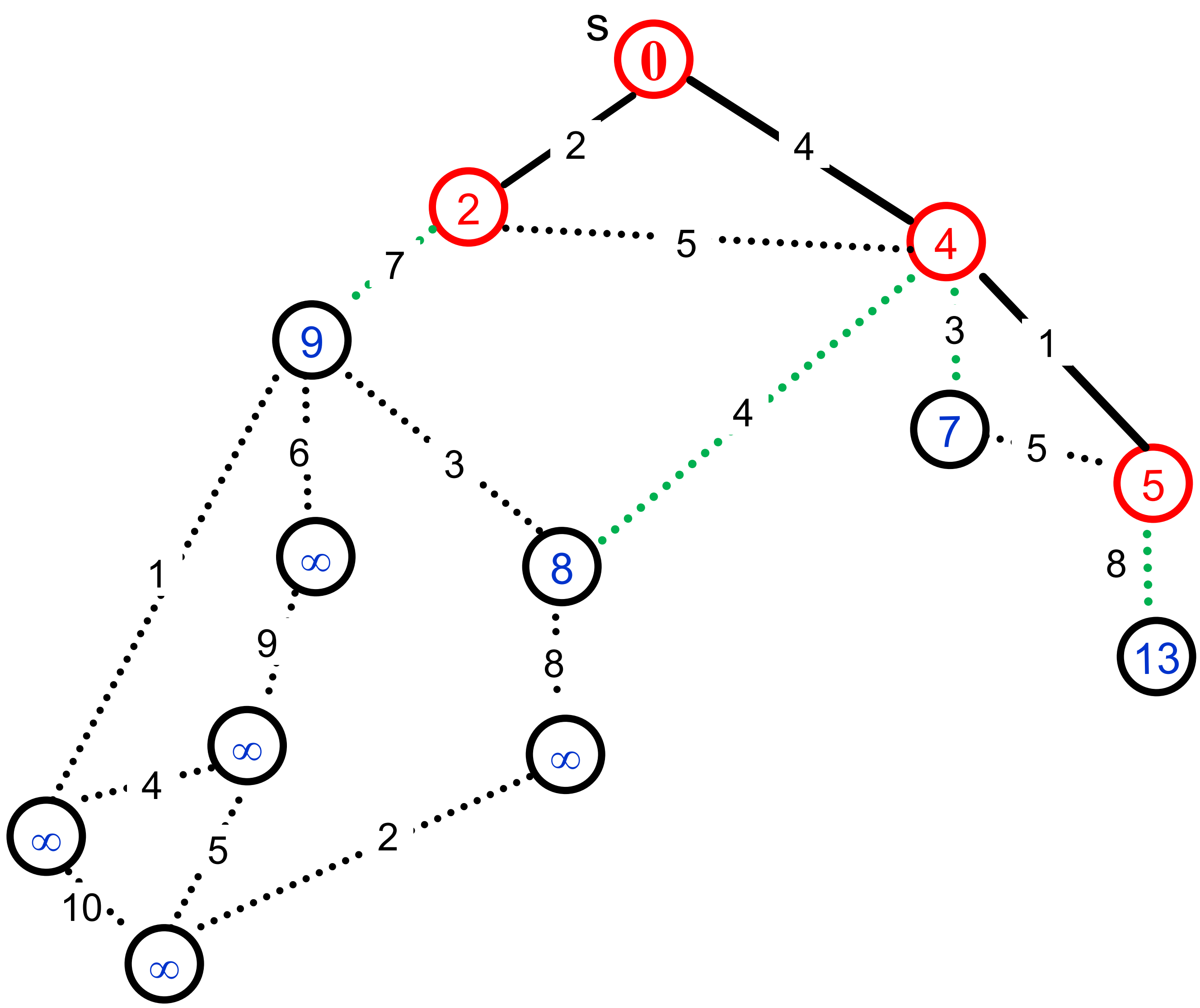

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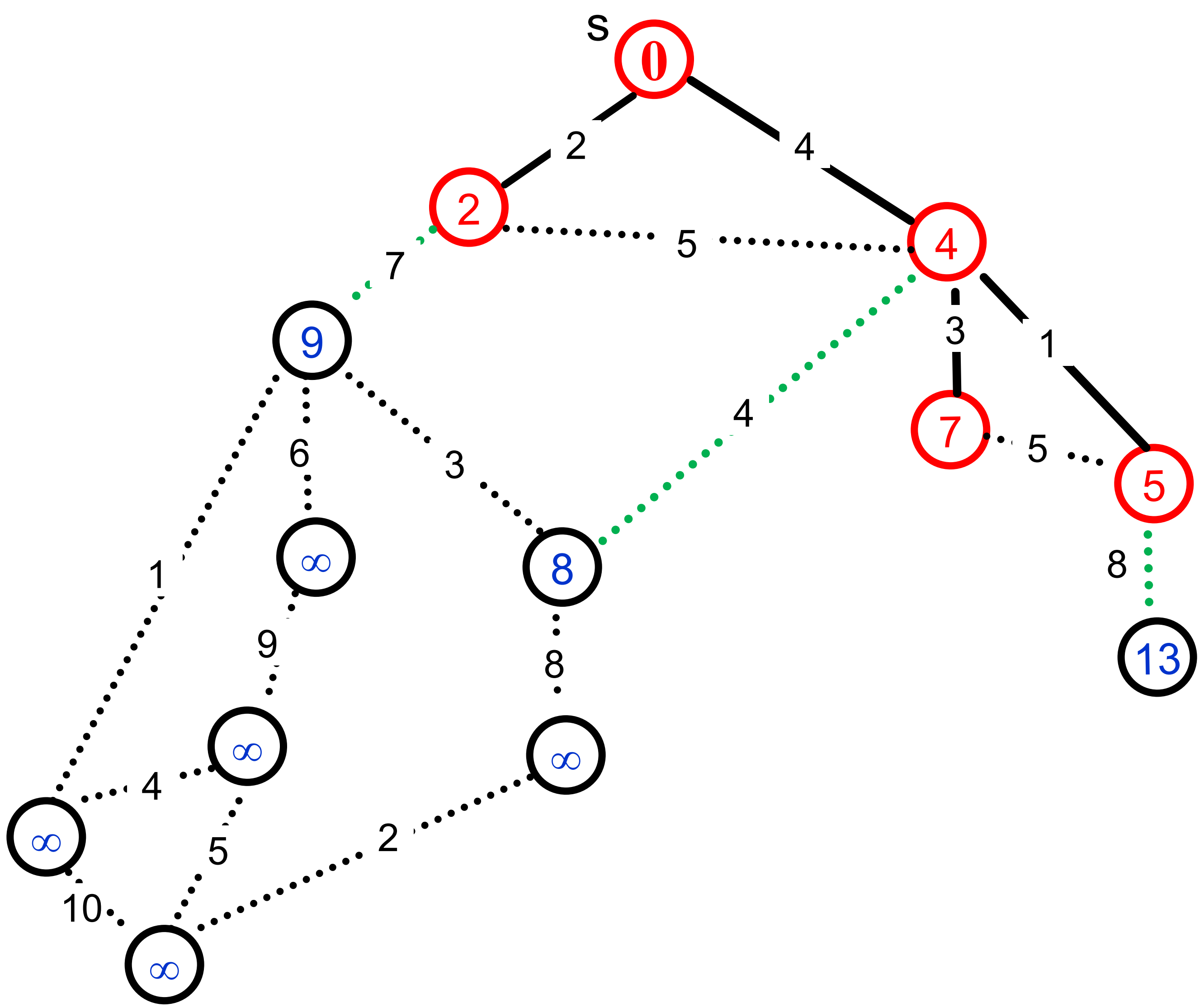

