

## Correctness of Dijkstra.

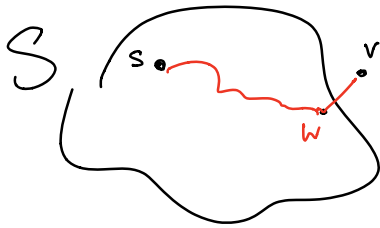
**Fact:** If  $P_v$  is shortest path from  $s$  to  $v$ .  
 $u$  is on the path. Then the part from  $s$  to  $u$  (in  $P_v$ )  
 is the shortest path from  $s$  to  $u$ .



$u$  is right before  $v$ .

If I know  $u$ , then first find  $P_u$  (shortest path to  $u$ )  
 and then add  $(u, v)$ .

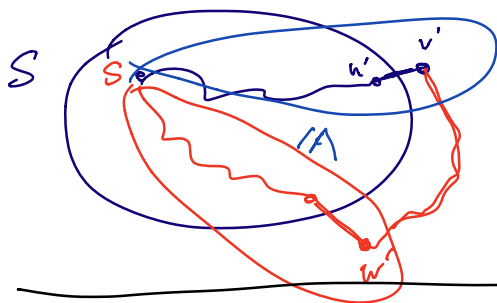
In Dijkstra Assume  $u$  have found shortest  $\phi$  to  $S$ .



If I know the vertex **before**  $v$   
 is in  $S \Rightarrow P_v = P_u + (u, v)$   
~~So just add  $v$  with this shortest path~~

$$P_v = \min_{w \in S} c(P_w) + c_{w,v}$$

Q: Find  $v \notin S$ , s.t. vertex before  $v$  in  $P_v$  (shortest path to  $v$ )  
 is in  $S$ .



The candidate for  $v$  is the  
 closest vertex to  $s$ .

Formal Pf. Induction.  $P(0) = S_{\text{supp}}$  we have found shortest path to a  $S$ , where  $|S|=k$ .

$$P(1) = |S| = 1, S = \{s\} \checkmark$$

IH:  $P(k)$ .

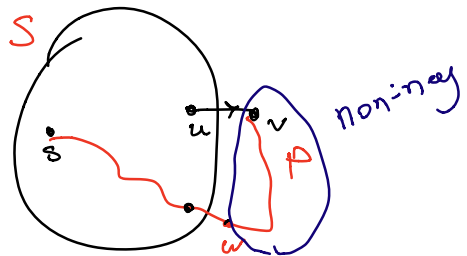
IS: Goal is to Prove  $P(k+1)$ . Assume that we have found shortest path to  $S$  s.t.  $|S|=k$ .

Assume  $v$  is the next vertex that we add,  $d[v] = d[u] + c_{uv}$

Supp  $c(P_v) < d[v]$  for contradiction

Let  $w$  be the first time shortest path

$P_v$  leaves  $S$



$$d[v] > c(P_v) \geq c(P_w) = d[w]$$

cost  
of edge  
edge  $Tr$

So  $d[v] > d[w]$  contradiction for choosing  $v$ .

$$d[v] = \min_{u \in S} c(P_u) + c(u,v)$$