Department of Computer Science and Engineering
CSE 421, Winter 2019

Homework 8, Due Wednesday, March 15, 2019

## Problem 1 (10 points):

Give an algorithm, which given a directed graph $G=(V, E)$, with vertices $s, t \in V$ and an integer $k$, determines the number of paths from $s$ to $t$ of length $k$. Your algorithm should be polynomial in $k,|V|$ and $|E|$.

## Problem 2 (10 points):

Give an $O\left(n^{2}\right)$ algorithm for finding the longest strictly increasing subsequence of a sequence of $n$ integers. (Note that this problem can be solved in $O(n \log n)$ time by a non-dynamic programming style algorithm, but you do not need find it. Use of dynamic programming for this problem is recommended, but not required.)

## Problem 3 (10 points):

Decide whether you think the following statement is true or false. If it is true, give a short explanation. If it is false, give a counterexample.
Let $G$ be an arbitrary flow network, with a source $s$, a sink $t$ and a positive integer capacity $c_{e}$ on every edge $e$; and let $(A, B)$ be a minimum $s-t$ cut with respect to these capacities $\left\{c_{e}: e \in E\right\}$. Now suppose we add 1 to every capacity; then $(A, B)$ is still a minimum $s-t$ cut with respect to these new capacities $\left\{1+c_{e}: e \in E\right\}$.

## Problem 4 (10 points):

You are given a flow network with unit-capacity edges: it consists of a directed graph $G=(V, E)$, a source $s \in V$, and a $\operatorname{sink} t \in V$; and $c_{e}=1$ for every $e \in E$. You are also given a parameter $k$.

The goal is to delete $k$ edges so as to reduce the maximum $s-t$ flow in $G$ by as much as possible. In other words, you should find a set of edges $F \subseteq E$ so that $|F|=k$ and the maximum $s-t$ flow in $G^{\prime}=(V, E-F)$ is as small as possible subject to this.

Give a polynomial time algorithm to solve this problem, and justify that your algorithm is correct.

## Problem 5 (10 points):

In a standard $s-t$ Maximum-Flow Problem, we assume edges have capacities, and there is no limit on how much flow is allowed to pass through a node. In this problem, we consider the variant of the Maximum-Flow problem with node capacities.

Let $G=(V, E)$ be a directed graph, with source $s \in V$, $\operatorname{sink} t \in V$, and nonnegative node capacities $\left\{c_{v} \geq 0\right\}$ for each $v \in V$. Given a flow $f$ in this graph, the flow through a node $v$ is defined as $f^{\mathrm{in}}(v)$. We say that a flow is feasible if it satisfies the usual flow-conservation constraints and the node-capacity constraints: $f^{\text {in }}(v) \leq c_{v}$ for all nodes.

Give a polynomial-time algorithm to find an $s-t$ maximum flow in such a node-capacitated network. Justify the correctness of your algorithm.

## Problem 6 (10 points):

Consider the following flow graph. Find a maximum flow.

a) What is the value of the maximum flow? Indicate the value of flow on each edge.
b) Prove that your flow is maximum.

## Problem 7 ( 10 points):

Suppose that you are a wealthy individual who wishes to make campaign donations to candidates $C_{1}, \ldots, C_{n}$. You want to donate $b_{i}$ dollars to candidate $C_{i}$. However, campaign finance laws prohibit you from donating to these candidates, so you decide that you will have your friends, $F_{1}, \ldots, F_{m}$ help you out by laundering the donations. You determine, for each friend $F_{i}$ and candidate $C_{j}$, the maximum amount $a_{i j}$ that $F_{i}$ is willing to contribute to $C_{j}$ for you. Since your friends are not fully trustworthy, you decide to cap the amount of money that you will give to each of them, so you set a bound of $M_{i}$ for the amount of money you will give to $F_{i}$ to distribute.

Give a network flow based algorithm which determines how the money is distributed. The algorithm should indicate how much money each friend gives to each candidate. Justify the correctness of your algorithm.