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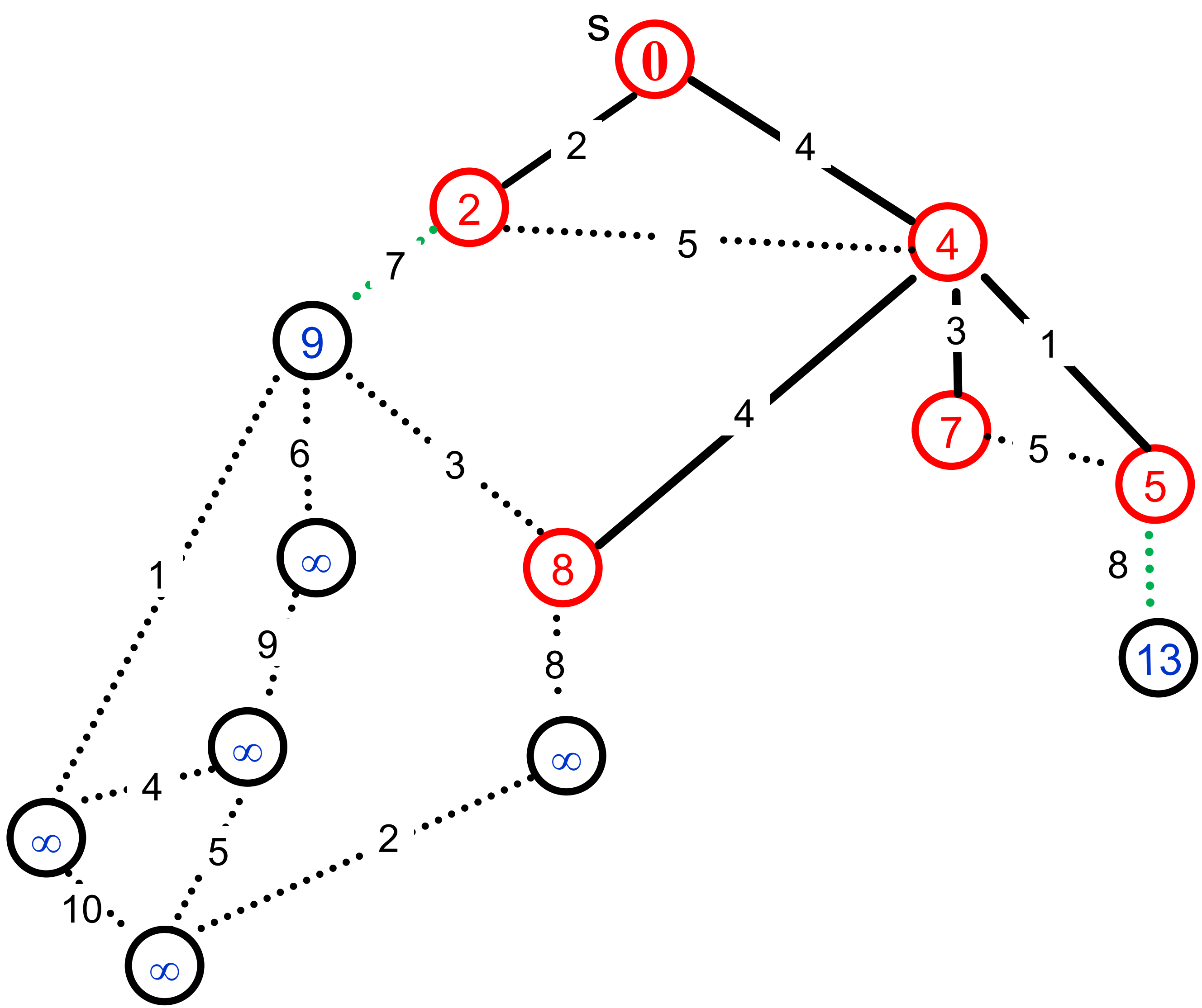

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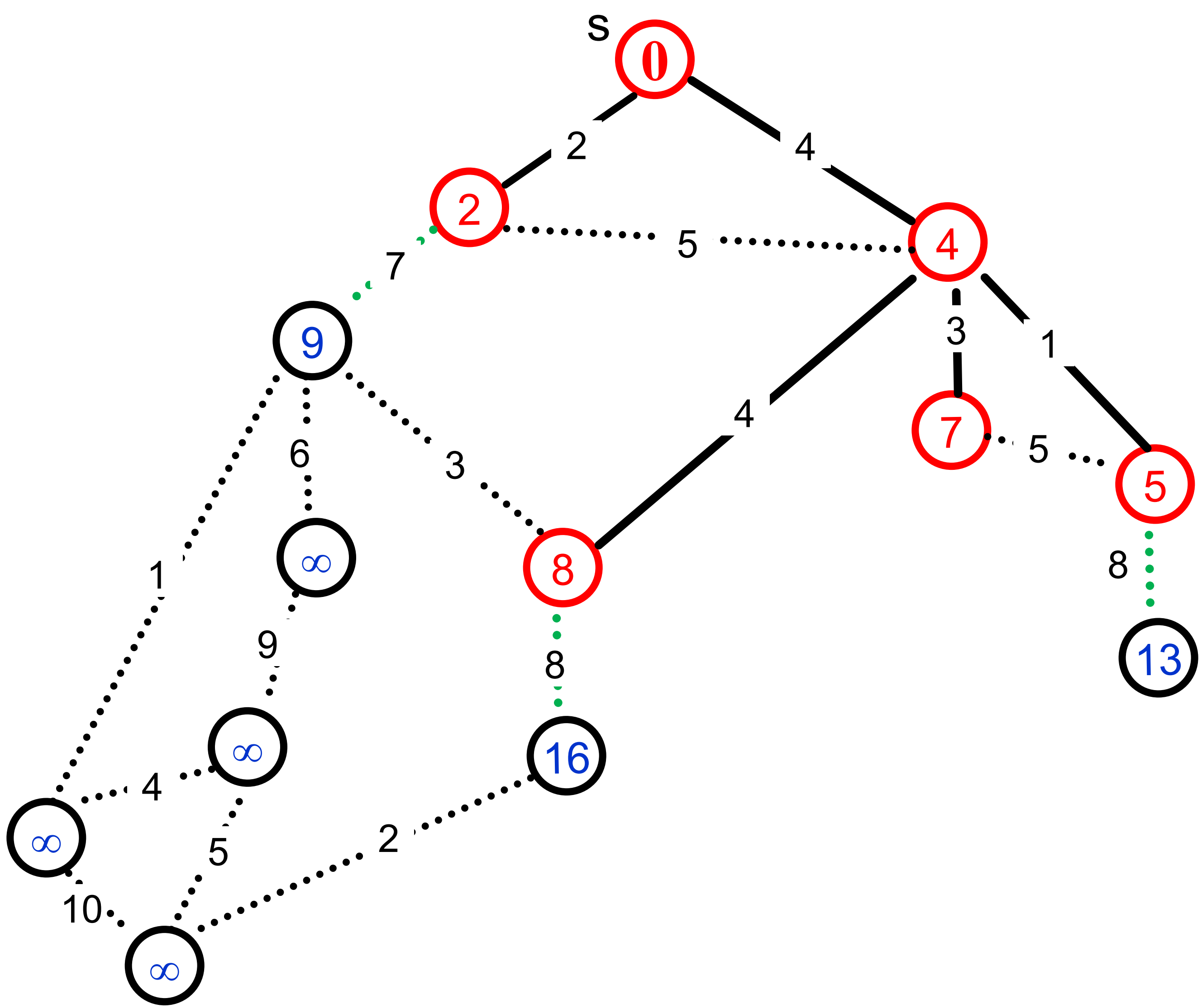

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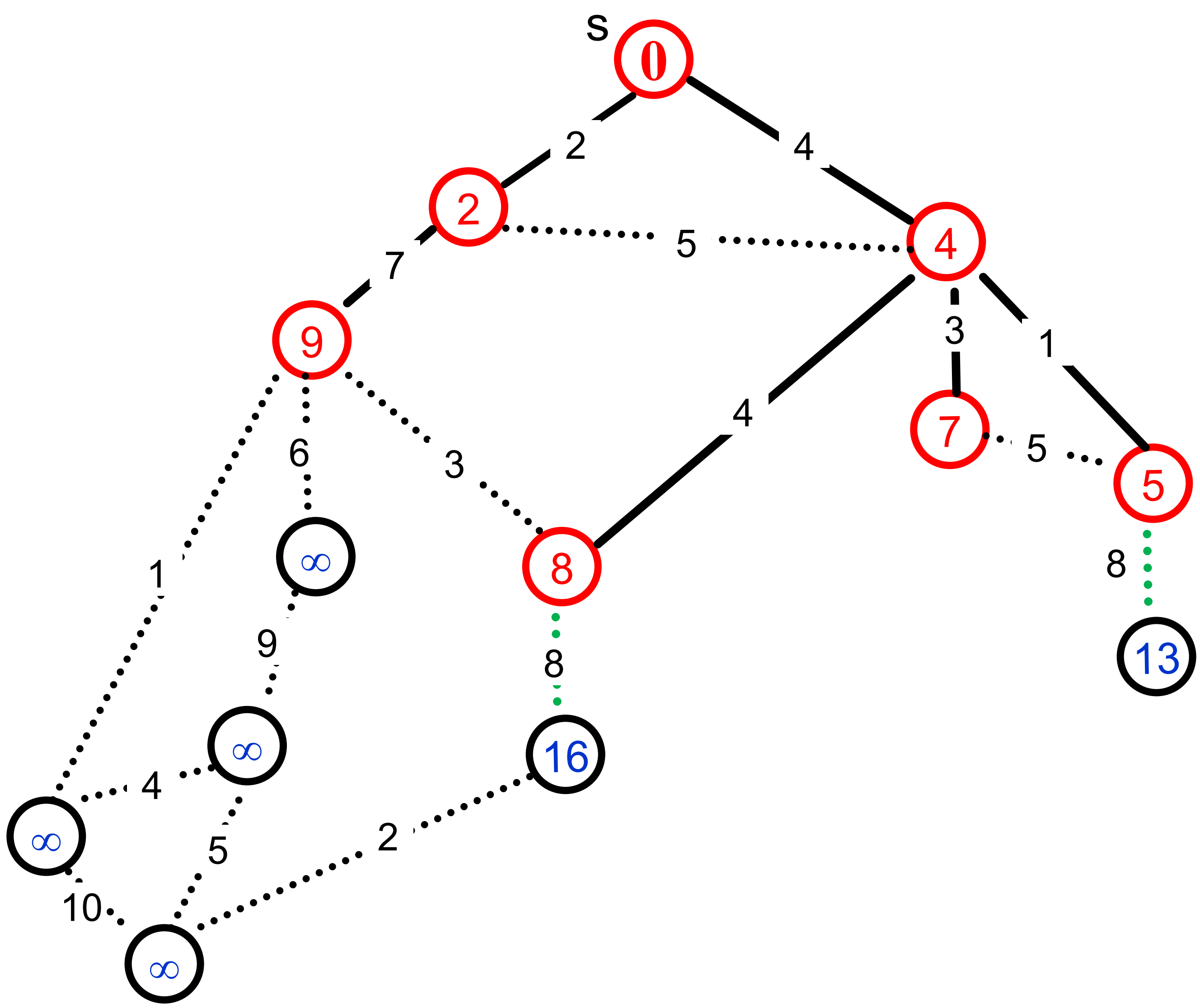

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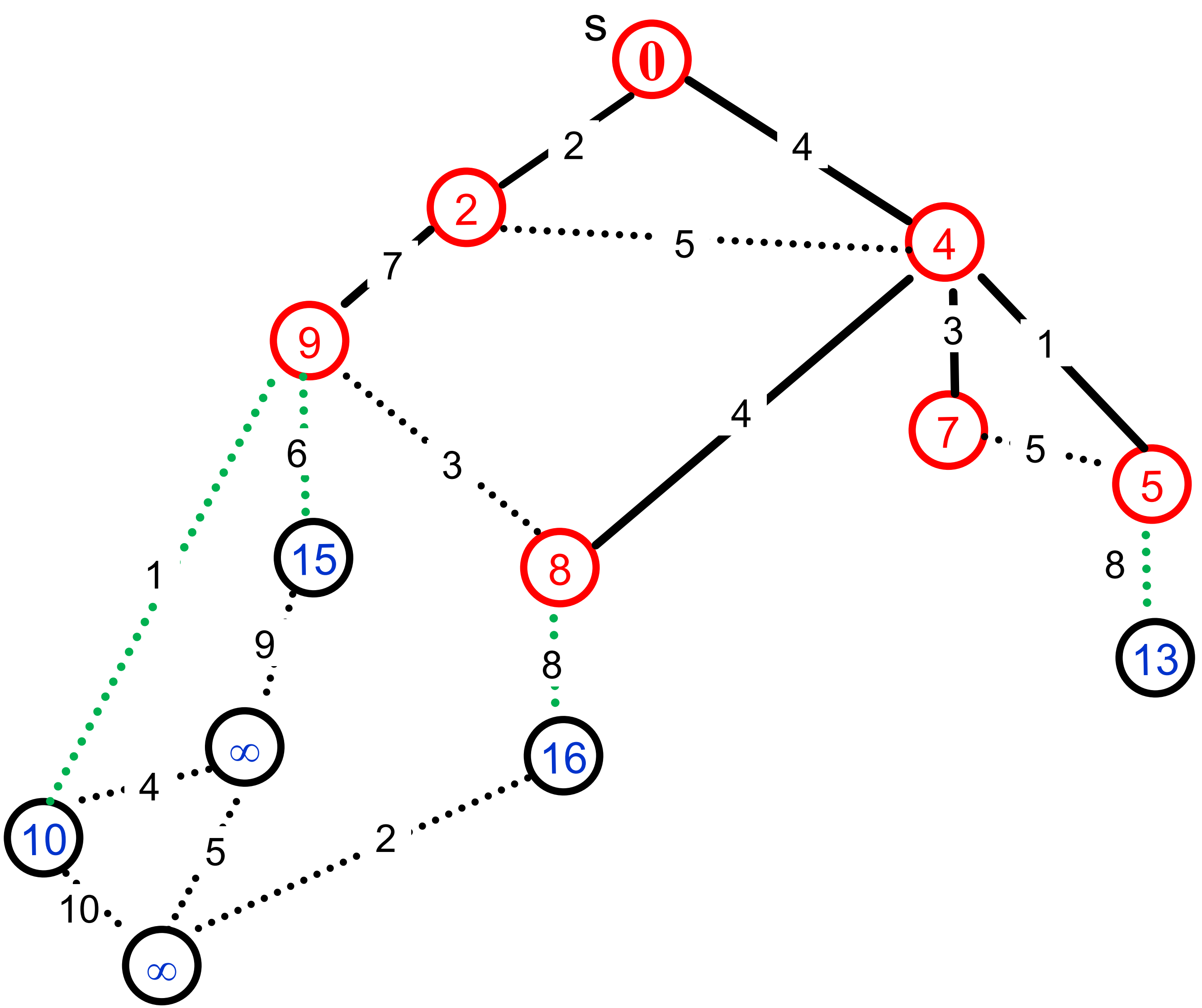

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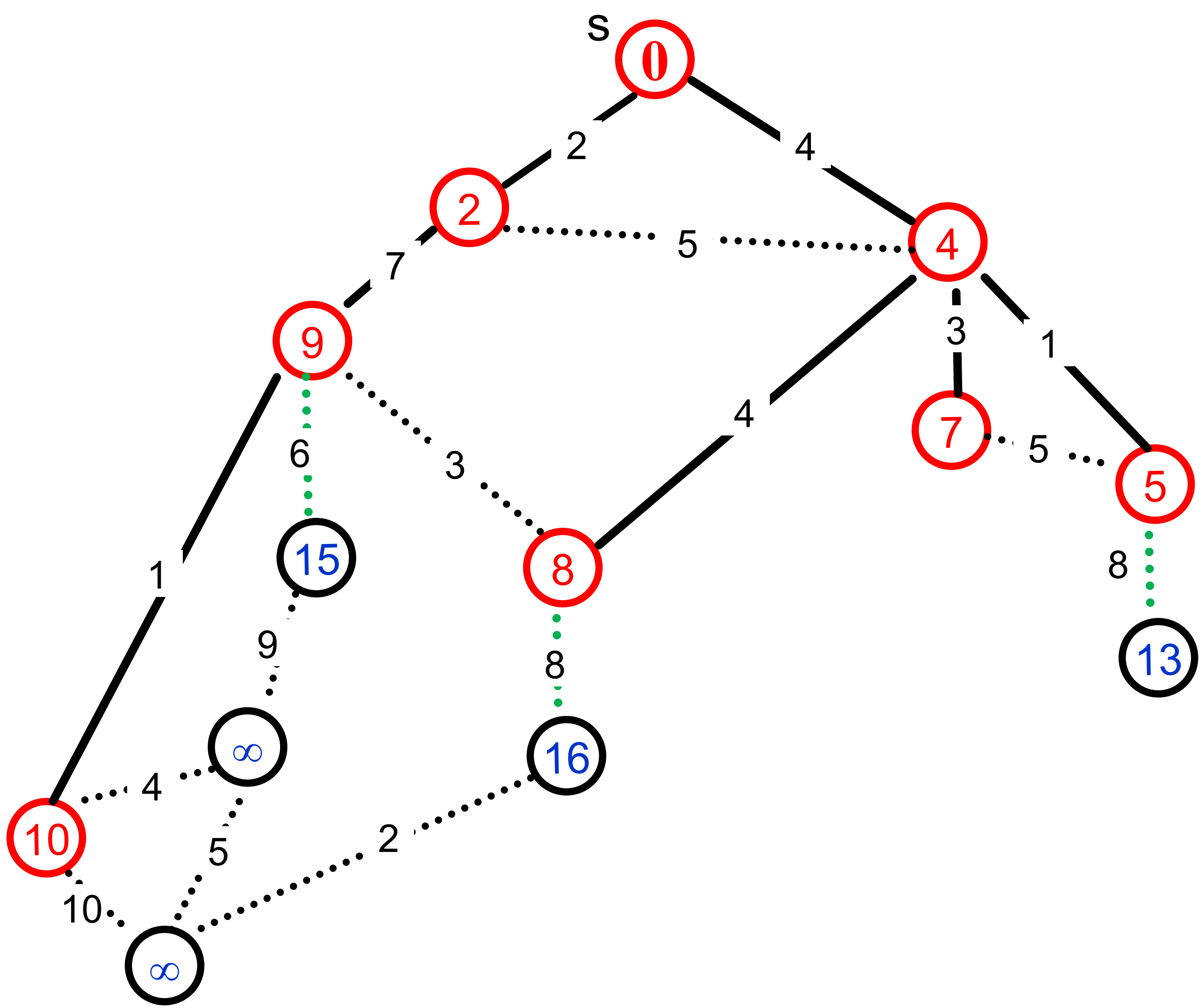

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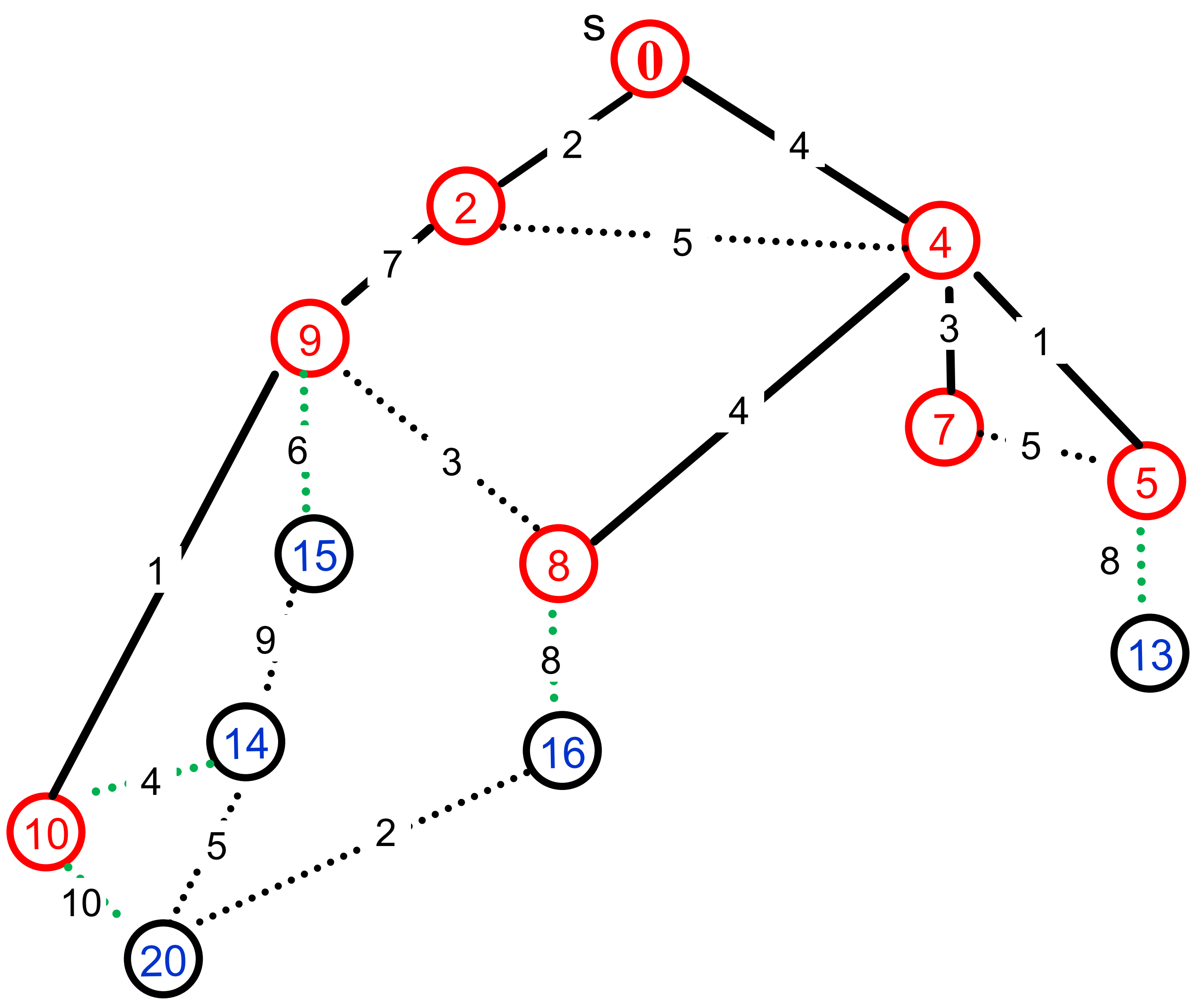

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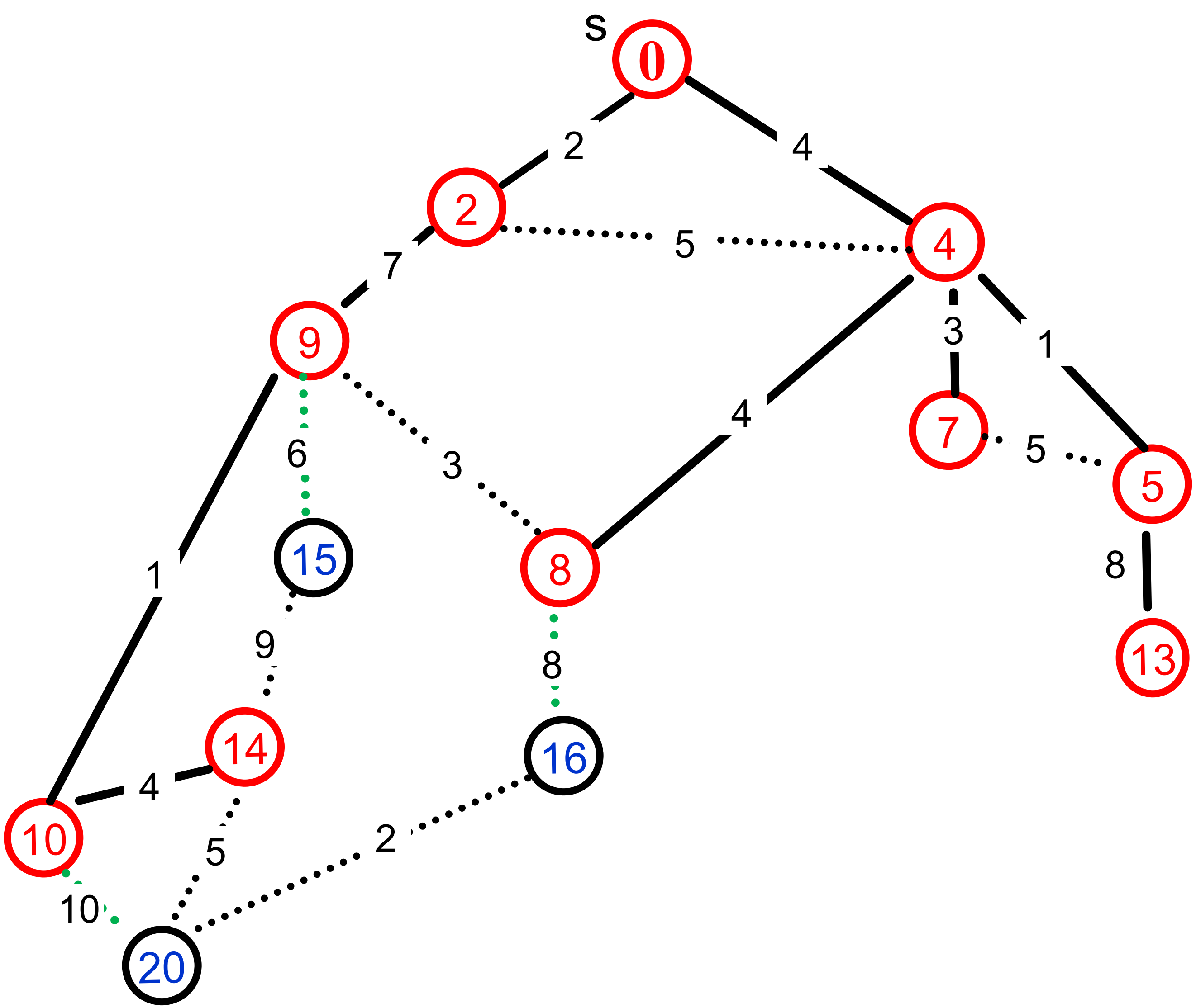

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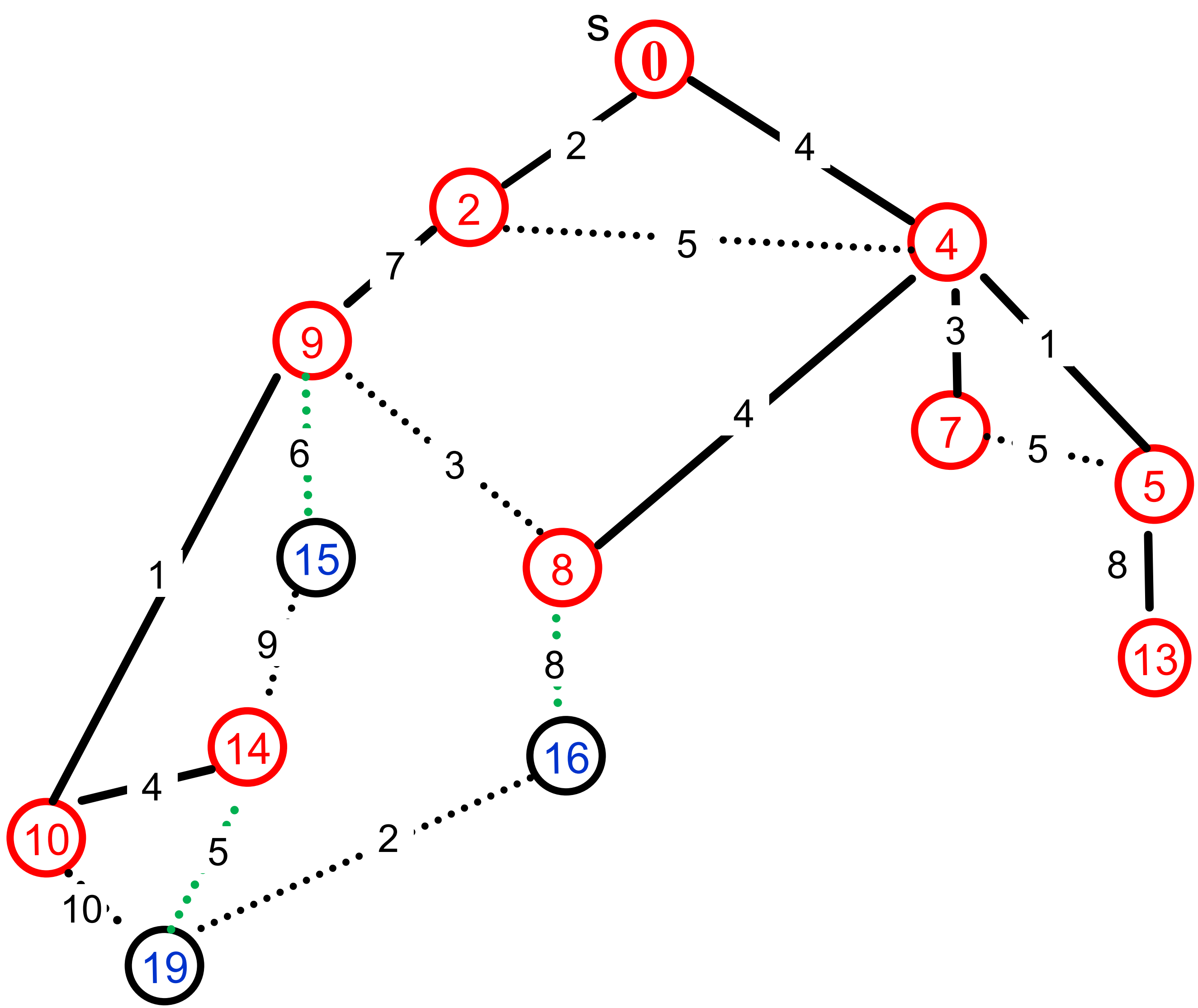

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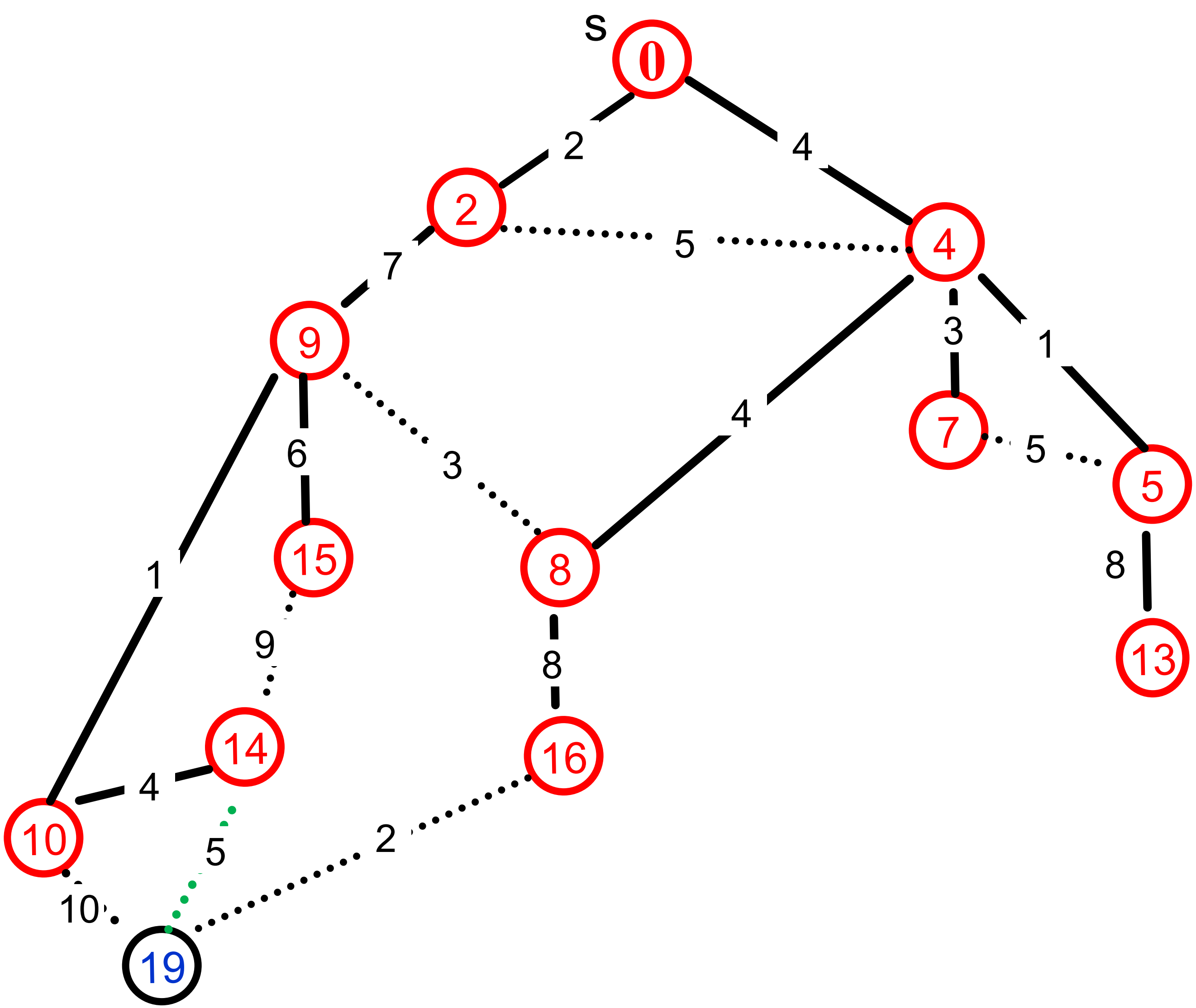

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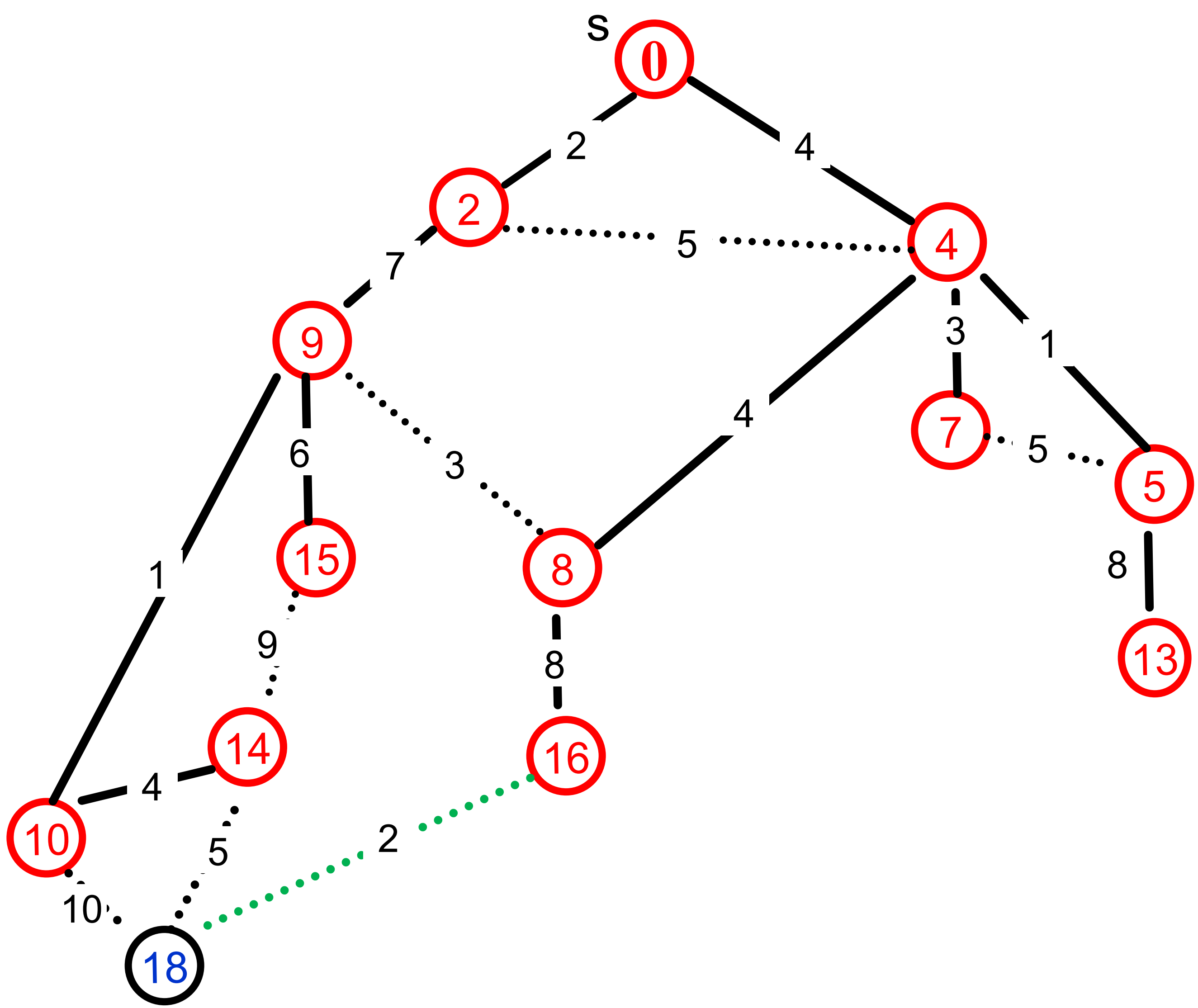

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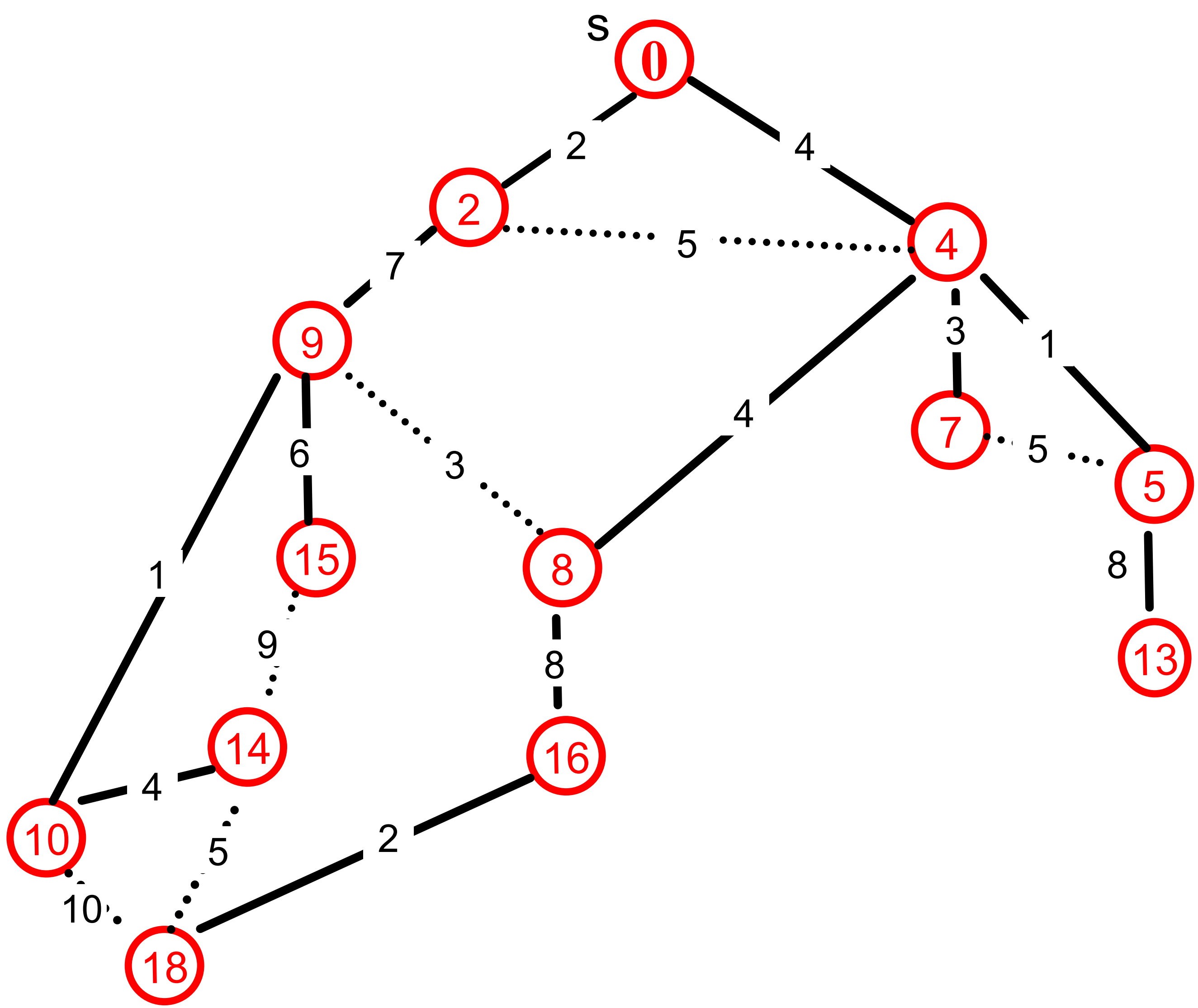


Disjkstra's Algorithm: Correctness

Prove by induction that throughout the algorithm, for any $u \in S$, the path P_u in the shortest from s to u.

Base Case: This is always true when $S = \{s\}$.

IH: Suppose $|S| = k$ and the claim holds for S

IS: Say v is the k+1-st vertex that we add to S. Let {u,v} be last edge on P_v.

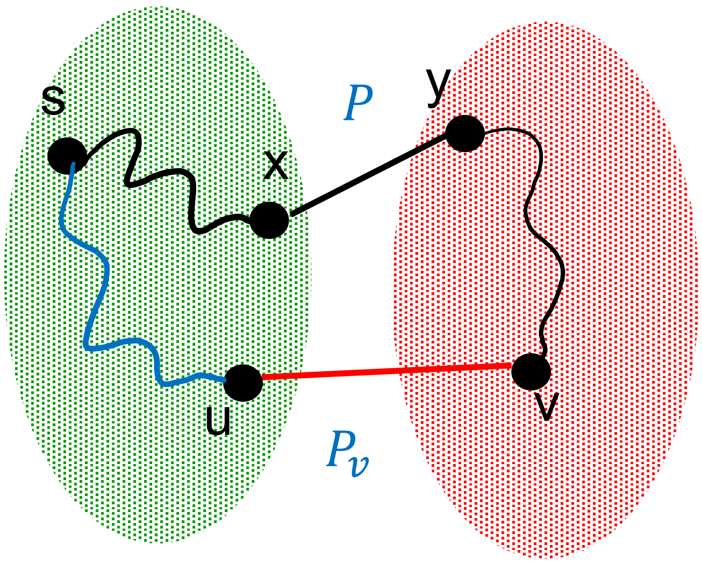
If P_v is not the shortest path there is a path P to s which is shorter.

Consider the first time that P leaves S (with edge {x,y}).

S -> x has weight (at least) d(x)

So, $c(P) \geq d(x) + c_{x,y} \geq d(v) = c(P_v)$.

A contradiction.

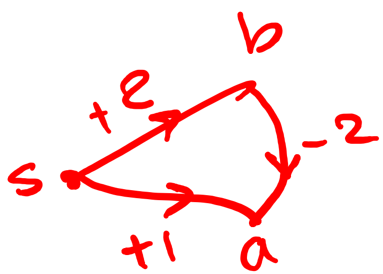


Remarks on Dijkstra's Algorithm

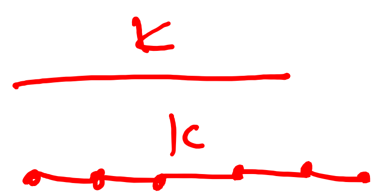
- Algorithm also produces a tree of shortest paths to s following Parent links
- Algorithm works on directed graph (with nonnegative weights)

- The algorithm fails with negative edge weights.
  - e.g., some airline tickets

Why does it fail?



- Dijkstra's algorithm is similar to BFS:
  - Subtitute every edge with $c_e = k$ with a path of length k, then run BFS.



dijk = BFS.

Implementing Dijkstra's Algorithm

Priority Queue: Elements each with an associated key Operations

- Insert
- Find-min
  - Return the element with the smallest key
- Delete-min
  - Return the element with the smallest key and delete it from the data structure
- Decrease-key
  - Decrease the key value of some element

Implementations

Arrays:

- O(n) time find/delete-min,
- O(1) time insert/decrease key

Binary Heaps:

- O(log n) time insert/decrease-key/delete-min,
- O(1) time find-min

Dijkstra's Algorithm

Runs in O((n+m)log n).

```
Dijkstra(G, c, s) {
   foreach (v ∈ V) d[v] ← ∞ //This is the key of node v
   d[s] ← 0
   foreach (v ∈ V) insert v onto a priority queue Q
   Initialize set of explored nodes S ← {s}

   while (Q is not empty) {   O(1)
      u ← delete min element from Q
      S ← S ∪ { u }
      foreach (edge e = (u, v) incident to u)
          if ((v ∉ S) and (d[u]+c_e < d[v]))
             d[v] ← d[u] + c_e
             Decrease key of v to d[v].
             Parent(v) ← u
}
```

$O(n)$ of delete min, each in O(log n)

$O(m)$ of decrease key, each runs in $O(\log n)$

Summary (Greedy Algorithms)

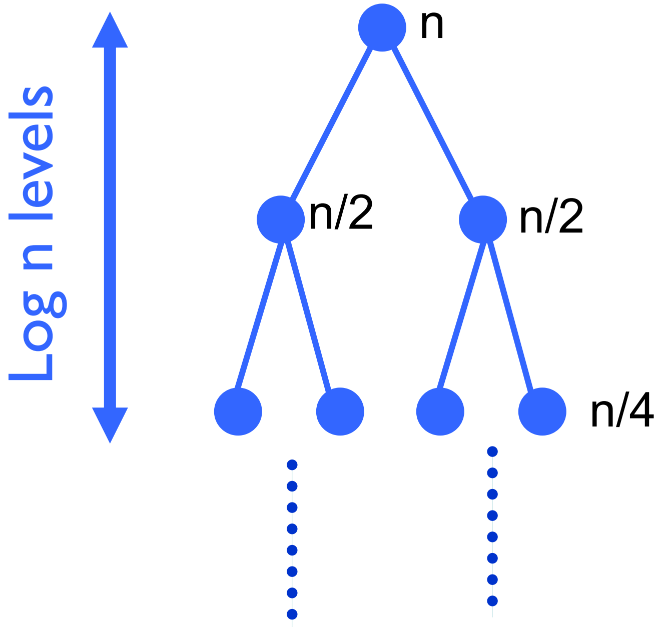
- Greedy Stays Ahead: Interval Scheduling, Dijkstra's algorithm
- Structural: Interval Partitioning
- Exchange Arguments: MST, Kruskal's Algorithm,
- Data Structures: Union Find, Heap

Divide and Conquer Approach

Divide and Conquer

Similar to algorithm design by induction, we reduce a problem to several subproblems.

Typically, each sub-problem is at most a constant fraction of the size of the original problem



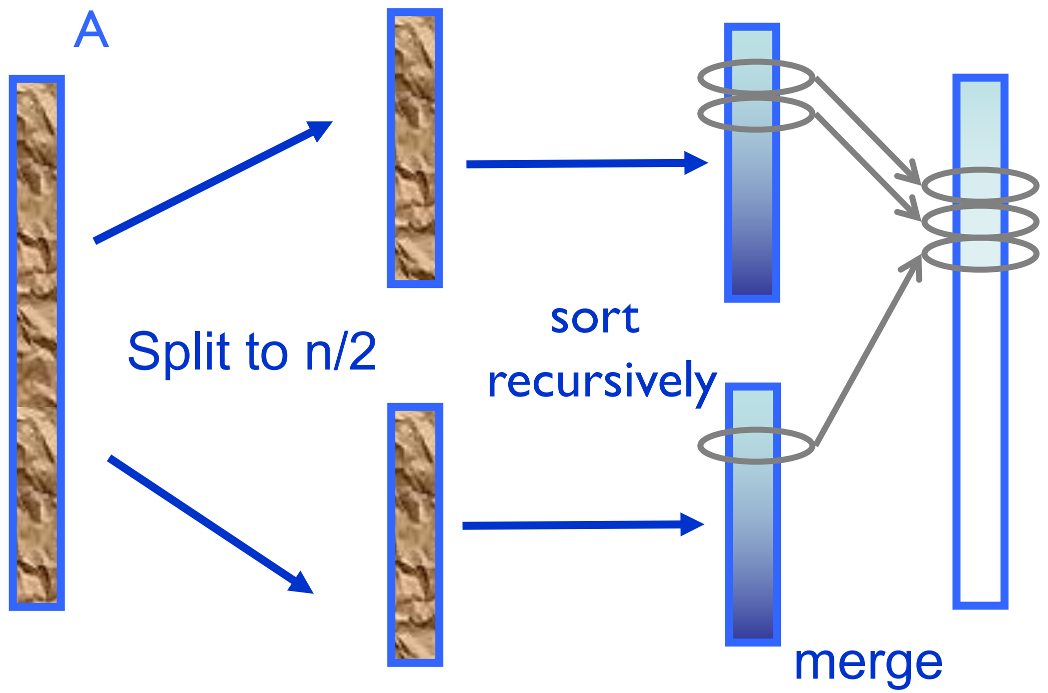
Recursively solve each subproblem

Merge the solutions

Examples:

- Mergesort, Binary Search, Strassen's Algorithm,

A Classical Example: Merge Sort



Why Balanced Partitioning?

An alternative "divide & conquer" algorithm:

- Split into n-1 and 1
- Sort each sub problem
- Merge them

Runtime

$$T(n) = T(n-1) + T(1) + n$$

Solution:

$$T(n) = n + T(n-1) + T(1)$$

$$= n + n - 1 + \ T(n-2)$$

$$= n + n - 1 + n - 2 + T(n-3)$$

$$= n + n - 1 + n - 2 + \cdots + 1 = O(n^2)$$

D&C: The Key Idea

Suppose we've already invented Bubble-Sort, and we know it takes n^2

Try just one level of divide & conquer:

Bubble-Sort(first n/2 elements)

Bubble-Sort(last n/2 elements)

Merge results

Time: $2\,\underbrace{T(n/2)}_{n^2/4} + n = n^2/2 + n \ll n^2$

Almost twice as fast!

D&C in a nutshell

D&C approach

- "the more dividing and conquering, the better"
  - Two levels of D&C would be almost 4 times faster, 3 levels almost 8, etc., even though overhead is growing.
  - Best is usually full recursion down to a small constant size (balancing "work" vs "overhead").

  In the limit: you've just rediscovered mergesort!
- Even unbalanced partitioning is good, but less good
  - Bubble-sort improved with a 0.1/0.9 split:

    $$(.1n)^2 + (.9n)^2 + n = .82n^2 + n$$

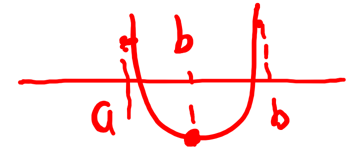
    The 18% savings compounds significantly if you carry recursion to more levels, actually giving $O(n \log n)$, but with a bigger constant.
- This is why Quicksort with random splitter is good – badly unbalanced splits are rare, and not instantly fatal.

Finding the Root of a Function

Finding the Root of a Function

Given a continuous function f and two points a < b such that

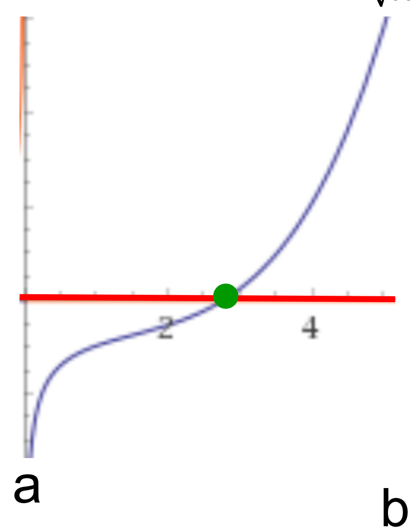
$$f(a) \leq 0$$
$$f(b) \geq 0$$



Find an approximate root of f (a point c where $f(c) = 0$).

f has a root in $[a, b]$ by intermediate value theorem

$$\mathrm{f}(x) = \sin(x) - \frac{100}{\sqrt{x}} + x^4$$



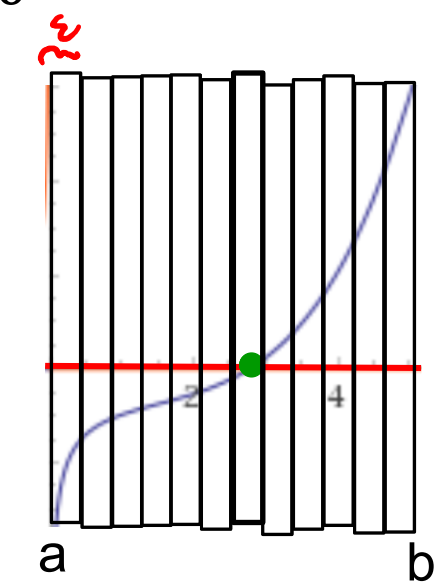
Note that roots of f may be irrational,
So, we want to approximate
the root with an arbitrary precision!

A Naiive Approch

Suppose we want ϵ approximation to a root.

Divide [a,b] into $n = \frac{b-a}{\epsilon}$ intervals. For each interval check

$$f(x) \leq 0, f(x+\epsilon) \geq 0$$



This runs in time $O(n) = O(\frac{b-a}{\epsilon})$

Can we do faster?

D&C Approach (Based on Binary Search)

Bisection(a,b, ε)

  if $(b - a) < \epsilon$ then

    return (a)

  else

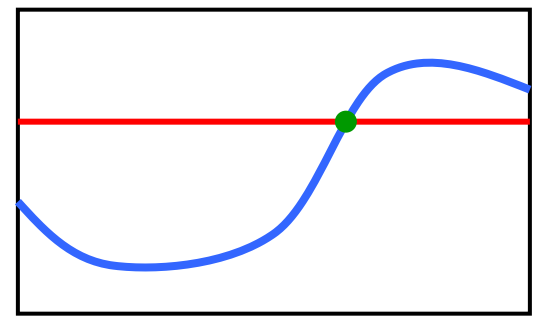
    $m \leftarrow (a + b)/2$

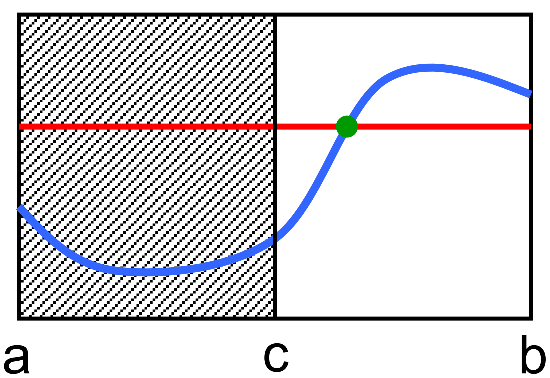
    if $f(m) \leq 0$ then

      return(Bisection(c, b, ε))

    else

      return(Bisection(a, c, ε))





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Time Analysis

Let $n = \frac{a-b}{\epsilon}$

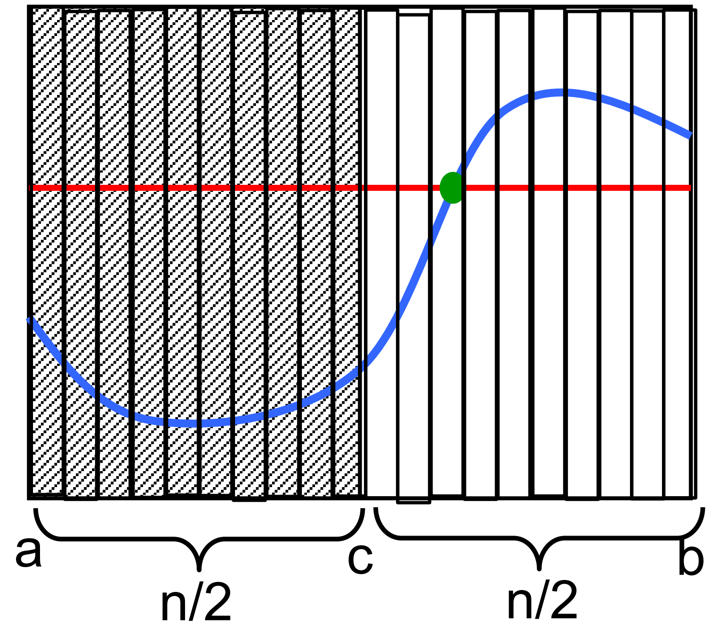
And $c = (a + b)/2$

Always half of the intervals lie to the left and half lie to the right of c

So,

$$T(n) = T\left(\frac{n}{2}\right) + O(1)$$

i.e., $T(n) = O(\log n) = O(\log \frac{a-b}{\epsilon})$



Correctness Proof

P(k) = "For any a, b such that $k\epsilon \leq |a - b| \leq (k + 1)\epsilon$ if $f(a)f(b) < 0$, then we find an ϵ approx to a root using $\log k$ queries to f"

Base Case: P(1): Output $a + \epsilon$

IH: Assume P(k).

IS: Show P(2k). Consider an arbitrary a, b s.t.,

$$2k\epsilon \leq |a - b| < (2k + 1)\epsilon$$

If $f(a + k\epsilon) = 0$ output $a + k\epsilon$.

If $f(\mathrm{a})f(a + k\epsilon) < 0$, solve for interval $a, a + k\epsilon$ using log(k) queries to f.

Otherwise, we must have $f(b)f(a + k\epsilon) < 0$ since $f(a)f(b) < 0$ and $f(a)f(a + k\epsilon) \geq 0$. Solve for interval $a + k\epsilon, b$.

Overall we use at most log(k)+1=log(2k) queries to *f*.